

# Supplementary Material

“Robust Submodular Maximization: A Non-Uniform Partitioning Approach” (ICML 2017)

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## A. Proof of Proposition 4.1

We have

$$\begin{aligned}
 |S_0| &= \sum_{i=0}^{\lceil \log \tau \rceil} \lceil \tau/2^i \rceil 2^i \eta \\
 &\leq \sum_{i=0}^{\lceil \log \tau \rceil} \left( \frac{\tau}{2^i} + 1 \right) 2^i \eta \\
 &\leq \eta (\lceil \log \tau \rceil + 1) (\tau + 2^{\lceil \log \tau \rceil}) \\
 &\leq 3\eta \tau (\lceil \log \tau \rceil + 1) \\
 &\leq 3\eta \tau (\log k + 2).
 \end{aligned}$$

## B. Proof of Proposition 4.4

Recalling that  $\mathcal{A}_j(T)$  denotes a set constructed by the algorithm after  $j$  iterations, we have

$$\begin{aligned}
 f(\mathcal{A}_j(T)) - f(\mathcal{A}_{j-1}(T)) &\geq \frac{1}{\beta} \max_{e \in T} f(e|\mathcal{A}_{j-1}(T)) \\
 &\geq \frac{1}{\beta} \max_{e \in T} f(e|\mathcal{A}_k(T)) \\
 &\geq \frac{1}{\beta} \max_{e \in T \setminus \mathcal{A}_k(T)} f(e|\mathcal{A}_k(T)),
 \end{aligned} \tag{13}$$

where the first inequality follows from the  $\beta$ -iterative property (6), and the second inequality follows from  $\mathcal{A}_{j-1}(S) \subseteq \mathcal{A}_k(S)$  and the submodularity of  $f$ .

Continuing, we have

$$\begin{aligned}
 f(\mathcal{A}_k(T)) &= \sum_{j=1}^k f(\mathcal{A}_j(T)) - f(\mathcal{A}_{j-1}(T)) \\
 &\geq \frac{k}{\beta} \max_{e \in T \setminus \mathcal{A}_k(T)} f(e|\mathcal{A}_k(T)),
 \end{aligned}$$

where the last inequality follows from (13).

By rearranging, we have for any  $e \in T \setminus \mathcal{A}_k(T)$  that

$$f(e|\mathcal{A}_k(T)) \leq \beta \frac{f(\mathcal{A}_k(T))}{k}.$$

## C. Proof of Lemma 4.3

Recalling that  $A_j(T)$  denotes the set constructed after  $j$  iterations when applied to  $T$ , we have

$$\begin{aligned}
 \max_{e \in T \setminus A_{j-1}(T)} f(e|A_{j-1}(T)) &\geq \frac{1}{k} \sum_{e \in \text{OPT}(k, T) \setminus A_{j-1}(T)} f(e|A_{j-1}(T)) \\
 &\geq \frac{1}{k} f(\text{OPT}(k, T)|A_{j-1}(T)) \\
 &\geq \frac{1}{k} (f(\text{OPT}(k, T)) - f(A_{j-1}(T))),
 \end{aligned} \tag{14}$$

where the first line holds since the maximum is lower bounded by the average, the line uses submodularity, and the last line uses monotonicity.

By combining the  $\beta$ -iterative property with (14), we obtain

$$\begin{aligned} f(\mathcal{A}_j(T)) - f(\mathcal{A}_{j-1}(T)) &\geq \frac{1}{\beta} \max_{e \in T \setminus \mathcal{A}_{j-1}(T)} f(e | \mathcal{A}_{j-1}(T)) \\ &\geq \frac{1}{k\beta} (f(\text{OPT}(k, T)) - f(\mathcal{A}_{j-1}(T))). \end{aligned}$$

By rearranging, we obtain

$$f(\text{OPT}(k, T)) - f(\mathcal{A}_{j-1}(T)) \leq \beta k (f(\mathcal{A}_j(T)) - f(\mathcal{A}_{j-1}(T))). \quad (15)$$

We proceed by following the steps from the proof of Theorem 1.5 in (Krause & Golovin, 2012). Defining  $\delta_j := f(\text{OPT}(k, T)) - f(\mathcal{A}_j(T))$ , we can rewrite (15) as  $\delta_{j-1} \leq \beta k (\delta_{j-1} - \delta_j)$ . By rearranging, we obtain

$$\delta_j \leq \left(1 - \frac{1}{\beta k}\right) \delta_{j-1}.$$

Applying this recursively, we obtain  $\delta_l \leq \left(1 - \frac{1}{\beta k}\right)^l \delta_0$ , where  $\delta_0 = f(\text{OPT}(k, T))$  since  $f$  is normalized (i.e.,  $f(\emptyset) = 0$ ). Finally, applying  $1 - x \leq e^{-x}$  and rearranging, we obtain

$$f(\mathcal{A}_l(T)) \geq \left(1 - e^{-\frac{l}{\beta k}}\right) f(\text{OPT}(k, T)).$$

## D. Proof of Theorem 4.5

### D.1. Technical Lemmas

We first provide several technical lemmas that will be used throughout the proof. We begin with a simple property of submodular functions.

