

## Appendix A. Preliminary results

We provide here two geometric results (Corollary 26 and Lemma 27) and a stochastic result (Proposition 28) that are used repeatedly in the computations. We start with the definition of covering numbers.

**Definition 24** (COVERING NUMBER AND  $\epsilon$ -COVER) *For any compact and convex set  $\mathcal{X} \subset \mathbb{R}^d$  and any  $\epsilon > 0$ , we say that a sequence  $x_1, \dots, x_n$  of  $n$  points in  $\mathcal{X}$  defines an  $\epsilon$ -cover of  $\mathcal{X}$  if and only if  $\mathcal{X} \subseteq \bigcup_{i=1}^n B(x_i, \epsilon)$ . The covering number  $\mathcal{N}_\epsilon(\mathcal{X})$  of  $\mathcal{X}$  is then defined as the minimal size of a sequence defining an  $\epsilon$ -cover of  $\mathcal{X}$ , i.e.*

$$\mathcal{N}_\epsilon(\mathcal{X}) := \inf \left\{ n \in \mathbb{N}^* : \exists (x_1, \dots, x_n) \in \mathcal{X}^n \text{ s.t. } \mathcal{X} \subseteq \bigcup_{i=1}^n B(x_i, \epsilon) \right\}.$$

The next result provides an upper bound on the covering numbers of hypercubes.

**Proposition 25** (COVERING NUMBER OF HYPERCUBES) *Let  $[0, R]^d$  be an hypercube of dimensionality  $d \geq 1$  whose side has length  $R > 0$ . Then, for all  $\epsilon > 0$ , we have that*

$$\mathcal{N}_\epsilon([0, R]^d) \leq (\sqrt{d}R/2\epsilon)^d \vee 1.$$

**Proof** Observe first that since  $[0, R]^d \subseteq B(c, \sqrt{d}R/2)$  where  $c$  denotes the center of the hypercube, then the result trivially holds for any  $\epsilon \geq \sqrt{d}R/2$ . Fix any  $\epsilon < \sqrt{d}R/2$ , set  $N_\epsilon = \lceil \sqrt{d}R/2\epsilon \rceil$  and define for all  $I \in \{0, \dots, N_\epsilon - 1\}^d$  the series  $H_I = I \times R/N_\epsilon + [0, R/N_\epsilon]^d$  of  $N_\epsilon^d$  hypercubes which cover  $[0, R]^d := \bigcup_{I \in \{0, \dots, N_\epsilon - 1\}^d} H_I$ . Denoting by  $c_I$  the center of  $H_I$  and since  $\max_{x \in H_I} \|x - c_I\|_2 \leq \epsilon$ , it necessarily follows that  $H_I \subseteq B(c_I, \epsilon)$  which implies that  $[0, R]^d \subseteq \bigcup_{I \in \{0, \dots, N_\epsilon - 1\}^d} B(c_I, \epsilon)$  and proves that  $\mathcal{N}_\epsilon([0, R]^d) \leq N_\epsilon^d \leq (\sqrt{d}R/2\epsilon)^d$ .  $\square$

This result can be extended to any compact and convex set of  $\mathbb{R}^d$  as shown below.

**Corollary 26** (COVERING NUMBER OF A CONVEX SET) *For any bounded compact and convex set  $\mathcal{X} \subset \mathbb{R}^d$ , we have that  $\forall \epsilon > 0$ ,*

$$\mathcal{N}_\epsilon(\mathcal{X}) \leq (\sqrt{d} \text{diam}(\mathcal{X}) / \epsilon)^d \vee 1.$$

**Proof** First, we show that  $\mathcal{N}_\epsilon(\mathcal{X}) \leq \mathcal{N}_\epsilon([0, 2 \text{diam}(\mathcal{X})]^d)$  and then, we use the bound of Proposition 25 to conclude the proof. By definition of  $\text{diam}(\mathcal{X})$ , we know that there exists some  $x \in \mathbb{R}^d$  such that  $\mathcal{X} \subseteq x + [0, 2 \text{diam}(\mathcal{X})]^d$ . Hence, we know from Proposition 25 that there exists a sequence  $c_1, \dots, c_{N_\epsilon}$  of  $N_\epsilon \leq \mathcal{N}_\epsilon([0, 2 \text{diam}(\mathcal{X})]^d)$  points in  $[0, 2 \text{diam}(\mathcal{X})]^d$  forming an  $\epsilon$ -cover of  $\mathcal{X}$ :

$$\mathcal{X} \subseteq [0, 2 \text{diam}(\mathcal{X})]^d \subseteq \bigcup_{i=1}^{N_\epsilon} B(c_i, \epsilon). \quad (1)$$

However, we do not have the guarantee at this point that the centers  $c_1, \dots, c_{N_\epsilon}$  belong to  $\mathcal{X}$ . To build an  $\epsilon$ -cover of  $\mathcal{X}$ , we project each of those centers on  $\mathcal{X}$ . More precisely, we show that  $\mathcal{X} \subseteq \bigcup_{i=1}^{N_\epsilon} B(\Pi_{\mathcal{X}}(c_i), \epsilon)$  where  $\Pi_{\mathcal{X}} : x \in \mathbb{R}^d \mapsto \arg \min_{x' \in \mathcal{X}} \|x - x'\|_2 \in \mathcal{X}$  denotes

the projection over the compact and convex set  $\mathcal{X}$ . Starting from (1), it is sufficient to show that  $B(c_i, \epsilon) \cap \mathcal{X} \subseteq B(\Pi_{\mathcal{X}}(c_i), \epsilon)$ , for all  $i \in \{1, \dots, N_\epsilon\}$  to prove that

$$\mathcal{X} \subseteq \bigcup_{i=1}^{N_\epsilon} B(c_i, \epsilon) \cap \mathcal{X} \subseteq \bigcup_{i=1}^{N_\epsilon} B(\Pi_{\mathcal{X}}(c_i), \epsilon).$$

Pick any  $c \in \{c_1, \dots, c_{N_\epsilon}\}$  and consider the following cases on the distance  $\|c - \Pi_{\mathcal{X}}(c)\|_2$  between the center and its projection: (i) if  $\|c - \Pi_{\mathcal{X}}(c)\|_2 = 0$ , then  $c = \Pi_{\mathcal{X}}(c)$  and we have  $B(c, \epsilon) \cap \mathcal{X} \subseteq B(\Pi_{\mathcal{X}}(c), \epsilon)$ , (ii) If  $\|c - \Pi_{\mathcal{X}}(c)\|_2 > \epsilon$ , then  $\mathcal{X} \cap B(c, \epsilon) = \emptyset$ , and we have  $\mathcal{X} \cap B(c, \epsilon) \subseteq B(\Pi_{\mathcal{X}}(c), \epsilon)$ . We now consider the non-trivial case where  $\|c - \Pi_{\mathcal{X}}(c)\|_2 \in (0, \epsilon)$ . Pick any  $x \in B(c, \epsilon) \cap \mathcal{X}$  and note that since  $x \in B(c, \epsilon)$ , then

$$\begin{aligned} \epsilon^2 &\geq \|x - c\|_2^2 \\ &= \|x - \Pi_{\mathcal{X}}(c) + \Pi_{\mathcal{X}}(c) - c\|_2^2 \\ &= \|x - \Pi_{\mathcal{X}}(c)\|_2^2 + \|c - \Pi_{\mathcal{X}}(c)\|_2^2 + 2 \cdot \langle x - \Pi_{\mathcal{X}}(c), \Pi_{\mathcal{X}}(c) - c \rangle \end{aligned}$$

which combined with the fact that  $\|c - \Pi_{\mathcal{X}}(c)\|_2^2 \geq 0$  gives

$$\|x - \Pi_{\mathcal{X}}(c)\|_2^2 \leq \epsilon^2 - 2 \cdot \langle x - \Pi_{\mathcal{X}}(c), \Pi_{\mathcal{X}}(c) - c \rangle. \quad (2)$$

We will simply show that the inner product  $\langle x - \Pi_{\mathcal{X}}(c), \Pi_{\mathcal{X}}(c) - c \rangle$  cannot be strictly negative to prove that  $\|x - \Pi_{\mathcal{X}}(c)\|_2 \leq \epsilon$ . Assume by contradiction that  $\langle x - \Pi_{\mathcal{X}}(c), \Pi_{\mathcal{X}}(c) - c \rangle < 0$ . Since  $\Pi_{\mathcal{X}}(c) \in \mathcal{X}$  and  $x \in \mathcal{X}$ , it follows the convexity of  $\mathcal{X}$  implies that  $\forall \lambda \in [0, 1]$ ,  $x_\lambda = \Pi_{\mathcal{X}}(c) + \lambda \cdot (x - \Pi_{\mathcal{X}}(c)) \in \mathcal{X}$ . However, for all  $\lambda \in (0, 1)$  we have that

$$\begin{aligned} \|x_\lambda - c\|_2^2 &= \|\Pi_{\mathcal{X}}(c) - c + \lambda \cdot (x - \Pi_{\mathcal{X}}(c))\|_2^2 \\ &= \|\Pi_{\mathcal{X}}(c) - c\|_2^2 + \lambda^2 \|x - \Pi_{\mathcal{X}}(c)\|_2^2 + 2\lambda \cdot \langle \Pi_{\mathcal{X}}(c) - c, x - \Pi_{\mathcal{X}}(c) \rangle \\ &= \|\Pi_{\mathcal{X}}(c) - c\|_2^2 + \lambda \cdot (\lambda \|x - \Pi_{\mathcal{X}}(c)\|_2^2 + 2 \cdot \langle \Pi_{\mathcal{X}}(c) - c, x - \Pi_{\mathcal{X}}(c) \rangle). \end{aligned}$$

