

A. Proofs

Theorem 5 Give any two finite sets $S_1, S_2 \in \Omega$, with $A = |S_1 \cup S_2| > a = |S_1 \cap S_2| > 0$ and $|\Omega| = D \rightarrow \infty$. The limiting variance of the estimators from densification and improved densification when $k = D \rightarrow \infty$ is given by:

$$\lim_{k \rightarrow \infty} \text{Var}(h) = \frac{a}{A} \left[\frac{A-a}{A(A+1)} \right] > 0 \quad (15)$$

$$\lim_{k \rightarrow \infty} \text{Var}(h^+) = \frac{a}{A} \left[\frac{3(A-1) + (2A-1)(a-1)}{2(A+1)(A-1)} - \frac{a}{A} \right] > 0 \quad (16)$$

Proof: When $k = D$, then $N_{emp} = D - A$. Substituting this value in the variance formulas from (Shrivastava & Li, 2014c) and taking the limit as $D = k \rightarrow \infty$, we get the above expression after manipulation. When $0 < R = \frac{a}{A} < 1$, they both are strictly positive. \square

Theorem 6

$$\text{Pr}(h^*(S_1) = h^*(S_2)) = \frac{|S_1 \cap S_2|}{|S_1 \cup S_2|} = R \quad (17)$$

$$\text{Var}(h^*) = \frac{R}{k} + A \frac{R}{k^2} + B \frac{R\tilde{R}}{k^2} - R^2 \quad (18)$$

$$\lim_{k \rightarrow \infty} \text{Var}(h^*) = 0 \quad (19)$$

where N_{emp} is the number of simultaneous empty bins between S_1 and S_2 and the quantities A and B are given by

$$A = \mathbb{E} \left[2N_{emp} + \frac{N_{emp}(N_{emp} - 1)}{k - N_{emp}} \right]$$

$$B = \mathbb{E} \left[(k - N_{emp})(k - N_{emp} - 1) + 2N_{emp}(k - N_{emp} - 1) + \frac{N_{emp}(N_{emp} - 1)(k - N_{emp} - 1)}{k - N_{emp}} \right]$$

Proof:

The collision probability is easy using a simple observation that values coming from different bin numbers can never match across S_1 and S_2 , i.e. $h_i^*(S_1) \neq h_j^*(S_2)$ if $i \neq j$, as they have disjoint different range. So whenever, for a simultaneous empty bin i , i.e. $E_i = 1$, we get $h_i^*(S_1) = h_i^*(S_2)$ after reassignment, the value must be coming from same non-empty bin, say numbers k which is not not empty. Thus,

$$\text{Pr}(h_i^*(S_1) = h_i^*(S_2)) = \text{Pr}(h_k^*(S_1) = h_k^*(S_2) | E_k = 0) = R \quad \text{By expansion and linearity of expectation, we obtain}$$

The variance is little involved. From the collision probability, we have the following is unbiased estimator.

$$\hat{R} = \frac{1}{k} \sum_{j=0}^{k-1} \mathbb{1}\{h_j^*(S_1) = h_j^*(S_2)\}. \quad (20)$$

For variance, define the number of simultaneously empty bins by

$$N_{emp} = \sum_{j=0}^{k-1} \mathbb{1}\{E_j = 1\}, \quad (21)$$

where $\mathbb{1}$ is the indicator function. We partition the event $(h_j^*(S_1) = h_j^*(S_2))$ into two cases depending on E_j . Let M_j^N (**Non-empty Match at j**) and M_j^E (**Empty Match at j**) be the events defined as:

$$M_j^N = \mathbb{1}\{E_j = 0 \text{ and } h_j^*(S_1) = h_j^*(S_2)\} \quad (22)$$

$$M_j^E = \mathbb{1}\{E_j = 1 \text{ and } h_j^*(S_1) = h_j^*(S_2)\} \quad (23)$$

Note that, $M_j^N = 1 \implies M_j^E = 0$ and $M_j^E = 1 \implies M_j^N = 0$. From the LSH property of estimator we have

$$\begin{aligned} \mathbb{E}(M_j^N | E_j = 0) &= \mathbb{E}(M_j^E | E_j = 1) \\ &= \mathbb{E}(M_j^E + M_j^N) = R \quad \forall j \end{aligned} \quad (24)$$

