

The Budan–Fourier Theorem and Counting Real Roots with Multiplicity

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Abstract

This entry is mainly about counting and approximating real roots (of a polynomial) with multiplicity. We have first formalised the Budan–Fourier theorem: given a polynomial with real coefficients, we can calculate sign variations on Fourier sequences to over-approximate the number of real roots (counting multiplicity) within an interval. When all roots are known to be real, the over-approximation becomes tight: we can utilise this theorem to count real roots exactly. It is also worth noting that Descartes’ rule of sign is a direct consequence of the Budan–Fourier theorem, and has been included in this entry. In addition, we have extended previous formalised Sturm’s theorem to count real roots with multiplicity, while the original Sturm’s theorem only counts distinct real roots. Compared to the Budan–Fourier theorem, our extended Sturm’s theorem always counts roots exactly but may suffer from greater computational cost.

Many problems in real algebraic geometry is about counting or approximating roots of a polynomial. Previous formalised results are mainly about counting distinct real roots (i.e. Sturm’s theorem in Isabelle/HOL [5, 2], HOL Light [4], PVS [9] and Coq [8]) and limited support for multiple real roots (i.e. Descartes’ rule of signs in Isabelle/HOL [3], HOL Light and Proof-Power¹). In comparison, this entry provides more comprehensive support for reasoning about multiple real roots.

The main motivation of this entry is to cope with the roots-on-the-border issue when counting complex roots [7, 6], but the results here should be beneficial to other developments.

Our proof of the Budan–Fourier theorem mainly follows Theorem 2.35 in the book by Basu et al. [1] and that of the extended Sturm’s theorem is inspired by Theorem 10.5.6 in Rahman and Schmeisser’s book [10].

¹According to Freek Wiedijk’s “Formalising 100 Theorems” (<http://www.cs.ru.nl/~freek/100/index.html>)

1 Misc results for polynomials and sign variations

```
theory BF-Misc imports  
  HOL-Computational-Algebra.Polynomial-Factorial  
  HOL-Computational-Algebra.Fundamental-Theorem-Algebra  
  Sturm-Tarski.Sturm-Tarski  
begin
```

1.1 Induction on polynomial roots

```
lemma poly-root-induct-alt [case-names 0 no-roots root]:  
  fixes  $p :: 'a :: \text{idom } \text{poly}$   
  assumes  $Q\ 0$   
  assumes  $\bigwedge p. (\bigwedge a. \text{poly } p\ a \neq 0) \implies Q\ p$   
  assumes  $\bigwedge a\ p. Q\ p \implies Q\ ([:-a, 1:] * p)$   
  shows  $Q\ p$   
<proof>
```

1.2 Misc

```
lemma lead-coeff-pderiv:  
  fixes  $p :: 'a :: \{\text{comm-semiring-1, semiring-no-zero-divisors, semiring-char-0}\}$  poly  
  shows  $\text{lead-coeff } (\text{pderiv } p) = \text{of-nat } (\text{degree } p) * \text{lead-coeff } p$   
<proof>
```

```
lemma gcd-degree-le-min:  
  assumes  $p \neq 0\ q \neq 0$   
  shows  $\text{degree } (\text{gcd } p\ q) \leq \min (\text{degree } p) (\text{degree } q)$   
<proof>
```

```
lemma lead-coeff-normalize-field:  
  fixes  $p :: 'a :: \{\text{field, semidom-divide-unit-factor}\}$  poly  
  assumes  $p \neq 0$   
  shows  $\text{lead-coeff } (\text{normalize } p) = 1$   
<proof>
```

```
lemma smult-normalize-field-eq:  
  fixes  $p :: 'a :: \{\text{field, semidom-divide-unit-factor}\}$  poly  
  shows  $p = \text{smult } (\text{lead-coeff } p) (\text{normalize } p)$   
<proof>
```

```
lemma lead-coeff-gcd-field:  
  fixes  $p\ q :: 'a :: \text{field-gcd } \text{poly}$   
  assumes  $p \neq 0 \vee q \neq 0$   
  shows  $\text{lead-coeff } (\text{gcd } p\ q) = 1$   
<proof>
```

```
lemma poly-gcd-0-iff:  
   $\text{poly } (\text{gcd } p\ q)\ x = 0 \iff \text{poly } p\ x = 0 \wedge \text{poly } q\ x = 0$   
<proof>
```

lemma *degree-eq-oneE*:
fixes $p :: 'a::zero\ poly$
assumes $degree\ p = 1$
obtains $a\ b$ **where** $p = [:a,b:]\ b \neq 0$
 $\langle proof \rangle$