Therefore, taking any  $0 < \lambda^* < |\langle \Pi_{\mathcal{X}}(c) - c, x - \Pi_{\mathcal{X}}(c) \rangle| / \|\Pi_{\mathcal{X}}(c) - c\|_2^2 \wedge 1$  so that the second term of the right hand term of the previous equation is strictly negative gives that  $\|x_{\lambda^*} - c\|_2^2 < \|\Pi_{\mathcal{X}}(c) - c\|_2^2$  leads us to the following contradiction  $\min_{x \in \mathcal{X}} \|x - c\|_2 \leq \|x_{\lambda^*} - c\|_2 < \|\Pi_{\mathcal{X}}(c) - c\|_2 = \min_{x \in \mathcal{X}} \|x - c\|_2$ . Hence,  $\langle x - \Pi_{\mathcal{X}}(c), \Pi_{\mathcal{X}}(c) - c \rangle \geq 0$  and we deduce from (2) that  $\mathcal{X} \cap B(c, \epsilon) \subseteq B(\Pi_{\mathcal{X}}(c), \epsilon)$ , which completes the proof.  $\square$

The next inequality will be useful to bound to bound volume of the intersection of a ball and a convex set.

**Lemma 27** (From Zabinsky and Smith (1992), see Appendix Section therein). *For any compact and convex set  $\mathcal{X} \subset \mathbb{R}^d$  with non-empty interior, we have that for any  $x^* \in \mathcal{X}$  and  $\epsilon \in (0, \text{diam}(\mathcal{X}))$ ,*

$$\frac{\mu(B(x^*, \epsilon) \cap \mathcal{X})}{\mu(\mathcal{X})} \geq \left( \frac{\epsilon}{\text{diam}(\mathcal{X})} \right)^d.$$

**Proof** We point out that a detailed proof of this result can be found in the Appendix Section of (Zabinsky and Smith (1992)). Nonetheless, we provide here a proof with less details for completeness. Introduce the similarity transformation  $S : \mathbb{R}^d \rightarrow \mathbb{R}^d$  defined by

$$S : x \mapsto x^* + \frac{r}{\text{diam}(\mathcal{X})}(x - x^*)$$

and let  $S(\mathcal{X}) := \{S(x) : x \in \mathcal{X}\}$  be the image of  $\mathcal{X}$  by  $S$ . Since  $x^* \in \mathcal{X}$  and  $\max_{x \in \mathcal{X}} \|x - x^*\|_2 \leq \text{diam}(\mathcal{X})$  by definition, it follows from the convexity of  $\mathcal{X}$  that  $S(\mathcal{X}) \subseteq B(x^*, r) \cap \mathcal{X}$  which implies that  $\mu(B(x^*, r) \cap \mathcal{X}) \geq \mu(S(\mathcal{X}))$ . However, as  $S$  is a similarity transformation conserves the ratios of the volumes before/after transformation, we thus deduce that

$$\frac{\mu(B(x^*, r) \cap \mathcal{X})}{\mu(\mathcal{X})} \geq \frac{\mu(S(\mathcal{X}))}{\mu(\mathcal{X})} = \frac{\mu(S(B(x^*, \text{diam}(\mathcal{X}))))}{\mu(B(x^*, \text{diam}(\mathcal{X})))} = \frac{\mu(B(x^*, r))}{\mu(B(x^*, \text{diam}(\mathcal{X})))}$$

and the result follows using the fact that  $\forall r \geq 0$ ,  $\mu(B(x^*, r)) = \pi^{d/2} r^d / \Gamma(d/2 + 1)$  where  $\Gamma(\cdot)$  stands for the standard gamma function.  $\square$

**Proposition 28** (PURE RANDOM SEARCH) *Let  $\mathcal{X} \subset \mathbb{R}^d$  be a compact and convex set with non-empty interior and let  $f \in \text{Lip}(k)$  be a  $k$ -Lipschitz functions defined on  $\mathcal{X}$  for some  $k \geq 0$ . Then, for any  $n \in \mathbb{N}^*$  and  $\delta \in (0, 1)$ , we have with probability at least  $1 - \delta$ ,*

$$\max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) \leq k \cdot \text{diam}(\mathcal{X}) \cdot \left( \frac{\ln(1/\delta)}{n} \right)^{\frac{1}{d}}$$

where  $X_1, \dots, X_n$  denotes a sequence of  $n$  independent copies of  $X \sim \mathcal{U}(\mathcal{X})$ .

**Proof** Fix any  $n \in \mathbb{N}^*$  and  $\delta \in (0, 1)$ , let  $\epsilon = k \text{diam}(\mathcal{X}) (\ln(1/\delta)/n)^{1/d}$  be the value of the upper bound and  $\mathcal{X}_\epsilon = \{x \in \mathcal{X} : f(x) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon\}$  the corresponding level set. As the result trivially holds whenever  $n \leq \ln(1/\delta)$ , we consider that  $n > \ln(1/\delta)$ . Observe now that since  $f \in \text{Lip}(k)$ , then for any  $x^* \in \arg \max_{x \in \mathcal{X}} f(x)$ , we have that  $\mathcal{X} \cap B(x^*, \epsilon/k) \subseteq \mathcal{X}_\epsilon$  since  $|f(x) - f(x^*)| \leq k \cdot \|x - x^*\|_2 = \epsilon$  for all  $x \in B(x^*, \epsilon/k) \cap \mathcal{X}$ . Therefore, by picking any  $x^* \in \arg \max_{x \in \mathcal{X}} f(x)$ , one gets

$$\begin{aligned} \mathbb{P} \left( \max_{i=1 \dots n} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon \right) &= \mathbb{P} \left( \bigcup_{i=1}^n \{X_i \in \mathcal{X}_\epsilon\} \right) && \text{(def. of } \mathcal{X}_\epsilon) \\ &= 1 - \mathbb{P}(X_1 \notin \mathcal{X}_\epsilon)^n && \text{(i.i.d. r.v.)} \\ &\geq 1 - \mathbb{P}(X_1 \notin \mathcal{X} \cap B(x^*, \epsilon/k))^n && (\mathcal{X} \cap B(x^*, \epsilon/k) \subseteq \mathcal{X}_\epsilon) \\ &= 1 - \left( 1 - \left( \frac{\mu(\mathcal{X} \cap B(x^*, \epsilon/k))}{\mu(\mathcal{X})} \right)^d \right)^n && (X_1 \sim \mathcal{U}(\mathcal{X})) \\ &\geq 1 - \left( 1 - \left( \frac{\epsilon}{k \text{diam}(\mathcal{X})} \right)^d \right)^n && \text{(Lemma 27)} \\ &= 1 - \left( 1 - \frac{\ln(1/\delta)}{n} \right)^n && \text{(def. of } \epsilon) \\ &\geq 1 - \delta. && (1 + x \leq e^x) \end{aligned}$$

$\square$

## Appendix B. Proofs of Section 3

In this section, we provide the proofs of Propositions 3, 5, 6 and Example 4.

**Proof of proposition 3.** ( $\Leftarrow$ ) Let  $A$  be any global optimization algorithm such that  $\forall f \in \bigcup_{k \geq 0} \text{Lip}(k)$ ,  $\sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|X_i - x\|_2 \xrightarrow{p} 0$ . Pick any  $\epsilon > 0$ , any  $f \in \bigcup_{k \geq 0} \text{Lip}(k)$  and let  $\mathcal{X}_\epsilon = \{x \in \mathcal{X} : f(x) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon\}$  be the corresponding level set. As  $\mathcal{X}_\epsilon$  is non-empty, there necessarily exists some  $x_\epsilon \in \mathcal{X}$  and  $r_\epsilon > 0$  such that  $B(x_\epsilon, r_\epsilon) \cap \mathcal{X} \subseteq \mathcal{X}_\epsilon$ . Thus, if  $X_1, \dots, X_n$  denotes a sequence a sequence of  $n$  evaluation points generated by  $A$  over  $f$ , we directly obtain from the convergence in probability of the mesh grid that

$$\begin{aligned} \mathbb{P} \left( \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) > \epsilon \right) &= \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin \mathcal{X}_\epsilon\} \right) \\ &\leq \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin B(x_\epsilon, r_\epsilon)\} \right) \\ &= \mathbb{P} \left( \min_{i=1 \dots n} \|X_i - x_\epsilon\|_2 > r_\epsilon \right) \\ &\leq \mathbb{P} \left( \sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|X_i - x\|_2 > r_\epsilon \right) \xrightarrow{n \rightarrow \infty} 0. \end{aligned}$$

( $\Rightarrow$ ) Let  $A$  be any global optimization algorithm consistent over the set of Lipschitz functions and assume by contradiction that there exists some  $f^* \in \bigcup_{k \geq 0} \text{Lip}(k)$  such that  $\sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|x - X_i\|_2 \not\xrightarrow{p} 0$ . The implication is proved in two steps: first, we show that there exists a ball  $B(c^*, \epsilon)$  for some  $c^* \in \mathcal{X}$  which is almost never hit by the algorithm and second, we build a Lipschitz function which admits its maximum over this ball.