**Lemma D.1** *For any submodular function  $f$  on a ground set  $V$ , and any sets  $A, B, R \subseteq V$ , we have*

$$f(A \cup B) - f(A \cup (B \setminus R)) \leq f(R | A).$$

*Proof.* Define  $R_2 := A \cap R$ , and  $R_1 := R \setminus A = R \setminus R_2$ . We have

$$\begin{aligned} f(A \cup B) - f(A \cup (B \setminus R)) &= f(A \cup B) - f((A \cup B) \setminus R_1) \\ &= f(R_1 | (A \cup B) \setminus R_1) \\ &\leq f(R_1 | (A \setminus R_1)) \end{aligned} \quad (16)$$

$$= f(R_1 | A) \quad (17)$$

$$= f(R_1 \cup R_2 | A) \quad (18)$$

$$= f(R | A),$$

where (16) follows from the submodularity of  $f$ , (17) follows since  $A$  and  $R_1$  are disjoint, and (18) follows since  $R_2 \subseteq A$ .  $\square$

The next lemma provides a simple lower bound on the maximum of two quantities; it is stated formally since it will be used on multiple occasions.

**Lemma D.2** *For any set function  $f$ , sets  $A, B$ , and constant  $\alpha > 0$ , we have*

$$\max\{f(A), f(B) - \alpha f(A)\} \geq \left(\frac{1}{1 + \alpha}\right) f(B), \quad (19)$$

and

$$\max\{\alpha f(A), f(B) - f(A)\} \geq \left(\frac{\alpha}{1 + \alpha}\right) f(B). \quad (20)$$

*Proof.* Starting with (19), we observe that one term is increasing in  $f(A)$  and the other is decreasing in  $f(A)$ . Hence, the maximum over all possible  $f(A)$  is achieved when the two terms are equal, i.e.,  $f(A) = \frac{1}{1+\alpha}f(B)$ . We obtain (20) via the same argument.  $\square$

The following lemma relates the function values associated with two buckets formed by PRO, denoted by  $X$  and  $Y$ . It is stated with respect to an arbitrary set  $E_Y$ , but when we apply the lemma, this will correspond to the elements of  $Y$  that are removed by the adversary.

**Lemma D.3** *Under the setup of Theorem 4.5, let  $X$  and  $Y$  be buckets of PRO such that  $Y$  is constructed at a later time than  $X$ . For any set  $E_Y \subseteq Y$ , we have*

$$f(X \cup (Y \setminus E_Y)) \geq \frac{1}{1+\alpha}f(Y),$$

and

$$f(E_Y | X) \leq \alpha f(X), \tag{21}$$

where  $\alpha = \beta \frac{|E_Y|}{|X|}$ .

*Proof.* Inequality (21) follows from the  $\beta$ -iterative property of  $\mathcal{A}$ ; specifically, we have from (8) that

$$f(e|X) \leq \beta \frac{f(X)}{|X|},$$

where  $e$  is any element of the ground set that is neither in  $X$  nor any bucket constructed before  $X$ . Hence, we can write

$$f(E_Y | X) \leq \sum_{e \in E_Y} f(e|X) \leq \beta \frac{|E_Y|}{|X|} f(X) = \alpha f(X),$$

where the first inequality is by submodularity. This proves (21).

Next, we write

$$f(Y) - f(X \cup (Y \setminus E_Y)) \leq f(X \cup Y) - f(X \cup (Y \setminus E_Y)) \tag{22}$$

$$\leq f(E_Y | X), \tag{23}$$

where (22) is by monotonicity, and (23) is by Lemma D.1 with  $A = X$ ,  $B = Y$ , and  $R = E_Y$ .

Combining (21) and (23), together with the fact that  $f(X \cup (Y \setminus E_Y)) \geq f(X)$  (by monotonicity), we have

$$\begin{aligned} f(X \cup (Y \setminus E_Y)) &\geq \max \{f(X), f(Y) - \alpha f(X)\} \\ &\geq \frac{1}{1+\alpha}f(Y), \end{aligned} \tag{24}$$

where (24) follows from (19).  $\square$

Finally, we provide a lemma that will later be used to take two bounds that are known regarding the previously-constructed buckets, and use them to infer bounds regarding the next bucket.