1.3 More results about sign variations (i.e. *changes*)

lemma *changes-0[simp]*: $changes\ (0\ \#\ xs) = changes\ xs$
 $\langle proof \rangle$

lemma *changes-Cons*: $changes\ (x\ \#\ xs) =$ (if $filter\ (\lambda x. x \neq 0)\ xs = []$ then
 0
else if $x * hd\ (filter\ (\lambda x. x \neq 0)\ xs) < 0$ then
 $1 + changes\ xs$
else $changes\ xs$)
 $\langle proof \rangle$

lemma *changes-filter-eq*:
 $changes\ (filter\ (\lambda x. x \neq 0)\ xs) = changes\ xs$
 $\langle proof \rangle$

lemma *changes-filter-empty*:
assumes $filter\ (\lambda x. x \neq 0)\ xs = []$
shows $changes\ xs = 0$ $changes\ (a\ \#\ xs) = 0$ $\langle proof \rangle$

lemma *changes-append*:
assumes $xs \neq [] \wedge ys \neq [] \longrightarrow (last\ xs = hd\ ys \wedge last\ xs \neq 0)$
shows $changes\ (xs@ys) = changes\ xs + changes\ ys$
 $\langle proof \rangle$

lemma *changes-drop-dup*:
assumes $xs \neq []\ ys \neq [] \longrightarrow last\ xs = hd\ ys$
shows $changes\ (xs@ys) = changes\ (xs@tl\ ys)$
 $\langle proof \rangle$

lemma *Im-poly-of-real*:
 $Im\ (poly\ p\ (of-real\ x)) = poly\ (map-poly\ Im\ p)\ x$
 $\langle proof \rangle$

lemma *Re-poly-of-real*:
 $Re\ (poly\ p\ (of-real\ x)) = poly\ (map-poly\ Re\ p)\ x$
 $\langle proof \rangle$

1.4 More about *map-poly* and *of-real*

lemma *of-real-poly-map-pCons*[simp]: *map-poly of-real* (pCons a p) = pCons (of-real a) (map-poly of-real p)
<proof>

lemma *of-real-poly-map-plus*[simp]: *map-poly of-real* (p + q) = *map-poly of-real* p + *map-poly of-real* q
<proof>

lemma *of-real-poly-map-smult*[simp]: *map-poly of-real* (smult s p) = smult (of-real s) (map-poly of-real p)
<proof>

lemma *of-real-poly-map-mult*[simp]: *map-poly of-real* (p*q) = *map-poly of-real* p * *map-poly of-real* q
<proof>

lemma *of-real-poly-map-poly*:
of-real (poly p x) = poly (map-poly of-real p) (of-real x)
<proof>

lemma *of-real-poly-map-power*: *map-poly of-real* (pⁿ) = (map-poly of-real p)ⁿ
<proof>

lemma *of-real-poly-eq-iff* [simp]: *map-poly of-real* p = *map-poly of-real* q \longleftrightarrow p = q
<proof>

lemma *of-real-poly-eq-0-iff* [simp]: *map-poly of-real* p = 0 \longleftrightarrow p = 0
<proof>

1.5 More about *order*

lemma *order-multiplicity-eq*:
assumes p \neq 0
shows order a p = multiplicity [-a,1:] p
<proof>

lemma *order-gcd*:
assumes p \neq 0 q \neq 0
shows order x (gcd p q) = min (order x p) (order x q)
<proof>

lemma *order-linear*[simp]: order x [-a,1:] = (if x=a then 1 else 0)
<proof>

lemma *map-poly-order-of-real*:

assumes $p \neq 0$
shows $\text{order } (of\text{-real } t) (\text{map-poly } of\text{-real } p) = \text{order } t \ p$ $\langle proof \rangle$

lemma *order-pcompose*:

assumes $p \text{compose } p \ q \neq 0$
shows $\text{order } x \ (p \text{compose } p \ q) = \text{order } x \ (q - [:poly \ q \ x:]) * \text{order } (poly \ q \ x) \ p$
 $\langle proof \rangle$

1.6 Polynomial roots / zeros

definition *roots-within*:: $'a::comm\text{-semiring-0}$ $poly \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set}$ **where**
 $\text{roots-within } p \ s = \{x \in s. \text{poly } p \ x = 0\}$

abbreviation *roots*:: $'a::comm\text{-semiring-0}$ $poly \Rightarrow 'a \text{ set}$ **where**
 $\text{roots } p \equiv \text{roots-within } p \ UNIV$

lemma *roots-def*: $\text{roots } p = \{x. \text{poly } p \ x = 0\}$
 $\langle proof \rangle$

lemma *roots-within-empty[simp]*:
 $\text{roots-within } p \ \{\} = \{\}$ $\langle proof \rangle$

lemma *roots-within-0[simp]*:
 $\text{roots-within } 0 \ s = s$ $\langle proof \rangle$

lemma *roots-withinI[intro,simp]*:
 $\text{poly } p \ x = 0 \Longrightarrow x \in s \Longrightarrow x \in \text{roots-within } p \ s$
 $\langle proof \rangle$