*First step.* Let  $\{X_i\}_{i \in \mathbb{N}^*}$  be a sequence of evaluation points generated by  $A$  over  $f^*$ . Observe first that since for all  $\epsilon > 0$ , the series  $n \in \mathbb{N}^* \mapsto \mathbb{P}(\sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|x - X_i\|_2 > \epsilon)$  is non-increasing, then the contradiction assumption necessarily implies that

$$\exists \epsilon_1, \epsilon_2 > 0 \text{ such that } \forall n \in \mathbb{N}^*, \mathbb{P} \left( \sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|x - X_i\|_2 > \epsilon_1 \right) > \epsilon_2. \quad (3)$$

Consider now any sequence  $c_1, \dots, c_{N_1}$  of  $N_1 = \mathcal{N}_{\epsilon_1}(\mathcal{X})$  points in  $\mathcal{X}$  defining an  $\epsilon_1$ -cover of  $\mathcal{X}$  and suppose by contradiction that

$$\forall c \in \{c_1, \dots, c_{N_1}\}, \exists n_c \in \mathbb{N}^* \text{ such that } \mathbb{P} \left( \bigcap_{i=1}^{n_c} \{X_i \notin B(c, \epsilon_1) \cap \mathcal{X}\} \right) \leq \frac{\epsilon_2}{2N_1}$$

which gives by setting  $N_2 = \max_{c \in \{c_1, \dots, c_{N_1}\}} n_c$  that

$$\forall c \in \{c_1, \dots, c_{N_1}\}, \mathbb{P} \left( \bigcap_{i=1}^{N_2} \{X_i \notin B(c, \epsilon) \cap \mathcal{X}\} \right) \leq \frac{\epsilon_2}{2N_1}.$$

However, as  $c_1, \dots, c_{N_1}$  form an  $\epsilon_1$ -cover of  $\mathcal{X}$ , it follows that

$$\begin{aligned}
 \mathbb{P} \left( \sup_{x \in \mathcal{X}} \min_{i=1 \dots N_2} \|x - X_i\|_2 \leq \epsilon_1 \right) &\geq \mathbb{P} \left( \bigcap_{j=1}^{N_1} \bigcup_{i=1}^{N_2} \{X_i \in B(c_j, \epsilon_1) \cap \mathcal{X}\} \right) \\
 &= 1 - \mathbb{P} \left( \bigcup_{j=1}^{N_1} \bigcap_{i=1}^{N_2} \{X_i \notin B(c_j, \epsilon_1) \cap \mathcal{X}\} \right) \\
 &\geq 1 - \sum_{j=1}^{N_1} \mathbb{P} \left( \bigcap_{i=1}^{N_2} \{X_i \notin B(c_j, \epsilon) \cap \mathcal{X}\} \right) \\
 &\geq 1 - N_1 \times \frac{\epsilon_2}{2N_1} \\
 &= 1 - \frac{\epsilon_2}{2}
 \end{aligned}$$

which contradicts (3). Hence, we deduce that

$$\exists c^* \in \{c_1, \dots, c_{N_\epsilon}\} \text{ such that } \forall n \in \mathbb{N}^*, \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin B(c^*, \epsilon_1) \cap \mathcal{X}\} \right) \geq \frac{\epsilon_2}{2N_1}.$$

*Second Step.* Based on this center  $c^* \in \mathcal{X}$ , one can introduce the function  $\tilde{f} : \mathcal{X} \mapsto \mathbb{R}$  defined for all  $x \in \mathcal{X}$  by

$$\tilde{f}(x) = \begin{cases} f^*(x) + 3 \left( 1 - \frac{\|c^* - x\|_2}{\epsilon_1} \right) \times (\max_{x \in \mathcal{X}} f^*(x) - \min_{x \in \mathcal{X}} f^*(x)) & \text{if } x \in B(c^*, \epsilon_1) \\ f^*(x) & \text{otherwise} \end{cases}$$

which is maximized over  $B(c^*, \epsilon_1)$  and Lipschitz continuous as both  $f^*$  and  $x \mapsto \|c^* - x\|_2$  are Lipschitz. However, since  $\tilde{f}$  and  $f^*$  can not be distinguished over  $\mathcal{X}/B(c, \epsilon_1)$ , we have that  $\forall n \in \mathbb{N}^*$ ,

$$\begin{aligned}
 \mathbb{P} \left( \max_{x \in \mathcal{X}} \tilde{f}(x) - \max_{i=1 \dots n} \tilde{f}(X'_i) > \max_{x \in \mathcal{X}} f(x) \right) &\geq \mathbb{P} \left( \bigcap_{i=1}^n \{X'_i \notin B(c, \epsilon_2) \cap \mathcal{X}\} \right) \\
 &= \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin B(c, \epsilon_2) \cap \mathcal{X}\} \right) \\
 &\geq \epsilon_2 / (2N_1) \\
 &> 0
 \end{aligned}$$

where  $X'_1, \dots, X'_n$  denotes a sequence of evaluation points generated by  $A$  over  $\tilde{f}$ , and we deduce that there exists  $\tilde{f} \in \bigcup_{k \geq 0} \text{Lip}(k)$  such that  $\max_{i=1 \dots n} \tilde{f}(X'_i) \xrightarrow{P} \max_{x \in \mathcal{X}} \tilde{f}(x)$ . Hence, it contradicts the fact that  $A$  is consistent over  $\bigcup_{k \geq 0} \text{Lip}(k)$  and we deduce that, necessarily,  $\sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|X_i - x\| \xrightarrow{P} 0$  for all  $f \in \bigcup_{k \geq 0} \text{Lip}(k)$ .  $\square$

**Proof of Example 4.** Fix any  $n \in \mathbb{N}^*$ , set  $\delta \in (0, 1)$ , define  $\epsilon = \text{diam}(\mathcal{X}) \cdot ((\ln(n/\delta) + d \ln(d))/n)^{1/d}$  and let  $X_1, \dots, X_n$  be a sequence of  $n$  independent copies of

$X \sim \mathcal{U}(\mathcal{X})$ . Since the result trivially holds whenever  $(\ln(n/\delta) + d \ln(d))/n \geq 1$ , we consider the case where  $(\ln(n/\delta) + d \ln(d))/n < 1$ . From Proposition 26, we know that there exists a sequence  $x_1, \dots, x_{N_\epsilon}$  of  $N_\epsilon = \mathcal{N}_\epsilon(\mathcal{X})$  points in  $\mathcal{X}$  such that  $\mathcal{X} \subseteq \bigcup_{j=1}^{N_\epsilon} B(x_j, \epsilon)$ . Therefore, using the bound on the covering number  $\mathcal{N}_\epsilon(\mathcal{X})$  of Corollary 26, we obtain that

$$\begin{aligned}
\mathbb{P} \left( \sup_{x \in \mathcal{X}} \min_{i=1 \dots n} \|x - X_i\|_2 \leq \epsilon \right) &\geq \mathbb{P} \left( \bigcap_{j=1}^{N_\epsilon} \bigcup_{i=1}^n \{X_i \in B(x_j, \epsilon) \cap \mathcal{X}\} \right) \\
&= 1 - \mathbb{P} \left( \bigcup_{j=1}^{N_\epsilon} \bigcap_{i=1}^n \{X_i \notin B(x_j, \epsilon) \cap \mathcal{X}\} \right) \\
&\geq 1 - \sum_{j=1}^{N_\epsilon} \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin B(x_j, \epsilon) \cap \mathcal{X}\} \right) \\
&\geq 1 - N_\epsilon \times \max_{j=1 \dots N_\epsilon} \mathbb{P}(X_1 \notin B(x_j, \epsilon) \cap \mathcal{X})^n \\
&= 1 - N_\epsilon \times \max_{j=1 \dots N_\epsilon} \left( 1 - \frac{\mu(\mathcal{X} \cap B(x_j, \epsilon))}{\mu(\mathcal{X})} \right)^n \\
&\geq 1 - N_\epsilon \times \left( 1 - \left( \frac{\epsilon}{\text{diam}(\mathcal{X})} \right)^d \right)^n \\
&\geq 1 - \left( \frac{\sqrt{d} \text{diam}(\mathcal{X})}{\epsilon} \right)^d \times \left( 1 - \left( \frac{\epsilon}{\text{diam}(\mathcal{X})} \right)^d \right)^n \\
&\geq 1 - \delta
\end{aligned}$$