**Lemma D.4** *Under the setup of Theorem 4.5, let  $Y$  and  $Z$  be buckets of PRO such that  $Z$  is constructed at a later time than  $Y$ , and let  $E_Y \subseteq Y$  and  $E_Z \subseteq Z$  be arbitrary sets. Moreover, let  $X$  be a set (not necessarily a bucket) such that*

$$f((Y \setminus E_Y) \cup X) \geq \frac{1}{1+\alpha}f(Y), \tag{25}$$

and

$$f(E_Y | X) \leq \alpha f(X). \tag{26}$$

Then, we have

$$f(E_Z | (Y \setminus E_Y) \cup X) \leq \alpha_{\text{next}} f((Y \setminus E_Y) \cup X), \tag{27}$$

and

$$f((Z \setminus E_Z) \cup (Y \setminus E_Y) \cup X) \geq \frac{1}{1 + \alpha_{\text{next}}} f(Z), \quad (28)$$

where

$$\alpha_{\text{next}} = \beta \frac{|E_Z|}{|Y|} (1 + \alpha) + \alpha. \quad (29)$$

*Proof.* We first prove (27):

$$\begin{aligned} f(E_Z | (Y \setminus E_Y) \cup X) &= f((Y \setminus E_Y) \cup X \cup E_Z) - f((Y \setminus E_Y) \cup X) \\ &\leq f(X \cup Y \cup E_Z) - f((Y \setminus E_Y) \cup X) \end{aligned} \quad (30)$$

$$\begin{aligned} &= f(E_Z | X \cup Y) + f(X \cup Y) - f((Y \setminus E_Y) \cup X) \\ &\leq f(E_Z | Y) + f(X \cup Y) - f((Y \setminus E_Y) \cup X) \end{aligned} \quad (31)$$

$$\leq \beta \frac{|E_Z|}{|Y|} f(Y) + f(X \cup Y) - f((Y \setminus E_Y) \cup X) \quad (32)$$

$$\leq \beta \frac{|E_Z|}{|Y|} (1 + \alpha) f((Y \setminus E_Y) \cup X) + f(X \cup Y) - f((Y \setminus E_Y) \cup X) \quad (33)$$

$$\leq \beta \frac{|E_Z|}{|Y|} (1 + \alpha) f((Y \setminus E_Y) \cup X) + f(E_Y | (Y \setminus E_Y) \cup X) \quad (34)$$

$$\leq \beta \frac{|E_Z|}{|Y|} (1 + \alpha) f((Y \setminus E_Y) \cup X) + f(E_Y | X) \quad (35)$$

$$\leq \beta \frac{|E_Z|}{|Y|} (1 + \alpha) f((Y \setminus E_Y) \cup X) + \alpha f(X) \quad (36)$$

$$\leq \beta \frac{|E_Z|}{|Y|} (1 + \alpha) f((Y \setminus E_Y) \cup X) + \alpha f((Y \setminus E_Y) \cup X) \quad (37)$$

$$= \left( \beta \frac{|E_Z|}{|Y|} (1 + \alpha) + \alpha \right) f((Y \setminus E_Y) \cup X), \quad (38)$$

where: (30) and (31) follow by monotonicity and submodularity, respectively; (32) follows from the second part of Lemma D.3; (33) follows from (25); (34) is obtained by applying Lemma D.1 for  $A = X$ ,  $B = Y$ , and  $R = E_Y$ ; (35) follows by submodularity; (36) follows from (26); (37) follows by monotonicity. Finally, by defining  $\alpha_{\text{next}} := \beta \frac{|E_Z|}{|Y|} (1 + \alpha) + \alpha$  in (38) we establish the bound in (27).

In the rest of the proof, we show that (28) holds as well. First, we have

$$f((Z \setminus E_Z) \cup (Y \setminus E_Y) \cup X) \geq f(Z) - f(E_Z | (Y \setminus E_Y) \cup X) \quad (39)$$

by Lemma D.1 with  $B = Z$ ,  $R = E_Z$  and  $A = (Y \setminus E_Y) \cup X$ . Now we can use the derived bounds (38) and (39) to obtain

$$\begin{aligned} f((Z \setminus E_Z) \cup (Y \setminus E_Y) \cup X) &\geq f(Z) - f(E_Z | (Y \setminus E_Y) \cup X) \\ &\geq f(Z) - \left( \beta \frac{|E_Z|}{|Y|} (1 + \alpha) + \alpha \right) f((Y \setminus E_Y) \cup X). \end{aligned}$$

Finally, we have

$$\begin{aligned} f((Z \setminus E_Z) \cup (Y \setminus E_Y) \cup X) &\geq \max \left\{ f((Y \setminus E_Y) \cup X), f(Z) - \left( \beta \frac{|E_Z|}{|Y|} (1 + \alpha) + \alpha \right) f((Y \setminus E_Y) \cup X) \right\} \\ &\geq \frac{1}{1 + \alpha_{\text{next}}} f(Z), \end{aligned}$$

where the last inequality follows from Lemma D.1.  $\square$

Observe that the results we obtain on  $f(E_Z | (Y \setminus E_Y) \cup X)$  and on  $f((Z \setminus E_Z) \cup (Y \setminus E_Y) \cup X)$  in Lemma D.4 are of the same form as the pre-conditions of the lemma. This will allow us to apply the lemma recursively.