lemma *roots-within-iff[simp]*:
 $x \in \text{roots-within } p \ s \longleftrightarrow \text{poly } p \ x = 0 \wedge x \in s$
 $\langle proof \rangle$

lemma *roots-within-union*:
 $\text{roots-within } p \ A \cup \text{roots-within } p \ B = \text{roots-within } p \ (A \cup B)$
 $\langle proof \rangle$

lemma *roots-within-times*:
fixes $s::'a::\{semiring\text{-no-zero-divisors}, comm\text{-semiring-0}\}$ set
shows $\text{roots-within } (p * q) \ s = \text{roots-within } p \ s \cup \text{roots-within } q \ s$
 $\langle proof \rangle$

lemma *roots-within-gcd*:
fixes $s::'a::\{factorial\text{-ring-gcd}, semiring\text{-gcd-mult-normalize}\}$ set
shows $\text{roots-within } (gcd \ p \ q) \ s = \text{roots-within } p \ s \cap \text{roots-within } q \ s$
 $\langle proof \rangle$

lemma *roots-within-inter*:
 $NO\text{-MATCH } UNIV \ s \Longrightarrow \text{roots-within } p \ s = \text{roots } p \ \cap \ s$

<proof>

lemma *roots-within-roots[simp]*:
roots-within p s \subseteq *roots p*
<proof>

lemma *finite-roots[simp]*:
fixes *p* :: 'a::idom poly
shows *p* $\neq 0 \implies$ *finite (roots-within p s)*
<proof>

lemma *roots-within-pCons-1-iff*:
fixes *a*::'a::idom
shows *roots-within* $[-a, 1:] s = (\text{if } a \in s \text{ then } \{a\} \text{ else } \{\})$
roots-within $[a, -1:] s = (\text{if } a \in s \text{ then } \{a\} \text{ else } \{\})$
<proof>

lemma *roots-within-uminus[simp]*:
fixes *p* :: 'a::comm-ring poly
shows *roots-within* $(- p) s = \text{roots-within } p s$
<proof>

lemma *roots-within-smult*:
fixes *a*::'a:: $\{\text{semiring-no-zero-divisors, comm-semiring-0}\}$
assumes *a* $\neq 0$
shows *roots-within* $(\text{smult } a p) s = \text{roots-within } p s$
<proof>

1.7 Polynomial roots counting multiplicities.

definition *roots-count*::'a::idom poly \Rightarrow 'a set \Rightarrow nat **where**
roots-count p s = $(\sum r \in \text{roots-within } p s. \text{order } r p)$

lemma *roots-count-empty[simp]*:*roots-count p* $\{\} = 0$
<proof>

lemma *roots-count-times*:
fixes *s* :: 'a::idom set
assumes *p***q* $\neq 0$
shows *roots-count* $(p*q) s = \text{roots-count } p s + \text{roots-count } q s$
<proof>

lemma *roots-count-power-n-n*:
shows *roots-count* $([: - a, 1:]^n) s = (\text{if } a \in s \wedge n > 0 \text{ then } n \text{ else } 0)$
<proof>

lemma *degree-roots-count*:
fixes *p*::complex poly
shows *degree p* = *roots-count p UNIV*

<proof>

lemma *proots-count-smult*:
 fixes $a::'a::\{\text{semiring-no-zero-divisors},\text{idom}\}$
 assumes $a\neq 0$
 shows $\text{proots-count } (\text{smult } a \ p) \ s = \text{proots-count } p \ s$
<proof>

lemma *proots-count-pCons-1-iff*:
 fixes $a::'a::\text{idom}$
 shows $\text{proots-count } [:-a,1:] \ s = (\text{if } a\in s \text{ then } 1 \text{ else } 0)$
<proof>

lemma *proots-count-uminus[simp]*:
 $\text{proots-count } (- \ p) \ s = \text{proots-count } p \ s$
<proof>

lemma *card-proots-within-leq*:
 assumes $p\neq 0$
 shows $\text{proots-count } p \ s \geq \text{card } (\text{proots-within } p \ s)$ *<proof>*

lemma *proots-count-0-imp-empty*:
 assumes $\text{proots-count } p \ s = 0 \ p\neq 0$
 shows $\text{proots-within } p \ s = \{\}$
<proof>

lemma *proots-count-leq-degree*:
 assumes $p\neq 0$
 shows $\text{proots-count } p \ s \leq \text{degree } p$ *<proof>*

lemma *proots-count-union-disjoint*:
 assumes $A \cap B = \{\} \ p\neq 0$
 shows $\text{proots-count } p \ (A \cup B) = \text{proots-count } p \ A + \text{proots-count } p \ B$
<proof>