and the proof is complete.  $\square$

**Proof of Proposition 5.** The proof heavily builds upon the arguments used in the proof of the Theorem 1 in (Bull (2011)). Pick any algorithm  $A \in \mathcal{A}$  and any constant  $C > 0$ . Fix any  $n \in \mathbb{N}^*$  and  $\delta \in (0, 1)$  and set  $N_\delta = \lceil (n/\delta)^{1/d} \rceil$ . By definition of  $\text{rad}(\mathcal{X})$ , we know there exists some  $x \in \mathcal{X}$  such that  $x + [0, 2\text{rad}(\mathcal{X})/\sqrt{d}]^d \subseteq \mathcal{X}$ . One can then define for all  $I \in \{1, \dots, N_\delta\}^d$ , the centers  $c_I$  of the hypercubes  $H_I$  whose side are equal to  $D = 2\text{rad}(\mathcal{X})/(\sqrt{d}N_\delta)$  and cover  $\mathcal{X}$ , i.e.,  $\bigcup_I H_I = x + [0, 2\text{rad}(\mathcal{X})/\sqrt{d}]^d \subseteq \mathcal{X}$ . Now, let  $X_1, \dots, X_n$  be a sequence of  $n$  evaluation points generated by the algorithm  $A$  over the constant function  $f_0 : x \in \mathcal{X} \mapsto 0$  and define for all  $I \in \{1, \dots, N_\delta\}^d$  the event

$$E_I = \bigcap_{i=1}^n \{X_i \notin \text{Int}(H_I)\}.$$

As the interiors of the  $N_\delta^d$  hypercubes are disjoint and we have  $n$  points, it necessarily follows that

$$N_\delta^d \times \max_I \mathbb{P}(E_I) \geq \sum_I \mathbb{P}(E_I) = \mathbb{E} \left[ \sum_I \mathbb{I}\{E_I\} \right] \geq N_\delta^d - n.$$

Hence, there exists some fixed  $I^\star$  only depending on  $\mathcal{A}$  which maximizes the above probability and thus satisfies

$$\mathbb{P}(E_{I^\star}) \geq \frac{N_\delta^d - n}{N_\delta^d} = 1 - \frac{n}{\lceil (n/\delta)^{1/d} \rceil^d} \geq 1 - \delta.$$

Now, using the center  $c_{I^\star}$  of the hypercube  $H_{I^\star}$ , one can then introduce the function  $\tilde{f} \in \bigcup_{k \geq 0} \text{Lip}(k)$  defined for all  $x \in \mathcal{X}$  by

$$\tilde{f}(x) = \begin{cases} C \times (1 - 2\|c_{I^\star} - x\|_2 / D) & \text{if } \|c_{I^\star} - x\|_2 \leq D/2 \\ 0 & \text{otherwise.} \end{cases}$$

However, since the functions  $\tilde{f}$  and  $f_0$  can not be distinguished over  $\mathcal{X}/H_{I^\star}$ , we have that

$$\mathbb{P}\left(\max_{x \in X} \tilde{f}(x) - \max_{i=1 \dots n} \tilde{f}(X'_i) \geq C\right) \geq \mathbb{P}\left(\bigcap_{i=1}^n \{X'_i \notin \text{Int}(H_{I^\star})\}\right) = \mathbb{P}(E_{I^\star}) \geq 1 - \delta$$

where  $X'_1, \dots, X'_n$  denotes a sequence of evaluation points generated by  $\mathcal{A}$  over  $\tilde{f}$ , which proves the result.  $\square$

**Proof of Proposition 6. (Lower bound).** Pick any  $n \in \mathbb{N}^\star$  and set  $D = 2 \text{rad}(\mathcal{X}) / (\sqrt{d} \lceil (2n)^{1/d} \rceil)$ . It can easily be shown by reproducing the same steps as in the proof of Proposition 5 with  $\delta$  set to  $1/2$ , that for any global optimization algorithm  $A$ , there exists a function  $\tilde{f}_A \in \text{Lip}(k)$  defined by

$$\tilde{f}_A(x) = \begin{cases} kD/2 - k \cdot \|c_A - x\|_2 & \text{if } \|c_A - x\|_2 \leq D/2 \\ 0 & \text{otherwise,} \end{cases}$$

for some center  $c_A \in \mathcal{X}$  only depending on  $A$ , for which we have  $\mathbb{P}(\max_{x \in X} \tilde{f}_A(x) - \max_{i=1 \dots n} \tilde{f}_A(X_i) \geq k \cdot D/2) \geq 1/2$  where  $X_1, \dots, X_n$  is a sequence of  $n$  evaluation points generated by  $A$  over  $\tilde{f}_A$ . Therefore, using the definition of the supremum and Markov's inequality gives that  $\forall A \in \mathcal{A}$ :

$$\begin{aligned} \sup_{f \in \text{Lip}(k)} \mathbb{E} \left[ \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) \right] &\geq \mathbb{E} \left[ \max_{x \in \mathcal{X}} \tilde{f}_A(x) - \max_{i=1 \dots n} \tilde{f}_A(X_i) \right] \\ &\geq \frac{kD}{2} \times \mathbb{P} \left( \max_{x \in X} \tilde{f}_A(x) - \max_{i=1 \dots n} \tilde{f}_A(X_i) \geq k \cdot \frac{D}{2} \right) \\ &\geq k \cdot \frac{\text{rad}(\mathcal{X})}{8\sqrt{d}} \cdot n^{-\frac{1}{d}}. \end{aligned}$$

As the previous inequality holds true for any algorithm  $A$ , the proof is complete.

**(Upper bound).** Sequentially using the fact that (i) the infimum minimax loss taken over all algorithms is necessarily upper bounded by the loss suffered by a Pure Random

Search, (ii) for any positive random variable,  $\mathbb{E}[X] = \int_{t=0}^{\infty} \mathbb{P}(X \geq t) dt$ , (iii) Proposition 28 and (iv) the change of variable  $u = n(t/\text{diam}(\mathcal{X}))^{1/d}$ , we obtain that

$$\begin{aligned} \inf_{A \in \mathcal{A}} \sup_{f \in \text{Lip}(k)} \mathbb{E} \left[ \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) \right] &\leq \sup_{f \in \text{Lip}(k)} \mathbb{E} \left[ \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X'_i) \right] \\ &\leq \int_0^{\infty} \exp \left\{ -n(t/k \cdot \text{diam}(\mathcal{X}))^{1/d} \right\} dt \\ &= k \cdot \text{diam}(\mathcal{X}) \cdot n^{-d} \cdot d \cdot \int_0^{\infty} u^{d-1} e^{-u} du \\ &= k \cdot \text{diam}(\mathcal{X}) \cdot n^{-d} \cdot d \cdot \Gamma(d) \end{aligned}$$

where  $X'_1, \dots, X'_n$  denotes a sequence of  $n$  independent copies of  $X' \sim \mathcal{U}(\mathcal{X})$  and  $\Gamma(\cdot)$  the Euler's Gamma function. Recalling that  $\Gamma(d) = (d-1)!$  for all  $d \in \mathbb{N}^*$  completes the proof.  $\square$

### Appendix C. Proofs of Section 3

In this section, we provide the proofs for Lemma 9, Proposition 11, Proposition 12, Corollary 13, Proposition 14, Theorem 15 and Theorem 16.

**Proof of Lemma 9.** The first implication ( $\Rightarrow$ ) is a direct consequence of the definition of  $\mathcal{X}_{k,t}$ . Noticing that the function  $\hat{f}: x \mapsto \min(\max_{i=1 \dots t} f(X_i), \min_{i=1 \dots t} f(X_i) + k \|x - X_i\|_2)$  belongs to  $\mathcal{F}_{k,t}$  and that  $\arg \max_{x \in \mathcal{X}} \hat{f}(x) = \{x \in \mathcal{X} : \min_{i=1 \dots t} f(X_i) + k \|x - X_i\|_2 \geq \max_{i=1 \dots t} f(X_i)\}$  proves the second implication.  $\square$

**Proof of Proposition 11.** Fix any  $f \in \text{Lip}(k)$ , pick any  $n \in \mathbb{N}^*$ , set  $\epsilon > 0$  and let  $\mathcal{X}_\epsilon = \{x \in \mathcal{X} : f(x) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon\}$  be the corresponding level set. Denoting by  $X'_1, \dots, X'_n$  a sequence of  $n$  random variable uniformly distributed over  $\mathcal{X}$  and observing that  $\mu(\mathcal{X}_\epsilon) > 0$ , we directly obtain from Proposition 12 that

$$\begin{aligned} \mathbb{P} \left( \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) > \epsilon \right) &\leq \mathbb{P} \left( \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X'_i) > \epsilon \right) \\ &= \mathbb{P} \left( \bigcap_{i=1}^n \{X'_i \notin \mathcal{X}_\epsilon\} \right) \\ &\leq \left( 1 - \frac{\mu(\mathcal{X}_\epsilon)}{\mu(\mathcal{X})} \right)^n \xrightarrow{n \rightarrow \infty} 0. \end{aligned}$$

$\square$

**Proof of Proposition 12.** The proof is similar to the one of Proposition 12 in (Malherbe and Vayatis (2016)).  $\square$

**Proof of Corollary 13.** Combining Proposition 12 and Proposition 28 stated at the beginning of the Appendix Section gives the result.  $\square$