## D.2. Characterizing the Adversary

Let  $E$  denote a set of elements removed by an adversary, where  $|E| \leq \tau$ . Within  $S_0$ , PRO constructs  $\lceil \log \tau \rceil + 1$  partitions. Each partition  $i \in \{0, \dots, \lceil \log \tau \rceil\}$  consists of  $\lceil \tau/2^i \rceil$  buckets, each of size  $2^i \eta$ , where  $\eta \in \mathbb{N}$  will be specified later. We let  $B$  denote a generic bucket, and define  $E_B$  to be all the elements removed from this bucket, i.e.  $E_B = B \cap E$ .

The following lemma identifies a bucket in each partition for which not too many elements are removed.

**Lemma D.5** *Under the setup of Theorem 4.5, suppose that an adversary removes a set  $E$  of size at most  $\tau$  from the set  $S$  constructed by PRO. Then for each partition  $i$ , there exists a bucket  $B_i$  such that  $|E_{B_i}| \leq 2^i$ , i.e., at most  $2^i$  elements are removed from this bucket.*

*Proof.* Towards contradiction, assume that this is not the case, i.e., assume  $|E_{B_i}| > 2^i$  for every bucket of the  $i$ -th partition. As the number of buckets in partition  $i$  is  $\lceil \tau/2^i \rceil$ , this implies that the adversary has to spend a budget of

$$|E| \geq 2^i |E_{B_i}| > 2^i \lceil \tau/2^i \rceil = \tau,$$

which is in contradiction with  $|E| \leq \tau$ .  $\square$

We consider  $B_0, \dots, B_{\lceil \log \tau \rceil}$  as above, and show that even in the worst case,  $f\left(\bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i})\right)$  is almost as large as  $f(B_{\lceil \log \tau \rceil})$  for appropriately set  $\eta$ . To achieve this, we apply Lemma D.4 multiple times, as illustrated in the following lemma. We henceforth write  $\eta_h := \eta/2$  for brevity.

**Lemma D.6** *Under the setup of Theorem 4.5, suppose that an adversary removes a set  $E$  of size at most  $\tau$  from the set  $S$  constructed by PRO, and let  $B_0, \dots, B_{\lceil \log \tau \rceil}$  be buckets such that  $|E_{B_i}| \leq 2^i$  for each  $i \in \{1, \dots, \lceil \log \tau \rceil\}$  (cf., Lemma D.5). Then,*

$$f\left(\bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i})\right) \geq \left(1 - \frac{1}{1 + \frac{1}{\alpha}}\right) f(B_{\lceil \log \tau \rceil}) = \frac{1}{1 + \alpha} f(B_{\lceil \log \tau \rceil}), \quad (40)$$

and

$$f\left(E_{B_{\lceil \log \tau \rceil}} \mid \bigcup_{i=0}^{\lceil \log \tau \rceil - 1} (B_i \setminus E_{B_i})\right) \leq \alpha f\left(\bigcup_{i=0}^{\lceil \log \tau \rceil - 1} (B_i \setminus E_{B_i})\right), \quad (41)$$

for some

$$\alpha \leq \beta^2 \frac{(1 + \eta_h)^{\lceil \log \tau \rceil} - \eta_h^{\lceil \log \tau \rceil}}{\eta_h^{\lceil \log \tau \rceil}}. \quad (42)$$

*Proof.* In what follows, we focus on the case where there exists some bucket  $B_0$  in partition  $i = 0$  such that  $B_0 \setminus E_{B_0} = B_0$ . If this is not true, then  $E$  must be contained entirely within this partition, since it contains  $\tau$  buckets. As a result, (i) we trivially obtain (40) even when  $\alpha$  is replaced by zero, since the union on the left-hand side contains  $B_{\lceil \log \tau \rceil}$ ; (ii) (41) becomes trivial since the left-hand side is zero is a result of  $E_{B_{\lceil \log \tau \rceil}} = \emptyset$ .

We proceed by induction. Namely, we show that

$$f\left(\bigcup_{i=0}^j (B_i \setminus E_{B_i})\right) \geq \left(1 - \frac{1}{1 + \frac{1}{\alpha_j}}\right) f(B_j) = \frac{1}{1 + \alpha_j} f(B_j), \quad (43)$$

and

$$f\left(E_{B_j} \mid \bigcup_{i=0}^{j-1} (B_i \setminus E_{B_i})\right) \leq \alpha_j f\left(\bigcup_{i=0}^{j-1} (B_i \setminus E_{B_i})\right), \quad (44)$$

for every  $j \geq 1$ , where

$$\alpha_j \leq \beta^2 \frac{(1 + \eta_h)^j - \eta_h^j}{\eta_h^j}. \quad (45)$$

Upon showing this, the lemma is concluded by setting  $j = \lceil \log \tau \rceil$ .