lemma *proots-count-cong*:
 assumes $\text{order-eq}:\forall x\in s. \text{order } x \ p = \text{order } x \ q$ **and** $p\neq 0$ **and** $q\neq 0$
 shows $\text{proots-count } p \ s = \text{proots-count } q \ s$ *<proof>*

lemma *proots-count-of-real*:
 assumes $p\neq 0$
 shows $\text{proots-count } (\text{map-poly } \text{of-real } p) \ ((\text{of-real}::\Rightarrow'a::\{\text{real-algebra-1},\text{idom}\}) \ s)$
 $= \text{proots-count } p \ s$
<proof>

lemma *roots-pcompose*:
fixes $p q :: 'a :: \text{field}$ *poly*
assumes $p \neq 0$ *degree* $q = 1$
shows *roots-count* (*pcompose* p q) $s = \text{roots-count } p$ (*poly* q ' s)
 $\langle \text{proof} \rangle$

1.8 Composition of a polynomial and a rational function

definition *fcompose* :: $'a :: \text{field}$ *poly* $\Rightarrow 'a$ *poly* $\Rightarrow 'a$ *poly* $\Rightarrow 'a$ *poly* **where**
fcompose p q $r = \text{fst}$ (*fold-coeffs* (λa (c, d). ($d * [:a:] + q * c, r * d$)) p ($0, 1$))

lemma *fcompose-0* [*simp*]: *fcompose* 0 q $r = 0$
 $\langle \text{proof} \rangle$

lemma *fcompose-const*[*simp*]: *fcompose* $[:a:]$ q $r = [:a:]$
 $\langle \text{proof} \rangle$

lemma *fcompose-pCons*:
fcompose (*pCons* a p) $q1$ $q2 = \text{smult } a$ ($q2^{\wedge}(\text{degree } (pCons$ a $p))$) + $q1 * \text{fcompose}$
 p $q1$ $q2$
 $\langle \text{proof} \rangle$

lemma *fcompose-uminus*:
fcompose ($-p$) q $r = - \text{fcompose } p$ q r
 $\langle \text{proof} \rangle$

lemma *fcompose-add-less*:
assumes *degree* $p1 > \text{degree } p2$
shows *fcompose* ($p1 + p2$) $q1$ $q2$
 $= \text{fcompose } p1$ $q1$ $q2 + q2^{\wedge}(\text{degree } p1 - \text{degree } p2) * \text{fcompose } p2$ $q1$ $q2$
 $\langle \text{proof} \rangle$

lemma *fcompose-add-eq*:
assumes *degree* $p1 = \text{degree } p2$
shows $q2^{\wedge}(\text{degree } p1 - \text{degree } (p1 + p2)) * \text{fcompose } (p1 + p2)$ $q1$ $q2$
 $= \text{fcompose } p1$ $q1$ $q2 + \text{fcompose } p2$ $q1$ $q2$
 $\langle \text{proof} \rangle$

lemma *fcompose-add-const*:
fcompose ($[:a:] + p$) $q1$ $q2 = \text{smult } a$ ($q2^{\wedge} \text{degree } p$) + *fcompose* p $q1$ $q2$
 $\langle \text{proof} \rangle$

lemma *fcompose-smult*: *fcompose* (*smult* a p) $q1$ $q2 = \text{smult } a$ (*fcompose* p $q1$ $q2$)
 $\langle \text{proof} \rangle$

lemma *fcompose-mult*: *fcompose* ($p1 * p2$) $q1$ $q2 = \text{fcompose } p1$ $q1$ $q2 * \text{fcompose}$
 $p2$ $q1$ $q2$
 $\langle \text{proof} \rangle$

lemma *fcompose-poly*:
assumes *poly q2 x ≠ 0*
shows $\text{poly } p (\text{poly } q1 \ x / \text{poly } q2 \ x) = \text{poly } (\text{fcompose } p \ q1 \ q2) \ x / \text{poly } (q2^{\wedge}(\text{degree } p)) \ x$
<proof>

lemma *poly-fcompose*:
assumes *poly q2 x ≠ 0*
shows $\text{poly } (\text{fcompose } p \ q1 \ q2) \ x = \text{poly } p (\text{poly } q1 \ x / \text{poly } q2 \ x) * (\text{poly } q2 \ x)^{\wedge}(\text{degree } p)$
<proof>

lemma *poly-fcompose-0-denominator*:
assumes *poly q2 x = 0*
shows $\text{poly } (\text{fcompose } p \ q1 \ q2) \ x = \text{poly } q1 \ x^{\wedge} \text{degree } p * \text{lead-coeff } p$
<proof>

lemma *fcompose-0-denominator*: $\text{fcompose } p \ q1 \ 0 = \text{smult } (\text{lead-coeff } p) (q1^{\wedge} \text{degree } p)$
<proof>

lemma *fcompose-nzero*:
fixes *p::'a::field poly*
assumes *p ≠ 0 and q2 ≠ 0 and nconst:∀ c. q1 ≠ smult c q2*
and *infi:infinite (UNIV::'a set)*
shows $\text{fcompose } p \ q1 \ q2 \neq 0$ *<proof>*