**Proof of Proposition 14.** Fix any  $\delta \in (0, 1)$ , set  $n \in \mathbb{N}^*$  and let  $r_{\delta,n} = \text{rad}(\mathcal{X}) (\delta/n)^{\frac{1}{d}}$  be

the value of the lower bound divided by  $k$ . As  $\text{rad}(\mathcal{X}) > 0$ , there necessarily exists some point  $x^* \in \mathcal{X}$  such that  $B(x^*, \text{rad}(\mathcal{X})) \subseteq \mathcal{X}$ . Based on this point, one can then introduce the function  $\tilde{f} \in \text{Lip}(k)$  defined for all  $x \in \mathcal{X}$  by

$$\tilde{f}(x) = \begin{cases} k \cdot r_{\delta,n} - k \cdot \|x - x^*\|_2 & \text{if } x \in B(x^*, r_{\delta,n}) \\ 0 & \text{otherwise.} \end{cases}$$

Denoting now by  $X_1, \dots, X_n$  a sequence of  $n$  evaluation points generated by LIPO tuned with a parameter  $k$  over  $\tilde{f}$  and observing that (i)  $X_1$  is uniformly distributed over  $\mathcal{X}$  and (ii)  $X_{i+1}$  is also uniformly distributed over  $\mathcal{X}$  for  $i \geq 1$  as soon as only constant evaluations have been recorded (*i.e.*  $\mathcal{X}_{k,i+1} = \mathcal{X}$  on the event  $\bigcap_{t \leq i} \{X_t \notin B(x^*, r_{\delta,n})\}$ ), we have that

$$\begin{aligned} \mathbb{P} \left( \max_{x \in \mathcal{X}} \tilde{f}(x) - \max_{i=1 \dots n} \tilde{f}(X_i) \geq k \cdot r_{\delta,n} \right) &\geq \mathbb{P} \left( \bigcap_{i=1}^n \{X_i \notin B(x^*, r_{\delta,n})\} \right) \\ &= \left[ \mathbb{P}(X_1 \notin B(x^*, r_{\delta,n})) \times \right. \\ &\quad \left. \prod_{i=1}^{n-1} \mathbb{P} \left( X_{i+1} \notin B(x^*, r_{\delta,n}) \mid \bigcap_{t=1}^i \{X_t \notin B(x^*, r_{\delta,n})\} \right) \right] \\ &= \left( 1 - \frac{\mu(B(x^*, r_{\delta,n}) \cap \mathcal{X})}{\mu(\mathcal{X})} \right)^n \\ &\geq \left( 1 - \left( \frac{r_{\delta,n}}{\text{rad}(\mathcal{X})} \right)^d \right)^n \\ &= \left( 1 - \frac{\delta}{n} \right)^n \\ &\geq 1 - \delta. \end{aligned}$$

□

**Proof of Theorem 15.** Pick any  $n \in \mathbb{N}^*$ , fix any  $\delta \in (0, 1)$  and let  $X_1, \dots, X_n$  be a sequence of  $n$  evaluation points generated by the LIPO algorithm over  $f$  after  $n$  iterations. To clarify the proof, we set some specific notations: let  $D = \max_{x \in \mathcal{X}} \|x - x^*\|_2$ , set

$$M = \begin{cases} \left\lfloor \left( \frac{c_\kappa}{8k} \right)^d \cdot \frac{n}{\ln(n/\delta) + 2(2\sqrt{d})^d} \right\rfloor & \text{if } \kappa = 1 \\ \left\lfloor \frac{1}{\ln(2)d(\kappa-1)} \ln \left( 1 + \left( \frac{c_\kappa D^\kappa}{8kD} \right)^d \frac{n(2^{d(\kappa-1)} - 1)}{\ln(n/\delta) + 2(2\sqrt{d})^d} \right) \right\rfloor & \text{otherwise,} \end{cases}$$

define for all  $m \in \{1 \dots M\}$  the series of integers:

$$N_m := \left\lceil \sqrt{d} \cdot \left( \frac{8kD}{c_\kappa D^\kappa} \right) \cdot 2^{m(\kappa-1)} \right\rceil^d \quad \text{and} \quad N'_m := \left\lceil \ln(M/\delta) \cdot \left( \frac{8kD}{c_\kappa D^\kappa} \right)^d \cdot 2^{md(\kappa-1)} \right\rceil$$

and let  $\tau_0, \dots, \tau_M$  be the series of stopping times initialized by  $\tau_0 = 0$  and defined for all  $m \geq 1$  by

$$\tau_m := \inf \left\{ t \geq \tau_{m-1} \mid \sum_{i=\tau_{m-1}+1}^t \mathbb{I}\{X_i \in B(x^*, 2 \cdot D \cdot 2^{-m})\} = N'_m \right\}.$$

The stopping time  $\tau_m$  correspond to the time after  $\tau_{m-1}$  where we have recorded at least  $N'_m$  random evaluation points inside the ball  $B(x^*, 2 \cdot D \cdot 2^{-m})$ . To prove the result, we show that each of the following events:

$$E_m := \left\{ \max_{i=1 \dots \tau_m} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - \frac{c_\kappa}{2} \cdot \left( \frac{D}{2^m} \right)^\kappa \right\} \cap \left\{ \tau_m \leq N'_1 + \sum_{l=1}^{m-1} (N'_{l+1} + N_l) \right\}.$$

holds true with probability at least  $1 - \delta/M$  on the event  $\cap_{l=1}^{m-1} E_l$  for all  $m \in \{2, \dots, M\}$  so that:

$$\mathbb{P}(E_M) \geq \mathbb{P}(E_1) \times \prod_{m=1}^{M-1} \mathbb{P} \left( E_{m+1} \mid \bigcap_{l=1}^m E_l \right) \geq \left( 1 - \frac{\delta}{M} \right)^M \geq 1 - \delta \quad (4)$$

that will leads us to the result by analyzing  $E_M$ .

**Analysis of  $\mathbb{P}(E_1)$ .** Observe first that since  $\mathcal{X} \subseteq B(x^*, D)$ , then  $\tau_1 = N'_1$ . Using now the fact that (i) the algorithm is faster than a Pure Random Search (Proposition 12) and (ii) the bound of Proposition 28, we directly get that with probability at least  $1 - \delta/M$ ,

$$\begin{aligned} \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots \tau_1} f(X_i) &\leq k \cdot 2D \cdot \left( \frac{\ln(M/\delta)}{N'_1} \right)^{\frac{1}{d}} \\ &\leq k \cdot 2D \cdot \left( \frac{\ln(M/\delta)}{\ln(M/\delta) 2^{d(\kappa-1)}} \left( \frac{c_\kappa D^\kappa}{8kD} \right)^d \right)^{\frac{1}{d}} \\ &= \frac{c_\kappa}{2} \cdot \left( \frac{D}{2} \right)^\kappa \end{aligned}$$

which proves that  $\mathbb{P}(E_1) \geq 1 - \delta/M$ .

**Analysis of  $\mathbb{P}(E_{m+1} \mid \cap_{l=1}^m E_l)$ .** To bound this term, we use (i) a deterministic covering argument to control the stopping time  $\tau_{m+1}$  (Lemma 29 and Corollary 30) and (ii) a stochastic argument to bound the maximum  $\max_{i=1 \dots \tau_{m+1}} f(X_i)$  (Lemma 31 and Corollary 32). The following lemma states that after  $\tau_m$  and on the event  $E_m$  there will be at most  $N_m$  evaluation points that will fall inside the area  $B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m})$ .

**Lemma 29** *For all  $m \in \{1, \dots, M-1\}$ , we have on the event  $E_m$ ,*

$$\sum_{t=\tau_m+1}^n \mathbb{I}\{X_t \in B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m})\} \leq N_m.$$

**Proof** Fix  $m \in \{1, \dots, M-1\}$  and assume that  $E_m = \{\max_{i=1 \dots \tau_m} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - c_k/2 \cdot (D/2^m)^\kappa\}$  holds true. Setting  $N = \lceil \sqrt{d} 8k D 2^{m(\kappa-1)} / (c_k D^\kappa) \rceil$  and observing that  $B(x^*, 2D \cdot 2^{-m}) \subseteq x^* + 2D \cdot 2^{-m} \cdot [-1, +1]^d$ , one can then introduce the sequence  $H_I$ , with  $I \in \{1, \dots, N\}^d$ , of the  $N^d = N_m$  hypercubes whose side have length  $4D \cdot 2^{-m}/N$  and cover  $x^* + 2D \cdot 2^{-m} \times [-1, +1]^d$ , so that

$$B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m}) \subseteq x^* + 2D \cdot 2^{-m} \cdot [-1, +1]^d = \bigcup_I H_I.$$

Based on these hypercubes, one can define the set

$$I_t = \{I \in \{1, \dots, N\}^d : H_I \cap B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m}) \cap \mathcal{X}_{k,t} \neq \emptyset\}$$

which contains the indexes of the hypercubes that still intersect the set of potential maximizers  $\mathcal{X}_{k,t}$  at time  $t$  and the target area  $B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m})$ . We show by contradiction that there cannot be more than  $N^d = N_m$  evaluation points falling inside this area, otherwise it would be empty. Suppose that, after  $\tau_m$ , there exists a sequence

$$\tau_m < t_1 < t_2 < \dots < t_{N^d+1} \leq n$$

of  $N^d + 1$  strictly increasing indexes for which the evaluation points  $X_{t_j}$ ,  $j \geq 1$ , belong to the target area, *i.e.*,

$$\forall j \in \{1, \dots, N^d + 1\}, X_{t_j} \in B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m}).$$