**Base case**  $j = 1$ . In the case that  $j = 1$ , taking into account that  $E_{B_0} = \emptyset$ , we observe from (43) that our goal is to bound  $f(B_0 \cup (B_1 \setminus E_{B_1}))$ . Applying Lemma D.3 with  $X = B_0$ ,  $Y = B_1$ , and  $E_Y = E_{B_1}$ , we obtain

$$f(B_0 \cup (B_1 \setminus E_{B_1})) \geq \frac{1}{1 + \alpha_1} f(B_1),$$

and

$$f(E_{B_1} | B_0) \leq \alpha_1 f(B_0),$$

where  $\alpha_1 = \beta \frac{|E_{B_1}|}{|B_0|}$ . We have  $|B_0| = \eta$ , while  $|E_{B_1}| \leq 2$  by assumption. Hence, we can upper bound  $\alpha_1$  and rewrite as

$$\alpha_1 \leq \beta \frac{2}{\eta} = \beta \frac{1}{\eta_h} = \beta \frac{(1 + \eta_h) - \eta_h}{\eta_h} \leq \beta^2 \frac{(1 + \eta_h) - \eta_h}{\eta_h},$$

where the last inequality follows since  $\beta \geq 1$  by definition.

**Inductive step.** Fix  $j \geq 2$ . Assuming that the inductive hypothesis holds for  $j - 1$ , we want to show that it holds for  $j$  as well.

We write

$$f\left(\bigcup_{i=0}^j (B_i \setminus E_{B_i})\right) = f\left(\left(\bigcup_{i=0}^{j-1} (B_i \setminus E_{B_i})\right) \cup (B_j \setminus E_{B_j})\right),$$

and apply Lemma D.4 with  $X = \bigcup_{i=0}^{j-1} (B_i \setminus E_{B_i})$ ,  $Y = B_{j-1}$ ,  $E_Y = E_{B_{j-1}}$ ,  $Z = B_j$ , and  $E_Z = E_{B_j}$ . Note that the conditions (25) and (26) of Lemma D.4 are satisfied by the inductive hypothesis. Hence, we conclude that (43) and (44) hold with

$$\alpha_j = \beta \frac{|E_{B_j}|}{|B_{j-1}|} (1 + \alpha_{j-1}) + \alpha_{j-1}.$$

It remains to show that (45) holds for  $\alpha_j$ , assuming it holds for  $\alpha_{j-1}$ . We have

$$\begin{aligned} \alpha_j &= \beta \frac{|E_{B_j}|}{|B_{j-1}|} (1 + \alpha_{j-1}) + \alpha_{j-1} \\ &\leq \beta \frac{1}{\eta_h} \left(1 + \beta \frac{(1 + \eta_h)^{j-1} - \eta_h^{j-1}}{\eta_h^{j-1}}\right) + \beta \frac{(1 + \eta_h)^{j-1} - \eta_h^{j-1}}{\eta_h^{j-1}} \end{aligned} \quad (46)$$

$$\leq \beta^2 \left(\frac{1}{\eta_h} \left(1 + \frac{(1 + \eta_h)^{j-1} - \eta_h^{j-1}}{\eta_h^{j-1}}\right) + \frac{(1 + \eta_h)^{j-1} - \eta_h^{j-1}}{\eta_h^{j-1}}\right) \quad (47)$$

$$= \beta^2 \left(\frac{1}{\eta_h} \frac{(1 + \eta_h)^{j-1}}{\eta_h^{j-1}} + \frac{(1 + \eta_h)^{j-1} - \eta_h^{j-1}}{\eta_h^{j-1}}\right)$$

$$= \beta^2 \left(\frac{(1 + \eta_h)^{j-1}}{\eta_h^j} + \frac{\eta_h(1 + \eta_h)^{j-1} - \eta_h^j}{\eta_h^j}\right)$$

$$= \beta^2 \frac{(1 + \eta_h)^j - \eta_h^j}{\eta_h^j},$$

where (46) follows from (45) and the fact that

$$\beta \frac{|E_{B_j}|}{|B_{j-1}|} \leq \beta \frac{2^j}{2^{j-1}\eta} = \beta \frac{2}{\eta} = \beta \frac{1}{\eta_h},$$

by  $|E_{B_j}| \leq 2^j$  and  $|B_{j-1}| = 2^{j-1}\eta$ ; and (47) follows since  $\beta \geq 1$ .  $\square$

Inequality (45) provides an upper bound on  $\alpha_j$ , but it is not immediately clear how the bound varies with  $j$ . The following lemma provides a more compact form.