1.9 Bijection (*bij-betw*) and the number of polynomial roots

lemma *proots-fcompose-bij-eq*:
fixes *p::'a::field poly*
assumes *bij:bij-betw (λx. poly q1 x / poly q2 x) A B and p ≠ 0*
and *nzero:∀ x ∈ A. poly q2 x ≠ 0*
and *max-deg: max (degree q1) (degree q2) ≤ 1*
and *nconst:∀ c. q1 ≠ smult c q2*
and *infi:infinite (UNIV::'a set)*
shows $\text{proots-count } p \ B = \text{proots-count } (\text{fcompose } p \ q1 \ q2) \ A$
<proof>

lemma *proots-card-fcompose-bij-eq*:
fixes *p::'a::field poly*
assumes *bij:bij-betw (λx. poly q1 x / poly q2 x) A B and p ≠ 0*
and *nzero:∀ x ∈ A. poly q2 x ≠ 0*
and *max-deg: max (degree q1) (degree q2) ≤ 1*
and *nconst:∀ c. q1 ≠ smult c q2*
and *infi:infinite (UNIV::'a set)*
shows $\text{card } (\text{proots-within } p \ B) = \text{card } (\text{proots-within } (\text{fcompose } p \ q1 \ q2) \ A)$
<proof>

lemma *proots-pcompose-bij-eq*:
fixes $p::'a::idom\ poly$
assumes $bij:bij\ betw\ (\lambda x. poly\ q\ x)\ A\ B$ **and** $p \neq 0$
and $q\ deg: degree\ q = 1$
shows $proots\ count\ p\ B = proots\ count\ (p\ \circ_p\ q)\ A$ $\langle proof \rangle$

lemma *proots-card-pcompose-bij-eq*:
fixes $p::'a::idom\ poly$
assumes $bij:bij\ betw\ (\lambda x. poly\ q\ x)\ A\ B$ **and** $p \neq 0$
and $q\ deg: degree\ q = 1$
shows $card\ (proots\ within\ p\ B) = card\ (proots\ within\ (p\ \circ_p\ q)\ A)$ $\langle proof \rangle$

end

2 Budan–Fourier theorem

theory *Budan-Fourier* **imports**
BF-Misc
begin

The Budan–Fourier theorem is a classic result in real algebraic geometry to over-approximate real roots of a polynomial (counting multiplicity) within an interval. When all roots of the the polynomial are known to be real, the over-approximation becomes tight – the number of roots are counted exactly. Also note that Descartes’ rule of sign is a direct consequence of the Budan–Fourier theorem.

The proof mainly follows Theorem 2.35 in Basu, S., Pollack, R., Roy, M.-F.: Algorithms in Real Algebraic Geometry. Springer Berlin Heidelberg, Berlin, Heidelberg (2006).

2.1 More results related to *sign-r-pos*

lemma *sign-r-pos-nzero-right*:
assumes $nzero:\forall x. c < x \wedge x \leq d \longrightarrow poly\ p\ x \neq 0$ **and** $c < d$
shows *if sign-r-pos p c then poly p d > 0 else poly p d < 0*
 $\langle proof \rangle$

lemma *sign-r-pos-at-left*:
assumes $p \neq 0$
shows *if even (order c p) \longleftrightarrow sign-r-pos p c then eventually $(\lambda x. poly\ p\ x > 0)$ (at-left c)*
else eventually $(\lambda x. poly\ p\ x < 0)$ (at-left c)
 $\langle proof \rangle$

lemma *sign-r-pos-nzero-left*:
assumes $nzero:\forall x. d \leq x \wedge x < c \longrightarrow poly\ p\ x \neq 0$ **and** $d < c$
shows *if even (order c p) \longleftrightarrow sign-r-pos p c then poly p d > 0 else poly p d < 0*
 $\langle proof \rangle$

2.2 Fourier sequences

function *pders*::*real poly* \Rightarrow *real poly list* **where**
pders *p* = (if *p* = 0 then [] else Cons *p* (*pders* (*pderiv* *p*)))
 ⟨*proof*⟩

termination
 ⟨*proof*⟩

declare *pders.simps*[*simp del*]

lemma *set-pders-nzero*:
assumes *p* ≠ 0 *q* ∈ *set* (*pders* *p*)
shows *q* ≠ 0
 ⟨*proof*⟩