Fix any  $j \geq 1$  and observe that since  $X_{t_j} \notin B(x^*, D \cdot 2^{-m})$ , then we have from Condition 1 that (i)  $f(X_{t_j}) < \max_{x \in \mathcal{X}} f(x) - c_k \cdot (D \cdot 2^{-m})^\kappa$ . Moreover, as  $X_{t_j} \in \mathcal{X}_{k,t_j-1} \cap B(x^*, 2D \cdot 2^{-m})/B(x^*, D \cdot 2^{-m})$ , it necessarily follows from the definition of the algorithm that (ii) there exists an index  $I^* \in I_{t_j-1}$  such that  $X_{t_j} \in H_{I^*}$ . Therefore, combining (i) and (ii) with  $E_m$ , gives that  $\forall x \in H_{I^*}$ :

$$\begin{aligned} f(x) &\leq f(X_{t_j}) + k \cdot \|X_{t_j} - x\|_2 && (f \in \text{Lip}(k)) \\ &\leq f(X_{t_j}) + k \cdot \max_{(x,x') \in H_I^2} \|x - x'\|_2 && ((X_{t_j}, x) \in H_I^2) \\ &= f(X_{t_j}) + k \cdot \sqrt{d} \cdot 4D \cdot 2^{-m}/N && (\text{def. of } H_I) \\ &\leq f(X_{t_j}) + \frac{c_k}{2} \cdot (D \cdot 2^{-m})^\kappa && (\text{def. of } N) \\ &< \max_{x \in \mathcal{X}} f(x) - c_k \cdot (D \cdot 2^{-m})^\kappa + \frac{c_k}{2} \cdot (D \cdot 2^{-m})^\kappa && (i) \\ &\leq \max_{i=1 \dots \tau_m} f(X_i) && (E_1) \\ &\leq \max_{i=1 \dots t_j} f(X_i). && (t_j > \tau_m) \end{aligned}$$

It has been shown that if  $X_{t_j}$  belongs to the target area then  $f(x) < \max_{i=1 \dots t_j} f(X_i)$  for all  $x \in H_{I^*}$ , which combined with the definition of the set of potential maximizers  $\mathcal{X}_{k,t_j}$  at time  $t_j$  implies that  $H_{I^*} \notin \mathcal{X}_{\tau_j}$ . Hence, once an evaluation has been made in  $H_{I^*}$ , there will not be

any future evaluation point falling inside this cube. We thus deduce that  $|I_{t_j}| \leq |I_{t_{j-1}}| - 1$  for all  $j \geq 1$  which leads us to the following contradiction:

$$0 \leq |I_{t_{N^d+1}}| = |I_{\tau_m}| + \sum_{j=\tau_m+1}^{t_{N^d+1}} |I_{t_j}| - |I_{t_{j-1}}| \leq |I_{\tau_m}| - (N^d + 1) \leq N^d - (N^d + 1) < 0$$

and proves the statement.  $\square$

Based on this lemma, one might then derive a bound on the stopping time  $\tau_{m+1}$ .

**Corollary 30** *For all  $m \in \{1, \dots, M-1\}$ , we have on the event  $\bigcap_{l=1}^m E_l$  that*

$$\tau_{m+1} \leq N'_1 + \sum_{l=1}^m (N'_{l+1} + N_l).$$

**Proof** The result is proved by induction. We start with the case where  $m = 1$ . Assuming that  $E_1$  holds true and observing that (i)  $\tau_1 = N'_1$  and (ii)  $\mathcal{X} \subseteq B(x^*, D) = B(x^*, D/2) \cup B(x^*, D)/B(x^*, D/2)$ , one can then write:

$$\begin{aligned} \tau_2 &= \tau_1 + \sum_{i=\tau_1+1}^{\tau_2} \mathbb{I}\{X_i \in B(x^*, D)\} \\ &= N'_1 + \sum_{i=\tau_1+1}^{\tau_2} \mathbb{I}\{X_i \in B(x^*, D/2)\} + \sum_{i=\tau_1+1}^{\tau_2} \mathbb{I}\{X_i \in B(x^*, D)/B(x^*, D/2)\}. \end{aligned}$$

However, since (i)  $\sum_{i=\tau_1+1}^{\tau_2} \mathbb{I}\{X_i \in B(x^*, D/2)\} = N'_2$  by definition of  $\tau_2$  and (ii)  $\sum_{i=\tau_1+1}^{\tau_2} \mathbb{I}\{X_i \in B(x^*, D)/B(x^*, D/2)\} \leq N_1$  by Lemma 29, the result holds true for  $m = 1$ . Consider now any  $m \geq 2$  and assume that the statement holds true for all  $l < m$ . Again, observing that  $\mathcal{X} \subseteq B(x^*, D \cdot 2^{-m}) \cup \bigcup_{l=1}^m B(x^*, D \cdot 2^{-(l-1)})/B(x^*, D \cdot 2^{-l})$  and keeping in mind that the stopping times are bounded by the induction assumption, one can write

$$\begin{aligned} \tau_{m+1} &= \tau_m + \sum_{i=\tau_m+1}^{\tau_{m+1}} \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-m})\} \\ &\quad + \sum_{i=\tau_m+1}^{\tau_{m+1}} \sum_{l=1}^m \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-(l-1)})/B(x^*, D \cdot 2^{-l})\}. \end{aligned}$$

Now, combining the telescopic representation  $\tau_{m+1} = \tau_1 + \sum_{l=1}^m (\tau_{l+1} - \tau_l)$  with the previous decomposition gives that

$$\begin{aligned} \tau_{m+1} &= \tau_1 + \sum_{l=1}^m \sum_{i=\tau_l+1}^{\tau_{l+1}} \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-l})\} \\ &\quad + \sum_{l=1}^m \sum_{i=\tau_l+1}^{\tau_{m+1}} \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-(l-1)})/B(x^*, D \cdot 2^{-l})\}. \end{aligned}$$

However, since (i)  $\tau_1 = N'_1$ , (ii)  $\sum_{i=\tau_l+1}^{\tau_{l+1}} \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-l})\} = N'_{l+1}$ , for all  $l \geq 1$  by definition of the stopping times and (iii)  $\sum_{i=\tau_l+1}^{\tau_{l+1}} \mathbb{I}\{X_i \in B(x^*, 2^{-(l-1)})/B(x^*, 2^{-l})\} \leq N_l$ , for all  $l \geq 1$  on the event  $\bigcap_{l=1}^m E_l$ , from Lemma 29, we finally get that

$$\tau_{m+1} \leq N'_1 + \sum_{l=1}^m (N'_{l+1} + N_l).$$

□

As Corollary 30 gives the desired bound on  $\tau_{m+1}$ , it remains to control the maximum  $\max_{i=1 \dots \tau_{m+1}} f(X_i)$ . The next lemma shows that i.i.d. results can actually be used to bound this term.

**Lemma 31** *For all  $m \in \{1, \dots, M-1\}$ , we have that  $\forall y \in \text{Im}(f)$ ,*

$$\mathbb{P}\left(\max_{i=1 \dots \tau_{m+1}} f(X_i) \geq y \mid \bigcap_{l=1}^m E_l\right) \geq \mathbb{P}\left(\max_{i=1 \dots N'_{m+1}} f(X'_i) \geq y\right).$$

where  $X'_1 \dots X'_{N'_{m+1}}$  denotes a sequence  $N'_{m+1}$  i.i.d. copies of  $X' \sim \mathcal{U}(\mathcal{X} \cap B(x^*, D \cdot 2^{-m}))$ .

**Proof** From Corollary 30, we know that on the event  $\bigcap_{l=1}^m E_l$  the stopping time  $\tau_{m+1}$  is finite. Moreover, as  $\sum_{i=\tau_m+1}^{\tau_{m+1}} \mathbb{I}\{X_i \in B(x^*, D \cdot 2^{-m})\} = N'_{m+1}$  by definition of  $\tau_{m+1}$ , it can then easily be shown by reproducing the same steps as in the proof of Proposition 12 with the evaluations points falling into  $B(x^*, D \cdot 2^{-m})$  after  $\tau_m$  that the algorithm is faster than a Pure Random Search performed over the subspace  $\mathcal{X} \cap B(x^*, D \cdot 2^{-m})$ , which proves the result. □

As a direct consequence of this lemma, one can get the desired bound on the maxima as shown in the next corollary.