**Lemma D.7** Under the setup of Lemma D.6, we have for  $2^{\lceil \log \tau \rceil} \leq \eta_h$  that

$$\alpha_j \leq 3\beta^2 \frac{j}{\eta} \quad (48)$$

*Proof.* We unfold the right-hand side of (45) in order to express it in a simpler way. First, consider  $j = 1$ . From (45) we obtain  $\alpha_1 \leq 2\beta^2 \frac{1}{\eta}$ , as required. For  $j \geq 2$ , we obtain the following:

$$\begin{aligned} \beta^{-2}\alpha_j &\leq \frac{(1 + \eta_h)^j - \eta_h^j}{\eta_h^j} \\ &= \sum_{i=0}^{j-1} \binom{j}{i} \frac{\eta_h^i}{\eta_h^j} \end{aligned} \quad (49)$$

$$= \frac{j}{\eta_h} + \sum_{i=0}^{j-2} \binom{j}{i} \frac{\eta_h^i}{\eta_h^j} \quad (50)$$

$$\begin{aligned} &= \frac{j}{\eta_h} + \sum_{i=0}^{j-2} \left( \frac{\prod_{t=1}^{j-i} (j-t+1)}{\prod_{t=1}^{j-i} t} \frac{\eta_h^i}{\eta_h^j} \right) \\ &\leq \frac{j}{\eta_h} + \frac{1}{2} \sum_{i=0}^{j-2} j^{j-i} \frac{\eta_h^i}{\eta_h^j} \end{aligned} \quad (51)$$

$$\begin{aligned} &= \frac{j}{\eta_h} + \frac{1}{2} \sum_{i=0}^{j-2} \left( \frac{j}{\eta_h} \right)^{j-i} \\ &= \frac{j}{\eta_h} + \frac{1}{2} \left( -1 - \frac{j}{\eta_h} + \sum_{i=0}^j \left( \frac{j}{\eta_h} \right)^{j-i} \right), \end{aligned}$$

where (49) is a standard summation identity, and (51) follows from  $\prod_{t=1}^{j-i} (j-t+1) \leq j^{j-i}$  and  $\prod_{t=1}^{j-i} t \geq 2$  for  $j-i \geq 2$ . Next, explicitly evaluating the summation of the last equality, we obtain

$$\begin{aligned} \beta^{-2}\alpha_j &\leq \frac{j}{\eta_h} + \frac{1}{2} \left( -1 - \frac{j}{\eta_h} + \frac{1 - \left( \frac{j}{\eta_h} \right)^{j+1}}{1 - \frac{j}{\eta_h}} \right) \\ &\leq \frac{j}{\eta_h} + \frac{1}{2} \left( -1 - \frac{j}{\eta_h} + \frac{1}{1 - \frac{j}{\eta_h}} \right) \\ &= \frac{j}{\eta_h} + \frac{1}{2} \left( \frac{\left( \frac{j}{\eta_h} \right)^2}{1 - \frac{j}{\eta_h}} \right) \end{aligned} \quad (52)$$

$$= \frac{j}{\eta_h} + \frac{j}{2\eta_h} \left( \frac{\frac{j}{\eta_h}}{1 - \frac{j}{\eta_h}} \right), \quad (53)$$

where (52) follows from  $(-a-1)(-a+1) = a^2 - 1$  with  $a = j/\eta_h$ .

Next, observe that if  $j/\eta_h \leq 1/2$ , or equivalently

$$2j \leq \eta_h, \quad (54)$$

then we can weaken (53) to

$$\beta^{-2}\alpha_j \leq \frac{j}{\eta_h} + \frac{j}{2\eta_h} = \frac{3}{2} \frac{j}{\eta_h} = 3 \frac{j}{\eta}, \quad (55)$$

which yields (48).

□

### D.3. Completing the Proof of Theorem 4.5

We now prove Theorem 4.5 in several steps. Throughout, we define  $\mu$  to be a constant such that  $f(E_1 | (S \setminus E)) = \mu f(S_1)$  holds, and we write  $E_0 := E_S^* \cap S_0$ ,  $E_1 := E_S^* \cap S_1$ , and  $E_{B_i} := E_S^* \cap B_i$ , where  $E_S^*$  is defined in (9). We also make use of the following lemma characterizing the optimal adversary. The proof is straightforward, and can be found in Lemma 2 of (Orlin et al., 2016)

**Lemma D.8** (Orlin et al., 2016) *Under the setup of Theorem 4.5, we have for all  $X \subset V$  with  $|X| \leq \tau$  that*

$$f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*) \leq f(\text{OPT}(k - \tau, V \setminus X)).$$

**Initial lower bounds:** We start by providing three lower bounds on  $f(S \setminus E_S^*)$ . First, we observe that  $f(S \setminus E_S^*) \geq f(S_0 \setminus E_0)$  and  $f(S \setminus E_S^*) \geq f\left(\bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i})\right)$ . We also have

$$\begin{aligned} f(S \setminus E) &= f(S) - f(S) + f(S \setminus E) \\ &= f(S_0 \cup S_1) + f(S \setminus E_0) - f(S \setminus E_0) - f(S) + f(S \setminus E) \end{aligned} \quad (56)$$

$$\begin{aligned} &= f(S_1) + f(S_0 | S_1) + f(S \setminus E_0) - f(S) - f(S \setminus E_0) + f(S \setminus E) \\ &= f(S_1) + f(S_0 | (S \setminus S_0)) + f(S \setminus E_0) - f(E_0 \cup (S \setminus E_0)) - f(S \setminus E_0) + f(S \setminus E) \end{aligned} \quad (57)$$