2.3 Sign variations for Fourier sequences

definition *changes-itv-der*:: *real* \Rightarrow *real* \Rightarrow *real poly* \Rightarrow *int* **where**
changes-itv-der *a* *b* *p* = (let *ps* = *pders* *p* in *changes-poly-at* *ps* *a* − *changes-poly-at* *ps* *b*)

definition *changes-gt-der*:: *real* \Rightarrow *real poly* \Rightarrow *int* **where**
changes-gt-der *a* *p* = *changes-poly-at* (*pders* *p*) *a*

definition *changes-le-der*:: *real* \Rightarrow *real poly* \Rightarrow *int* **where**
changes-le-der *b* *p* = (*degree* *p* − *changes-poly-at* (*pders* *p*) *b*)

lemma *changes-poly-pos-inf-pders*[*simp*]: *changes-poly-pos-inf* (*pders* *p*) = 0
 ⟨*proof*⟩

lemma *changes-poly-neg-inf-pders*[*simp*]: *changes-poly-neg-inf* (*pders* *p*) = *degree* *p*
 ⟨*proof*⟩

lemma *pders-coeffs-sgn-eq*: *map* ($\lambda p. \text{sgn}(\text{poly } p \ 0)$) (*pders* *p*) = *map* *sgn* (*coeffs* *p*)
 ⟨*proof*⟩

lemma *changes-poly-at-pders-0*: *changes-poly-at* (*pders* *p*) 0 = *changes* (*coeffs* *p*)
 ⟨*proof*⟩

2.4 Budan–Fourier theorem

lemma *budan-fourier-aux-right*:
assumes *c* < *d2* and *p* ≠ 0
assumes $\forall x. c < x \wedge x \leq d2 \longrightarrow (\forall q \in \text{set } (\text{pders } p). \text{poly } q \ x \neq 0)$
shows *changes-itv-der* *c* *d2* *p* = 0
 ⟨*proof*⟩

lemma *budan-fourier-aux-left'*:
assumes *d1* < *c* and *p* ≠ 0

assumes $\forall x. d1 \leq x \wedge x < c \longrightarrow (\forall q \in \text{set } (pders\ p). \text{poly } q\ x \neq 0)$
shows $\text{changes-itv-der } d1\ c\ p \geq \text{order } c\ p \wedge \text{even } (\text{changes-itv-der } d1\ c\ p - \text{order } c\ p)$
 $\langle \text{proof} \rangle$

lemma *budan-fourier-aux-left*:

assumes $d1 < c$ **and** $p \neq 0$
assumes $nzero: \forall x. d1 < x \wedge x < c \longrightarrow (\forall q \in \text{set } (pders\ p). \text{poly } q\ x \neq 0)$
shows $\text{changes-itv-der } d1\ c\ p \geq \text{order } c\ p$ **even** $(\text{changes-itv-der } d1\ c\ p - \text{order } c\ p)$
 $\langle \text{proof} \rangle$

theorem *budan-fourier-interval*:

assumes $a < b$ $p \neq 0$
shows $\text{changes-itv-der } a\ b\ p \geq \text{roots-count } p\ \{x. a < x \wedge x \leq b\} \wedge$
even $(\text{changes-itv-der } a\ b\ p - \text{roots-count } p\ \{x. a < x \wedge x \leq b\})$
 $\langle \text{proof} \rangle$

theorem *budan-fourier-gt*:

assumes $p \neq 0$
shows $\text{changes-gt-der } a\ p \geq \text{roots-count } p\ \{x. a < x\} \wedge$
even $(\text{changes-gt-der } a\ p - \text{roots-count } p\ \{x. a < x\})$
 $\langle \text{proof} \rangle$

Descartes' rule of signs is a direct consequence of the Budan–Fourier theorem

theorem *descartes-sign*:

fixes $p::\text{real poly}$
assumes $p \neq 0$
shows $\text{changes } (\text{coeffs } p) \geq \text{roots-count } p\ \{x. 0 < x\} \wedge$
even $(\text{changes } (\text{coeffs } p) - \text{roots-count } p\ \{x. 0 < x\})$
 $\langle \text{proof} \rangle$

theorem *budan-fourier-le*:

assumes $p \neq 0$
shows $\text{changes-le-der } b\ p \geq \text{roots-count } p\ \{x. x \leq b\} \wedge$
even $(\text{changes-le-der } b\ p - \text{roots-count } p\ \{x. x \leq b\})$
 $\langle \text{proof} \rangle$

2.5 Count exactly when all roots are real

definition *all-roots-real*:: $\text{real poly} \Rightarrow \text{bool}$ **where**

$\text{all-roots-real } p = (\forall r \in \text{roots } (\text{map-poly of-real } p). \text{Im } r = 0)$

lemma *all-roots-real-mult[simp]*:

$\text{all-roots-real } (p * q) \longleftrightarrow \text{all-roots-real } p \wedge \text{all-roots-real } q$
 $\langle \text{proof} \rangle$

lemma *all-roots-real-const-iff*:

assumes *all-real:all-roots-real* p
shows $\text{degree } p \neq 0 \iff (\exists x. \text{poly } p \ x = 0)$
 $\langle \text{proof} \rangle$

lemma *all-roots-real-degree:*
assumes *all-roots-real* p
shows $\text{roots-count } p \ \text{UNIV} = \text{degree } p \ \langle \text{proof} \rangle$

lemma *all-real-roots-mobius:*
fixes $a \ b :: \text{real}$
assumes *all-roots-real* p **and** $a < b$
shows *all-roots-real* $(fcompose \ p \ [:a,b:] \ [:1,1:]) \ \langle \text{proof} \rangle$

If all roots are real, we can use the Budan–Fourier theorem to EXACTLY count the number of real roots.

corollary *budan-fourier-real:*
assumes $p \neq 0$
assumes *all-roots-real* p
shows $\text{roots-count } p \ \{x. \ x \leq a\} = \text{changes-le-der } a \ p$
 $a < b \implies \text{roots-count } p \ \{x. \ a < x \wedge x \leq b\} = \text{changes-itv-der } a \ b \ p$
 $\text{roots-count } p \ \{x. \ b < x\} = \text{changes-gt-der } b \ p$
 $\langle \text{proof} \rangle$

Similarly, Descartes’ rule of sign counts exactly when all roots are real.

corollary *descartes-sign-real:*
fixes $p :: \text{real poly}$ **and** $a \ b :: \text{real}$
assumes $p \neq 0$
assumes *all-roots-real* p
shows $\text{roots-count } p \ \{x. \ 0 < x\} = \text{changes } (\text{coeffs } p)$
 $\langle \text{proof} \rangle$

end

3 Extension of Sturm’s theorem for multiple roots

theory *Sturm-Multiple-Roots*
imports
BF-Misc
begin

The classic Sturm’s theorem is used to count real roots WITHOUT multiplicity of a polynomial within an interval. Surprisingly, we can also extend Sturm’s theorem to count real roots WITH multiplicity by modifying the signed remainder sequence, which seems to be overlooked by many textbooks.

Our formal proof is inspired by Theorem 10.5.6 in Rahman, Q.I., Schmeisser, G.: Analytic Theory of Polynomials. Oxford University Press (2002).

3.1 More results for *smods*

lemma *last-smods-gcd*:
fixes $p\ q :: \text{real poly}$
defines $pp \equiv \text{last } (\text{smods } p\ q)$
assumes $p \neq 0$
shows $pp = \text{smult } (\text{lead-coeff } pp) (\text{gcd } p\ q)$
 $\langle \text{proof} \rangle$

lemma *last-smods-nzero*:
assumes $p \neq 0$
shows $\text{last } (\text{smods } p\ q) \neq 0$
 $\langle \text{proof} \rangle$

3.2 Alternative signed remainder sequences

function *smods-ext*:: $\text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{real poly list}$ **where**
 $\text{smods-ext } p\ q = (\text{if } p=0 \text{ then } [] \text{ else}$
 $\quad (\text{if } p \bmod q \neq 0$
 $\quad \quad \text{then } \text{Cons } p (\text{smods-ext } q (- (p \bmod q)))$
 $\quad \quad \text{else } \text{Cons } p (\text{smods-ext } q (pderiv\ q)))$
 $\quad)$
 $\langle \text{proof} \rangle$
termination
 $\langle \text{proof} \rangle$

lemma *smods-ext-prefix*:
fixes $p\ q :: \text{real poly}$
defines $pp \equiv \text{last } (\text{smods } p\ q)$
assumes $p \neq 0\ q \neq 0$
shows $\text{smods-ext } p\ q = \text{smods } p\ q @ \text{tl } (\text{smods-ext } pp (pderiv\ pp))$
 $\langle \text{proof} \rangle$

lemma *no-0-in-smods-ext*: $0 \notin \text{set } (\text{smods-ext } p\ q)$
 $\langle \text{proof} \rangle$

3.3 Sign variations on the alternative signed remainder sequences

definition *changes-itv-smods-ext*:: $\text{real} \Rightarrow \text{real} \Rightarrow \text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{int}$
where
 $\text{changes-itv-smods-ext } a\ b\ p\ q = (\text{let } ps = \text{smods-ext } p\ q \text{ in } \text{changes-poly-at } ps\ a$
 $\quad - \text{changes-poly-at } ps\ b)$

definition *changes-gt-smods-ext*:: $\text{real} \Rightarrow \text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{int}$ **where**
 $\text{changes-gt-smods-ext } a\ p\ q = (\text{let } ps = \text{smods-ext } p\ q \text{ in } \text{changes-poly-at } ps\ a$
 $\quad - \text{changes-poly-pos-inf } ps)$

definition *changes-le-smods-ext*:: $\text{real} \Rightarrow \text{real poly} \Rightarrow \text{real poly} \Rightarrow \text{int}$ **where**
 $\text{changes-le-smods-ext } b\ p\ q = (\text{let } ps = \text{smods-ext } p\ q \text{ in } \text{changes-poly-neg-inf } ps)$