**Corollary 32** *For all  $m \in \{1, \dots, M-1\}$ , we have that*

$$\mathbb{P}\left(\max_{i=1 \dots \tau_{m+1}} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - \frac{c_k}{2} \cdot \left(\frac{D}{2^{m+1}}\right)^\kappa \mid \bigcap_{l=1}^m E_l\right) \geq 1 - \delta/M.$$

**Proof** Omitting the conditionning upon  $\bigcap_{l=1}^m E_l$ , we obtain from the combination of Lemma 31 and Proposition 28 that with probability at least  $1 - \delta/M$ :

$$\begin{aligned} \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots \tau_{m+1}} f(X_i) &\leq k \cdot 2D \cdot 2^{-m} \cdot \left(\frac{\ln(M/\delta)}{N'_{m+1}}\right)^{\frac{1}{d}} \\ &\leq k \cdot 2D \cdot 2^{-m} \cdot \left(\frac{\ln(M/\delta)}{\ln(M/\delta) 2^{d(m+1)(\kappa-1)}} \left(\frac{c_\kappa D^\kappa}{8kD}\right)^d\right)^{\frac{1}{d}} \\ &= \frac{c_\kappa}{2} \cdot \left(\frac{D}{2^{m+1}}\right)^\kappa. \end{aligned}$$

□

At this point, we know from the combination of Corollary 30 and Corollary 32 that

$$\forall m \in \{1, \dots, M-1\}, \quad \mathbb{P} \left( E_{m+1} \mid \bigcap_{l=1}^m E_l \right) \geq 1 - \delta/M$$

which proves from (4) that  $\mathbb{P}(E_M) \geq 1 - \delta$ .

**Analysis of  $E_M$ .** As  $\max_{i=1 \dots \tau_M} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - \frac{c_\kappa}{2} \cdot D^\kappa \cdot 2^{-M\kappa}$  and  $\tau_M \leq N'_1 + \sum_{l=1}^{M-1} (N'_{l+1} + N_l)$  on the event  $E_M$ , it remains to show that  $N'_1 + \sum_{l=1}^{M-1} (N'_{l+1} + N_l) \leq n$  to conclude the proof. Consider first the case  $\kappa = 1$ . Setting  $C = (8k/c_\kappa)^d$  and observing that (i)  $N'_l \leq \ln(M/\delta)C + 1$ , (ii)  $N_l \leq 2 \cdot C \cdot (2\sqrt{d})^d - 1$  for all  $l \leq M$  and (iii)  $M \leq n$ , one gets:

$$\begin{aligned} N'_1 + \sum_{l=1}^{M-1} (N'_{l+1} + N_l) &\leq C \cdot M \left( \ln(M/\delta) + 2(2\sqrt{d})^d \right) \\ &\leq n \cdot \frac{\ln(M/\delta) + 2(2\sqrt{d})^d}{\ln(n/\delta) + 2(2\sqrt{d})^d} \\ &\leq n. \end{aligned}$$

For  $\kappa > 1$ , since (i)  $M$  was chosen so that  $\frac{2^{d(\kappa-1)M}-1}{2^{d(\kappa-1)}-1} \leq \frac{n}{C} \cdot \frac{1}{\ln(n/\delta) + 2(2\sqrt{d})^d}$  and (ii)  $M \leq n$ , we obtain:

$$\begin{aligned} N'_1 + \sum_{l=1}^{M-1} (N'_{l+1} + N_l) &\leq C \cdot \left( \ln(M/\delta) + 2(2\sqrt{d})^d \right) \sum_{l=1}^M (2^{d(\kappa-1)})^l \\ &\leq C \cdot \left( \ln(M/\delta) + 2(2\sqrt{d})^d \right) \cdot \frac{2^{d(\kappa-1)M} - 1}{2^{d(\kappa-1)} - 1} \\ &\leq n. \end{aligned}$$

Finally, using the elementary inequality  $\lfloor x \rfloor \geq x - 1$  over  $M$  and the inequality  $c_\kappa D^\kappa \leq k \text{diam}(\mathcal{X})$  (by Condition 1) leads to the desired result and completes the proof. □

**Proof of Theorem 16.** (LOWER BOUND) Pick any  $n \in \mathbb{N}^*$  and  $\delta \in (0, 1)$ , set  $\epsilon = c_\kappa \text{rad}(\mathcal{X})^\kappa \delta^{\kappa/d} \exp(-\kappa(n - \sqrt{2n \ln(1/\delta)})/d)$ , let  $\mathcal{X}_\epsilon = \{x \in \mathcal{X} : f(x) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon\}$  be the corresponding level set. Observe first that since (i)  $\mathcal{X}_\epsilon = \{x \in \mathcal{X} : \epsilon \geq f(x^*) - f(x)\} \subseteq \{x \in \mathcal{X} : \epsilon \geq c_\kappa \|x^* - x\|_2^\kappa\} = \mathcal{X} \cap B(x^*, (\epsilon/c_\kappa)^{1/\kappa})$  and (ii) there exists  $x \in \mathcal{X}$  such that  $B(x, \text{rad}(\mathcal{X})) \subseteq \mathcal{X}$ , then  $\mu(\mathcal{X}_\epsilon)/\mu(\mathcal{X}) \leq ((\epsilon/c_\kappa)^{1/\kappa}/\text{rad}(\mathcal{X}))^d = \delta e^{-n - \sqrt{2n \ln(1/\delta)}}$ . It can then easily be shown by reproducing the same steps as in the proof of the Lower bound of

Theorem 17 in (Malherbe and Vayatis (2016)) that

$$\begin{aligned}
 \mathbb{P}\left(\max_{i=1\dots n} f(X_i) \geq \max_{x \in \mathcal{X}} f(x) - \epsilon\right) &= \mathbb{P}\left(\frac{\mu(\{x \in \mathcal{X} : f(x) \geq \max_{i=1\dots n} f(X_i)\})}{\mu(\mathcal{X})} \leq \frac{\mu(\mathcal{X}_\epsilon)}{\mu(\mathcal{X})}\right) \\
 &\leq \mathbb{P}\left(\prod_{i=1}^n U_i \leq \frac{\mu(\mathcal{X}_\epsilon)}{\mu(\mathcal{X})}\right) \\
 &\leq \mathbb{P}\left(\prod_{i=1}^n U_i \leq \delta \cdot e^{-n - \sqrt{2n \ln(1/\delta)}}\right) \\
 &= \mathbb{P}\left(\sum_{i=1}^n -\ln(U_i) > n + \sqrt{2n \ln(1/\delta)} + \ln(1/\delta)\right) \\
 &\leq \delta
 \end{aligned}$$

where  $U_1, \dots, U_n$  denotes a sequence of  $n$  i.i.d. copies of  $U \sim \mathcal{U}([0, 1])$ . We point out that a concentration results for gamma random variable was used on the last line (see Lemma 37 and Lemma 38 in (Malherbe and Vayatis (2016)) for more details).  $\square$

#### Appendix D. Analysis of AdaLipOpt (proofs of Section 4)

**Proof of Proposition 18.** Pick any  $t \geq 2$ , consider any non-constant  $f \in \bigcup_{k \geq 0} \text{Lip}(k)$  and set  $i^* = \min\{i \in \mathbb{Z} : f \in \text{Lip}(k_i)\}$ . To prove the result, we decorrelate the sample and use the fact that  $(X_1, X_{\lfloor t/2 \rfloor + 1}, \dots, (X_{\lfloor t/2 \rfloor}, X_{2\lfloor t/2 \rfloor}))$  forms a sequence of  $\lfloor t/2 \rfloor$  i.i.d. copies of  $(X, X') \sim \mathcal{U}(\mathcal{X} \times \mathcal{X})$ :

$$\begin{aligned}
 \mathbb{P}(f \in \text{Lip}(\hat{k}_t)) &= \mathbb{P}(\hat{k}_t = k_{i^*}) \\
 &= \mathbb{P}\left(\bigcup_{i \neq j}^t \{|f(X_i) - f(X_j)| > k_{i^*-1} \cdot \|X_i - X_j\|_2\}\right) \\
 &\geq \mathbb{P}\left(\bigcup_{i=1}^{\lfloor t/2 \rfloor} \left\{|f(X_i) - f(X_{\lfloor t/2 \rfloor + i})| > k_{i^*-1} \cdot \|X_i - X_{\lfloor t/2 \rfloor + i}\|_2\right\}\right) \\
 &= 1 - \mathbb{P}\left(\frac{|f(X_1) - f(X_2)|}{\|X_1 - X_2\|_2} \leq k_{i^*-1}\right)^{\lfloor t/2 \rfloor} \\
 &= 1 - (1 - \Gamma(f, k_{i^*-1}))^{\lfloor t/2 \rfloor}.
 \end{aligned}$$

It remains to show that  $\Gamma(f, k_{i^*-1}) > 0$ . Observe first that since  $f \in \text{Lip}(k_{i^*})$ , then the function  $F : (x, x') \mapsto |f(x) - f(x')| - k_{i^*-1} \cdot \|x - x'\|_2$  is also continuous. However, as  $f \notin \text{Lip}(k_{i^*-1})$ , we know that there exists some  $(x_1, x_2) \in \mathcal{X} \times \mathcal{X}$  such that  $F(x_1, x_2) > 0$ . Hence, it follows from the continuity of  $F$  that there necessarily exists some  $\epsilon > 0$  such that  $\forall (x, x') \in B(x_1, \epsilon) \cap \mathcal{X} \times B(x_2, \epsilon) \cap \mathcal{X}$ ,  $F(x, x') > 0$  which proves the proof.  $\square$

**Proof of Proposition 21.** Combining the consistency equivalence of Proposition 3 with the upper bound on the covering rate obtained in Example 4 gives the result.  $\square$

**Proof of Proposition 22.** Fix any  $\delta \in (0, 1)$ , set  $N_1 = 2 + \lceil 2 \ln(\delta/3) / \ln(1 - \Gamma(f, k_{i^*-1})) \rceil$  and  $N_2 = \lceil ((\sqrt{\ln(3/\delta)}/2 + 4N_1p - \sqrt{\ln(3/\delta)/2})/2p)^2 \rceil$ . Considering any  $n > N_2$ , we prove the result in three steps.

