

A. Organization of the Appendix

Appendix B describes the Greedy algorithm to solve Noisy $\mathbf{M}k\mathbf{SC}$ and its approximation ratio.

Appendix C provides omitted proofs from Section 3.

Appendix D provides omitted proofs from Section 4.

Appendix E presents details of the sampling method for Influence Maximization with k topics.

B. Greedy Algorithm

In this part, we investigate performance guarantee of Greedy algorithm. Greedy has been proven to obtain an approximation ratio of 2 for $\mathbf{M}k\mathbf{SC}$ with monotone non-noisy objective function (Ohsaka & Yoshida, 2015). We extend the authors' proof to $\mathbf{M}k\mathbf{SC}$ under noise and show that Greedy is able to obtain approximation of $\frac{2+2\epsilon B}{1-\epsilon}$ when f is monotone and $\frac{(2+2\epsilon+4\epsilon B)(1+\epsilon)}{(1-\epsilon)^2} + 1$ in case of non-monotonicity.

The pseudo code of Greedy is presented by Alg. 4

Algorithm 4 Greedy Algorithm

Input F, k, B

- 1: $\mathbf{s}^0 = \{\emptyset, \emptyset, \dots, \emptyset\}$
- 2: **for** $t = 1 \rightarrow B$ **do**
- 3: $e, i = \operatorname{argmax}_{e \in V; i \in [k]} F(\mathbf{s}^{t-1} \sqcup \langle e, i \rangle)$
- 4: $\mathbf{s}^t = \mathbf{s}^{t-1} \sqcup \langle e, i \rangle$

Return \mathbf{s}^B if f is **monotone**; $\operatorname{argmax}_{\mathbf{s}^t; t \in \{1, \dots, B\}} F(\mathbf{s}^t)$ if f is **non-monotone**.

Theorem 5. Given an instance of Noisy $\mathbf{M}k\mathbf{SC}$ with input V, k, B and F is an ϵ -estimation of the **monotone** k -submodular objective function f . If \mathbf{s} is an output of the Greedy algorithm and \mathbf{o} is an optimal solution, then

$$f(\mathbf{o}) \leq \frac{2 + 2\epsilon B}{1 - \epsilon} f(\mathbf{s}) \quad (3)$$

Proof. Let e^j and i^j be a selection in iteration j of the Greedy algorithm, we construct a sequence $\{\mathbf{o}^j\}$ as follows:

- $\mathbf{o}^0 = \mathbf{o}$
- With $j > 0$, let $S^j = \operatorname{supp}(\mathbf{o}^{j-1}) \setminus \operatorname{supp}(\mathbf{s}^{j-1})$. Let $\mathbf{o}^j = e^j$ if $e^j \in S^j$, otherwise let \mathbf{o}^j to be an arbitrary element in S^j .
 - Define $\mathbf{o}^{j-1/2}$ as a k -set that $\mathbf{o}^{j-1/2}(e) = \mathbf{o}^{j-1}(e) \forall e \in V \setminus \{\mathbf{o}^j\}$ and $\mathbf{o}^{j-1/2}(\mathbf{o}^j) = 0$.
 - Define \mathbf{o}^j as a k -set that $\mathbf{o}^j(e) = \mathbf{o}^{j-1/2}(e) \forall e \in V \setminus \{e^j\}$ and $\mathbf{o}^j(e^j) = i^j$
 - Define $\mathbf{s}^{j-1/2}$ as a k -set that $\mathbf{s}^{j-1/2}(e) = \mathbf{s}^{j-1}(e) \forall e \in V \setminus \{\mathbf{o}^j\}$ and $\mathbf{s}^{j-1/2}(\mathbf{o}^j) = i^j$.