– *changes-poly-at ps b*)

definition *changes-R-smods-ext*:: *real poly* \Rightarrow *real poly* \Rightarrow *int* **where**
changes-R-smods-ext p q = (*let ps* = *smods-ext p q* in *changes-poly-neg-inf ps*
– *changes-poly-pos-inf ps*)

3.4 Extension of Sturm’s theorem for multiple roots

theorem *sturm-ext-interval*:

assumes *a < b poly p a \neq 0 poly p b \neq 0*

shows *roots-count p {x. a < x \wedge x < b}* = *changes-itv-smods-ext a b p (pderiv p)*
(*proof*)

theorem *sturm-ext-above*:

assumes *poly p a \neq 0*

shows *roots-count p {x. a < x}* = *changes-gt-smods-ext a p (pderiv p)*
(*proof*)

theorem *sturm-ext-below*:

assumes *poly p b \neq 0*

shows *roots-count p {x. x < b}* = *changes-le-smods-ext b p (pderiv p)*
(*proof*)

theorem *sturm-ext-R*:

assumes *p \neq 0*

shows *roots-count p UNIV* = *changes-R-smods-ext p (pderiv p)*
(*proof*)

end

4 Descartes Roots Test

theory *Descartes-Roots-Test* **imports** *Budan-Fourier*
begin

The Descartes roots test is a consequence of Descartes’ rule of signs: through counting sign variations on coefficients of a base-transformed (i.e. Taylor shifted) polynomial, it can over-approximate the number of real roots (counting multiplicity) within an interval. Its ability is similar to the Budan–Fourier theorem, but is far more efficient in practice. Therefore, this test is widely used in modern root isolation procedures.

More information can be found in the wiki page about Vincent’s theorem: https://en.wikipedia.org/wiki/Vincent%27s_theorem and Collins and Akritas’s classic paper of root isolation: Collins, G.E., Akritas, A.G.: Polynomial real root isolation using Descarte’s rule of signs. SYMSACC. 272–275 (1976). A more modern treatment is available from a recent implementation of isolating real roots: Kobel, A., Rouillier, F., Sagraloff, M.: Computing Real Roots of Real Polynomials ... and now For Real! Proceedings of ISSAC

'16, New York, New York, USA (2016).

lemma *bij-betw-pos-interval*:

fixes $a\ b::\text{real}$

assumes $a < b$

shows $\text{bij-betw } (\lambda x. (a+b * x) / (1+x)) \{x. x > 0\} \{x. a < x \wedge x < b\}$

<proof>

lemma *proots-sphere-pos-interval*:

fixes $a\ b::\text{real}$

defines $q1 \equiv [a, b]$ **and** $q2 \equiv [1, 1]$

assumes $p \neq 0$ $a < b$

shows $\text{proots-count } p \{x. a < x \wedge x < b\} = \text{proots-count } (\text{fcompose } p\ q1\ q2) \{x. 0 < x\}$

<proof>

definition *descartes-roots-test::real \Rightarrow real \Rightarrow real poly \Rightarrow nat where*

descartes-roots-test a b p = nat (changes (coeffs (fcompose p [a,b] [1,1:])))

theorem *descartes-roots-test*:

fixes $p::\text{real poly}$

assumes $p \neq 0$ $a < b$

shows $\text{proots-count } p \{x. a < x \wedge x < b\} \leq \text{descartes-roots-test } a\ b\ p \wedge$
even (descartes-roots-test a b p - proots-count p {x. a < x \wedge x < b})

<proof>

The roots test *descartes-roots-test* is exact if its result is 0 or 1.

corollary *descartes-roots-test-zero*:

fixes $p::\text{real poly}$

assumes $p \neq 0$ $a < b$ *descartes-roots-test a b p = 0*

shows $\forall x. a < x \wedge x < b \longrightarrow \text{poly } p\ x \neq 0$

<proof>

corollary *descartes-roots-test-one*:

fixes $p::\text{real poly}$

assumes $p \neq 0$ $a < b$ *descartes-roots-test a b p = 1*

shows $\text{proots-count } p \{x. a < x \wedge x < b\} = 1$

<proof>

Similar to the Budan–Fourier theorem, the Descartes roots test result is exact when all roots are real.

corollary *descartes-roots-test-real*:

fixes $p::\text{real poly}$

assumes $p \neq 0$ $a < b$

assumes *all-roots-real p*

shows $\text{proots-count } p \{x. a < x \wedge x < b\} = \text{descartes-roots-test } a\ b\ p$

<proof>

end

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