**Step 1.** As the constant  $N_1$  and  $N_2$  were chosen so that Hoeffding's inequality ensures that  $\mathbb{P}\left(\sum_{i=1}^{N_2} B_i \geq N_1\right) \geq 1 - \delta/3$ , we know that after  $N_2$  iterations and with probability  $1 - \delta/3$  we have collected at least  $N_1$  evaluation points randomly and uniformly distributed over  $\mathcal{X}$  due to the exploration step.

**Step 2.** Using Proposition 18 and the fist  $N_1$  evaluation points which have been sampled independently and uniformly over  $\mathcal{X}$ , we know that after  $N_2$  iterations and on the event  $\{\sum_{i=1}^{N_2} B_i \geq N_1\}$  the Lipschitz constant  $k_{i^*}$  has been estimated with probability at least  $1 - \delta/3$ , i.e.,  $\mathbb{P}\left(\forall t \geq N_2 + 1, \hat{k}_t = k_{i^*} \mid \sum_{i=1}^{N_2} B_i \geq N_1\right) \geq 1 - \delta/3$ .

**Step 3.** Finally, as the Lipschitz constant estimate  $\hat{k}_t$  satisfies  $f \in \text{Lip}(\hat{k}_t)$  for all  $t \geq N_2 + 1$  on the above event, one can easily show by reproducing the same steps as in Proposition 12 that conditioned upon the event  $\{\forall t \geq N_2 + 1, \hat{k}_t = k_{i^*}\} \cap \{\sum_{i=1}^{N_2} B_i \geq N_1\}$  the algorithm is always faster or equal to a Pure Random Search ran with  $n - N_2$  i.i.d. copies of  $X' \sim \mathcal{U}(\mathcal{X})$ . Therefore, using (i) the bound of Proposition 28, (ii) the elementaries inequalities  $\lfloor x \rfloor \leq x + 1$ ,  $\lfloor x \rfloor \geq x - 1$ ,  $\sqrt{x+y} - \sqrt{x} \leq \sqrt{y}$  and (iv) the definition of  $N_2 < n$ , we obtain that with probability at least  $(1 - \delta/3)^3 \geq 1 - \delta$ ,

$$\begin{aligned} \max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) &\leq k_{i^*} \cdot \text{diam}(\mathcal{X}) \cdot \left( \frac{\ln(3/\delta)}{n - N_2} \right)^{\frac{1}{d}} \\ &= k_{i^*} \cdot \text{diam}(\mathcal{X}) \cdot \left( \frac{n}{n - N_2} \right)^{\frac{1}{d}} \cdot \left( \frac{\ln(3/\delta)}{n} \right)^{\frac{1}{d}} \\ &\leq k_{i^*} \cdot \text{diam}(\mathcal{X}) \cdot (1 + N_2)^{\frac{1}{d}} \left( \frac{\ln(3/\delta)}{n} \right)^{\frac{1}{d}} \\ &\leq k_{i^*} \cdot \text{diam}(\mathcal{X}) \cdot \left( \frac{5}{p} + \frac{2 \ln(\delta/3)}{p \ln(1 - \Gamma(f, k_{i^*-1}))} \right)^{\frac{1}{d}} \cdot \left( \frac{\ln(3/\delta)}{n} \right)^{\frac{1}{d}}. \end{aligned}$$

The result is extended to the case where  $n \leq N_2$  by noticing that the bound is superior to  $k_{i^*} \cdot \text{diam}(\mathcal{X})$  in that case, and thus trivial.  $\square$

**Proof of Theorem 23.** Fix  $\delta \in (0, 1)$ , set  $N_1 = 2 + \lceil 2 \ln(4/\delta) / \ln(1 - \Gamma) \rceil$  and  $N_2 = \lceil ((\sqrt{\ln(4/\delta)}/2 + 4N_1p - \sqrt{\ln(4/\delta)/2})/2p)^2 \rceil$  and let  $N_3 = N_2 + \lceil 2 \ln(4/\delta) / (1 - p)^2 \rceil$ . Picking any  $n > N_3$ , we proceed similarly as in the proof of Proposition 22 in four steps:

**Steps 1 & 2.** As in the above prove, by definition of  $N_1$  and  $N_2$  and due to Hoeffding's inequality and Proposition 18, we know that the following event:  $\{\forall t \geq N_2 + 1, \hat{k}_t = k_{i^*}\} \cap \{\sum_{i=1}^{N_2} B_i \geq N_1\}$  holds true with probability at least  $(1 - \delta/4)^2$ .

**Step 3.** Again, using Hoeffding's inequality and the definition of  $N_2$  and  $N_3$ , we know that after the iteration  $N_2 + 1$  we have collected with probability at least  $1 - \delta/4$  at least

$(1-p)(n-N_3)/2$  exploitative evaluation points:

$$\sum_{i=N_2+1}^n \mathbb{I}\{B_i = 0\} \geq (1-p)(n-N_2) - \sqrt{\frac{(n-N_2)\ln(4/\delta)}{2}} \geq \frac{1-p}{2} \cdot (n-N_3).$$

**Step 4.** Reproducing the same steps as in the proof of the fast rate of Theorem 15 with the  $(1-p) \cdot (n-N_3)/2$  previous exploitative points and putting the previous results altogether gives that with probability at least  $(1-\delta/4)^4 \geq 1-\delta$ ,

$$\max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) \leq k_{i^*} \times \text{diam}(\mathcal{X}) \times$$

$$\begin{cases} \exp \left\{ -C_{k,\kappa} \cdot \frac{(1-p)(n-N_3)\ln(2)}{2\ln(4n/\delta) + 4(2\sqrt{d})^d} \right\}, & \kappa = 1 \\ \frac{2^\kappa}{2} \left( 1 + C_{k,\kappa} \cdot \frac{(1-p)(n-N_3)(2^{d(\kappa-1)} - 1)}{2\ln(4n/\delta) + 4(2\sqrt{d})^d} \right)^{-\frac{\kappa}{d(\kappa-1)}}, & \kappa > 1. \end{cases}$$

We now take out the term  $N_3$ . Since  $C_{k_{i^*},\kappa}(1-p) \leq 1$ , when  $\kappa = 1$ , we have

$$\max_{x \in \mathcal{X}} f(x) - \max_{i=1 \dots n} f(X_i) \leq k_{i^*} \cdot \text{diam}(\mathcal{X}) \cdot \exp(5N_3/2) \exp \left\{ -C_{k,\kappa} \cdot \frac{(1-p)n\ln(2)}{2\ln(4n/\delta) + 4(2\sqrt{d})^d} \right\}.$$

For  $\kappa > 1$ , setting  $C = C_{k_{i^*},\kappa}(1-p)/(2\ln(4n/\delta) + 4(2\sqrt{d})^d)$  and using the decomposition  $n = (n-N_3) + N_3$ , we bound the ratio:

$$\left( \frac{1 + Cn(2^{d(\kappa-1)} - 1)}{1 + C(n-N_3)(2^{d(\kappa-1)} - 1)} \right)^{\frac{\kappa}{d(\kappa-1)}} \leq \left( 1 + \frac{CN_3(2^{d(\kappa-1)} - 1)}{1 + C(2^{d(\kappa-1)} - 1)} \right)^{\frac{\kappa}{d(\kappa-1)}}.$$

In the case where  $\kappa/d(\kappa-1) \leq 1$ , one directly obtains

$$\left( 1 + \frac{CN_3(2^{d(\kappa-1)} - 1)}{1 + C(2^{d(\kappa-1)} - 1)} \right)^{\frac{\kappa}{d(\kappa-1)}} \leq (1 + N_3)^{\frac{\kappa}{d(\kappa-1)}} \leq (1 + N_3) \leq e^{N_3}$$

Considering the case where  $\kappa/d(\kappa-1) > 1$  and setting  $\kappa = 1 + \epsilon/d$  with  $\epsilon \in (0, 1)$ , we obtain from the inequalities (i)  $\kappa \leq 1 + 1/d \leq 2$ , (ii)  $\forall \epsilon \in (0, 1)$ ,  $2^\epsilon - 1 \leq \epsilon$  and (iii)  $C \leq 1/2$  that

$$\left( 1 + \frac{CN_3(2^{d(\kappa-1)} - 1)}{1 + C(2^{d(\kappa-1)} - 1)} \right)^{\frac{\kappa}{d(\kappa-1)}} \leq (1 + CN_3(2^\epsilon - 1))^{\frac{2}{\epsilon}} \leq (1 + CN_3\epsilon)^{\frac{2}{\epsilon}} \leq e^{2CN_3} \leq e^{N_3}$$

Finally, using standard bounds on  $N_3$  and noticing that the previous bound is superior to  $k_{i^*} \text{diam}(\mathcal{X})$  whenever  $n \leq N_3$ , the previous result remains valid for any  $n \in \mathbb{N}^*$ .  $\square$