It is trivial that $\mathbf{o}^B = \mathbf{s}^B$. We have:

$$f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \leq f(\mathbf{o}^{j-1}) - f(\mathbf{o}^{j-1/2}) \quad (4)$$

$$\leq f(\mathbf{s}^{j-1/2}) - f(\mathbf{s}^{j-1}) \quad (5)$$

$$\leq \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1/2}) - f(\mathbf{s}^{j-1}) \quad (6)$$

$$\leq \frac{1}{1-\epsilon} F(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \quad (7)$$

$$\leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \quad (8)$$

The inequality (4) is due to f is monotone, thus $f(\mathbf{o}^j) \geq f(\mathbf{o}^{j-1/2})$. The inequality (5) is from k -submodularity of f . The inequality (7) is due to greedy selection. Therefore:

$$\begin{aligned} f(\mathbf{o}) - f(\mathbf{s}) &= \sum_{j=1}^B \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \right) \leq \sum_{j=1}^B \left(\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \right) \\ &\leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}) + \sum_{j=1}^B \frac{2\epsilon}{1-\epsilon} f(\mathbf{s}^j) \leq \frac{1+\epsilon+2\epsilon B}{1-\epsilon} f(\mathbf{s}) \end{aligned}$$

Thus $f(\mathbf{o}) \leq \frac{2+2\epsilon B}{1-\epsilon} f(\mathbf{s})$, which completes the proof. \square

Theorem 6. Given an instance of Noisy \mathbf{MkSC} with input V, k, B and F is an ϵ -estimation of the **non-monotone** k -submodular objective function f . If \mathbf{s} is an output of the Greedy algorithm and \mathbf{o} is an optimal solution, then

$$f(\mathbf{o}) \leq \frac{3+\epsilon+4\epsilon B}{1-\epsilon} f(\mathbf{s}) \quad (9)$$

Proof. We use the same definition of $\mathbf{o}^j, \mathbf{o}^{j-1/2}, \mathbf{s}^j$ as in proof of Theorem 5.

Although f is non-monotone, f is *pairwise-monotone* due to k -submodularity (Ward & Živný, 2016). To be specific, given a k -set \mathbf{x} and $e \notin \text{supp}(\mathbf{x})$, we have

$$\Delta_{e,i} f(\mathbf{x}) + \Delta_{e,j} f(\mathbf{x}) \geq 0 \quad \forall i, j \in [k] \text{ and } i \neq j \quad (10)$$

Let consider a value of $j \in [1, B]$, due to pairwise-monotonicity, there should exist $i' \in [k]$ that $f(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) \geq f(\mathbf{s}^{j-1})$. Moreover, due to greedy selection, $f(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) \leq \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) \leq \frac{1}{1-\epsilon} F(\mathbf{s}^j) \leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j)$. Thus $\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) \geq f(\mathbf{s}^{j-1})$. We consider 2 following cases:

- $\mathbf{o}^j = \mathbf{o}^{j-1}$. In this case $f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) = 0 \leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1})$

- $\mathbf{o}^j \neq \mathbf{o}^{j-1}$. There would be 2 sub-cases:

- $\mathbf{o}^{j-1}(e^j) = 0$. Then let $i' \in [k]$ be an arbitrary number and $i' \neq i^j$. Then

$$f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) = f(\mathbf{o}^{j-1}) + f(\mathbf{o}^{j-1/2} \sqcup \langle e^j, i' \rangle) - 2f(\mathbf{o}^{j-1/2}) - \left(f(\mathbf{o}^{j-1/2} \sqcup \langle e^j, i' \rangle) + f(\mathbf{o}^j) - 2f(\mathbf{o}^{j-1/2}) \right) \quad (11)$$

$$\leq f(\mathbf{o}^{j-1}) + f(\mathbf{o}^{j-1/2} \sqcup \langle e^j, i' \rangle) - 2f(\mathbf{o}^{j-1/2}) \quad (12)$$

$$\leq f(\mathbf{s}^{j-1/2}) + f(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) - 2f(\mathbf{s}^{j-1}) \quad (13)$$

$$\leq \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1/2}) + \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) - 2f(\mathbf{s}^{j-1}) \quad (14)$$

$$\leq \frac{2}{1-\epsilon} F(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1}) \leq 2\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1}) \quad (15)$$

The inequality 12 is due to pairwise-monotonicity of f . The inequality 13 is from k -submodularity and inequality 15 is due to greedy selection.