$$\begin{aligned} &= f(S_1) + f(S_0 | (S \setminus S_0)) - f(E_0 | (S \setminus E_0)) - f(S \setminus E_0) + f(S \setminus E) \\ &= f(S_1) + f(S_0 | (S \setminus S_0)) - f(E_0 | (S \setminus E_0)) - f(E_1 \cup (S \setminus E)) + f(S \setminus E) \end{aligned} \quad (58)$$

$$\begin{aligned} &= f(S_1) + f(S_0 | (S \setminus S_0)) - f(E_0 | (S \setminus E_0)) - f(E_1 | S \setminus E) \\ &= f(S_1) - f(E_1 | S \setminus E) + f(S_0 | (S \setminus S_0)) - f(E_0 | (S \setminus E_0)) \\ &\geq (1 - \mu)f(S_1), \end{aligned} \quad (59)$$

where (56) and (57) follow from  $S = S_0 \cup S_1$ , (58) follows from  $E_S^* = E_0 \cup E_1$ , and (59) follows from  $f(S_0 | (S \setminus S_0)) - f(E_0 | (S \setminus E_0)) \geq 0$  (due to  $E_0 \subseteq S_0$  and  $S \setminus S_0 \subseteq S \setminus E_0$ ), along with the definition of  $\mu$ .

By combining the above three bounds on  $f(S \setminus E_S^*)$ , we obtain

$$f(S \setminus E_S^*) \geq \max \left\{ f(S_0 \setminus E_0), (1 - \mu)f(S_1), f\left(\bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i})\right) \right\}. \quad (60)$$

We proceed by further bounding these terms.

**Bounding the first term in (60):** Defining  $S'_0 := \text{OPT}(k - \tau, V \setminus E_0) \cap (S_0 \setminus E_0)$  and  $X := \text{OPT}(k - \tau, V \setminus E_0) \setminus S'_0$ , we have

$$f(S_0 \setminus E_0) + f(\text{OPT}(k - \tau, V \setminus S_0)) \geq f(S'_0) + f(X) \quad (61)$$

$$\geq f(\text{OPT}(k - \tau, V \setminus E_0)) \quad (62)$$

$$\geq f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*), \quad (63)$$

where (61) follows from monotonicity, i.e.  $(S_0 \setminus E_0) \subseteq S'_0$  and  $(V \setminus S_0) \subseteq (V \setminus E_0)$ , (62) follows from the fact that  $\text{OPT}(k - \tau, V \setminus E_0) = S'_0 \cup X$  and submodularity,<sup>2</sup> and (63) follows from Lemma D.8 and  $|E_0| \leq \tau$ . We rewrite (63) as

$$f(S_0 \setminus E_0) \geq f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*) - f(\text{OPT}(k - \tau, V \setminus S_0)). \quad (64)$$

**Bounding the second term in (60):** Note that  $S_1$  is obtained by using  $\mathcal{A}$  that satisfies the  $\beta$ -iterative property on the set  $V \setminus S_0$ , and its size is  $|S_1| = k - |S_0|$ . Hence, from Lemma 4.3 with  $k - \tau$  in place of  $k$ , we have

$$f(S_1) \geq \left(1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}}\right) f(\text{OPT}(k - \tau, V \setminus S_0)). \quad (65)$$

<sup>2</sup>The submodularity property can equivalently be written as  $f(A) + f(B) \geq f(A \cup B) + f(A \cap B)$ .



**Bounding the third term in (60):** We can view  $S_1$  as a large bucket created by our algorithm after creating the buckets in  $S_0$ . Therefore, we can apply Lemma D.4 with  $X = \bigcup_{i=0}^{\lceil \log \tau \rceil - 1} (B_i \setminus E_{B_i})$ ,  $Y = B_{\lceil \log \tau \rceil}$ ,  $Z = S_1$ ,  $E_Y = E_S^* \cap Y$ , and  $E_Z = E_1$ . Conditions (25) and (26) needed to apply Lemma D.4 are provided by Lemma D.6. From Lemma D.4, we obtain the following with  $\alpha$  as in (42):

$$f \left( E_1 \mid \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) \cup (S_1 \setminus E_1) \right) \leq \left( \beta \frac{|E_1|}{|B_{\lceil \log \tau \rceil}|} (1 + \alpha) + \alpha \right) f \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right). \quad (66)$$