It is not difficult to show that,

$$\mathbb{E}(M_j^N M_i^N | i \neq j, E_j = 0 \text{ and } E_i = 0) = R\tilde{R},$$

where $\tilde{R} = \frac{a-1}{f_1 + f_2 - a - 1}$. Using these new events, we have

$$\hat{R} = \frac{1}{k} \sum_{j=0}^{k-1} [M_j^E + M_j^N] \quad (25)$$

We are interested in computing

$$\text{Var}(\hat{R}) = \mathbb{E} \left(\left(\frac{1}{k} \sum_{j=0}^{k-1} [M_j^E + M_j^N] \right)^2 \right) - R^2 \quad (26)$$

For notational convenience we will use m to denote the event $k - N_{emp} = m$, i.e., the expression $\mathbb{E}(\cdot | m)$ means $\mathbb{E}(\cdot | k - N_{emp} = m)$. To simplify the analysis, we will first compute the conditional expectation

$$f(m) = \mathbb{E} \left(\left(\frac{1}{k} \sum_{j=0}^{k-1} [M_j^E + M_j^N] \right)^2 \middle| m \right) \quad (27)$$

By expansion and linearity of expectation, we obtain

$$\begin{aligned} k^2 f(m) &= \mathbb{E} \left[\sum_{i \neq j} M_i^N M_j^N \middle| m \right] + \mathbb{E} \left[\sum_{i \neq j} M_i^E M_j^E \middle| m \right] \\ &+ \mathbb{E} \left[\sum_{i \neq j} M_i^E M_j^N \middle| m \right] + \mathbb{E} \left[\sum_{i=1}^k [(M_j^N)^2 + (M_j^E)^2] \middle| m \right] \end{aligned}$$

$M_j^N = (M_j^N)^2$ and $M_j^E = (M_j^E)^2$ as they are indicator functions and can only take values 0 and 1. Hence,

$$\mathbb{E} \left[\sum_{j=0}^{k-1} [(M_j^N)^2 + (M_j^E)^2] \middle| m \right] = kR \quad (28)$$

The values of the first three terms are given by the following 3 expression using simple binomial expansion and using the fact that we are dealing with indicator random variable which can only take values 0 or 1.

$$\mathbb{E} \left[\sum_{i \neq j} M_i^N M_j^N \middle| m \right] = m(m-1)R\tilde{R} \quad (29)$$

$$\mathbb{E} \left[\sum_{i \neq j} M_i^N M_j^E \middle| m \right] = 2m(k-m) \left[\frac{R}{m} + \frac{(m-1)R\tilde{R}}{m} \right] \quad (30)$$

Let p be the probability that two simultaneously empty bins i and j finally picks the same non-empty bin for reassignment. Then we have

$$\mathbb{E} \left[\sum_{i \neq j} M_i^E M_j^E \middle| m \right] = (k-m)(k-m-1) [pR + (1-p)R\tilde{R}] \quad (31)$$

because with probability $(1-p)$, it uses estimators from different simultaneous non-empty bin and in that case the $M_i^E M_j^E = 1$ with probability $R\tilde{R}$. We know that Algorithm 1 which uses 2-universal hashing the value of $p = \frac{1}{m}$. This is because any pairwise assignment is perfectly random with 2-universal hashing.

Substituting for all terms with value of p and rearranging terms gives the required expression.

When $k = D$, then $N_{emp} = D - A$. Substituting this value in the variance formulas and taking the limit as $D = k \rightarrow \infty$, we get 0 for all R .

Theorem 7

$$\text{Var}(h^*) \leq \text{Var}(h^+) \leq \text{Var}(h) \quad (32)$$

Proof: We have $p^* = \frac{1}{m} \leq p^+ = \frac{1.5}{m+1} \leq p = \frac{2}{m+1}$. The value of p^+ and p comes from analysis in (Shrivastava & Li, 2014c)

Theorem 8 Among all densification schemes, where the reassignment process for bin i is independent of the reassignment process of any other bin j , Algorithm 1 achieves the best possible variance.

Under any independent re-assignment, the probability that two empty bins chooses the same non-empty bin out of m non-empty bins is lower bounded by $\frac{1}{m}$ which is achieved by optimal densification.