Furthermore, noting that the assumption  $\eta \geq 4(\log k + 1)$  implies  $2\lceil \log \tau \rceil \leq \eta h$ , we can upper-bound  $\alpha$  as in Lemma D.7 by (48) for  $j = \lceil \log \tau \rceil$ . Also, we have  $\beta \frac{|E_1|}{|B_{\lceil \log \tau \rceil}|} \leq \beta \frac{\tau}{2^{\lceil \log \tau \rceil} \eta} \leq \frac{\beta}{\eta}$ . Putting these together, we upper bound (66) as follows:

$$\begin{aligned} f \left( E_1 \mid \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) \cup (S_1 \setminus E_1) \right) &\leq \left( \frac{\beta}{\eta} \left( 1 + \frac{3\beta^2 \lceil \log \tau \rceil}{\eta} \right) + \frac{3\beta^2 \lceil \log \tau \rceil}{\eta} \right) f \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) \\ &\leq \frac{5\beta^3 \lceil \log \tau \rceil}{\eta} f \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right), \end{aligned}$$

where we have used  $\eta \geq 1$  and  $\lceil \log \tau \rceil \geq 1$  (since  $\tau \geq 2$  by assumption). We rewrite the previous equation as

$$\begin{aligned} f \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) &\geq \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} f \left( E_1 \mid \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) \cup (S_1 \setminus E_1) \right) \\ &\geq \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} f(E_1 \mid (S \setminus E)) \end{aligned} \quad (67)$$

$$= \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} \mu f(S_1), \quad (68)$$

where (67) follows from submodularity, and (68) follows from the definition of  $\mu$ .

**Combining the bounds:** Returning to (60), we have

$$\begin{aligned} f(S \setminus E_S^*) &\geq \max \left\{ f(S_0 \setminus E_0), (1 - \mu) f(S_1), f \left( \bigcup_{i=0}^{\lceil \log \tau \rceil} (B_i \setminus E_{B_i}) \right) \right\} \\ &\geq \max \left\{ f(S_0 \setminus E_0), (1 - \mu) f(S_1), \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} \mu f(S_1) \right\} \end{aligned} \quad (69)$$

$$\begin{aligned} &\geq \max \{ f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*) - f(\text{OPT}(k - \tau, V \setminus S_0)), \\ &\quad (1 - \mu) \left( 1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}} \right) f(\text{OPT}(k - \tau, V \setminus S_0)), \\ &\quad \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} \mu \left( 1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}} \right) f(\text{OPT}(k - \tau, V \setminus S_0)) \} \end{aligned} \quad (70)$$

$$\begin{aligned} &\geq \max \{ f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*) - f(\text{OPT}(k - \tau, V \setminus S_0)), \\ &\quad \frac{\eta}{5\beta^3 \lceil \log \tau \rceil} \left( 1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}} \right) f(\text{OPT}(k - \tau, V \setminus S_0)) \} \end{aligned} \quad (71)$$

$$\begin{aligned} &= \max \{ f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*) - f(\text{OPT}(k - \tau, V \setminus S_0)), \\ &\quad \frac{\eta}{5\beta^3 \lceil \log \tau \rceil + \eta} \left( 1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}} \right) f(\text{OPT}(k - \tau, V \setminus S_0)) \} \\ &\geq \frac{\eta}{5\beta^3 \lceil \log \tau \rceil + \eta} \left( 1 - e^{-\frac{k - |S_0|}{\beta(k - \tau)}} \right) f(\text{OPT}(k, V, \tau) \setminus E_{\text{OPT}(k, V, \tau)}^*), \end{aligned} \quad (72)$$

where (69) follows from (68), (70) follows from (64) and (65), (71) follows since  $\max\{1 - \mu, c\mu\} \geq \frac{c}{1+c}$  analogously to (19), and (72) follows from (20). Hence, we have established (72).

Turning to the permitted values of  $\tau$ , we have from Proposition 4.1 that

$$|S_0| \leq 3\eta\tau(\log k + 2).$$

For the choice of  $\tau$  to yield valid set sizes, we only require  $|S_0| \leq k$ ; hence, it suffices that

$$\tau \leq \frac{k}{3\eta(\log k + 2)}. \tag{73}$$

Finally, we consider the second claim of the lemma. For  $\tau \in o\left(\frac{k}{\eta(\log k)}\right)$  we have  $|S_0| \in o(k)$ . Furthermore, by setting  $\eta \geq \log^2 k$  (which satisfies the assumption  $\eta \geq 4(\log k + 1)$  for large  $k$ ), we get  $\frac{k-|S_0|}{\beta(k-\tau)} \rightarrow \beta^{-1}$  and  $\frac{\eta}{5\beta^3\lceil\log \tau\rceil + \eta} \rightarrow 1$  as  $k \rightarrow \infty$ . Hence, the constant factor converges to  $\frac{1-e^{-1/\beta}}{2-e^{-1/\beta}}$ , yielding (11). In the case that GREEDY is used as the subroutine, we have  $\beta = 1$ , and hence the constant factor converges to  $\frac{1-e^{-1}}{2-e^{-1}} \geq 0.387$ . If THRESHOLDING-GREEDY is used, we have  $\beta = \frac{1}{1-\epsilon}$ , and hence the constant factor converges to  $\frac{1-e^{\epsilon-1}}{2-e^{\epsilon-1}} \geq (1-\epsilon)\frac{1-e^{-1}}{2-e^{-1}} \geq (1-\epsilon)0.387$ .