- $\mathbf{o}^{j-1}(e^j) \neq i^j$. Then:

$$\begin{aligned} f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) &= 2f(\mathbf{o}^{j-1}) - 2f(\mathbf{o}^{j-1/2}) - \left(f(\mathbf{o}^{j-1}) + f(\mathbf{o}^j) - 2f(\mathbf{o}^{j-1/2}) \right) \\ &\leq 2f(\mathbf{o}^{j-1}) - 2f(\mathbf{o}^{j-1/2}) \leq 2f(\mathbf{s}^{j-1/2}) - 2f(\mathbf{s}^{j-1}) \leq \frac{2}{1-\epsilon} F(\mathbf{s}^{j-1/2}) - 2f(\mathbf{s}^{j-1}) \\ &\leq \frac{2}{1-\epsilon} F(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1}) \leq 2\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1}) \end{aligned}$$

Both cases imply that $f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \leq 2\frac{1+\epsilon}{1-\epsilon}f(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1})$. Therefore:

$$\begin{aligned} f(\mathbf{o}) - f(\mathbf{s}) &= \sum_{j=1}^B \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \right) \leq 2 \sum_{j=1}^B \left(\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \right) \\ &\leq 2 \times \left(\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^B) + \sum_{j=1}^B \frac{2\epsilon}{1-\epsilon} f(\mathbf{s}^j) \right) \leq \frac{2+2\epsilon+4\epsilon B}{1-\epsilon} \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}) \end{aligned}$$

Thus, $f(\mathbf{o}) \leq \left(\frac{(2+2\epsilon+4\epsilon B)(1+\epsilon)}{(1-\epsilon)^2} + 1 \right) f(\mathbf{s})$, which completes the proof. \square

C. DSTREAM Algorithm

In this part, we present in detail omitted proofs of DSTREAM's approximation ratio. There would be 2 separate sub-parts: one is for when f is **monotone** - which we describe omitted proofs of claims in Proposition 1; the other is for when f is **non-monotone** - which we would fully prove the Theorem 2.

C.1. Approximation ratio of DSTREAM when f is monotone

Proof of Claim 1. Denote $\langle e^j, i^j \rangle$ as j -th addition of Alg. 1, i.e $\mathbf{s}^j = \mathbf{s}^{j-1} \sqcup \langle e^j, i^j \rangle$. Define:

$$\mathbf{o}^{j-1/2} = (\mathbf{o} \sqcup \mathbf{s}^j) \sqcup \mathbf{s}^{j-1}$$

Furthermore, define $\mathbf{s}^{j-1/2}$ as:

- If $e^j \in \text{supp}(\mathbf{o})$, let $i^j = \mathbf{o}(e^j)$. Then $\mathbf{s}^{j-1/2} = \mathbf{s}^{j-1} \sqcup \langle e^j, i^j \rangle$
- Otherwise, $\mathbf{s}^{j-1/2} = \mathbf{s}^{j-1}$

We have:

$$f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \leq f(\mathbf{o}^{j-1}) - f(\mathbf{o}^{j-1/2}) \quad (16)$$

$$\leq f(\mathbf{s}^{j-1/2}) - f(\mathbf{s}^{j-1}) \quad (17)$$

$$\leq \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1/2}) - f(\mathbf{s}^{j-1}) \quad (18)$$

$$\leq \frac{1}{1-\epsilon} F(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \quad (19)$$

$$\leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \quad (20)$$

The Inequality 16 is due to monotonicity of f ; the inequality 17 is due to its k -submodularity and the inequality 19 is from selection of Alg. 1 that guarantees $i^j = \text{argmax}_{i \in [k]} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle)$. Therefore:

$$\begin{aligned} f(\mathbf{o}) - f(\mathbf{o}^t) &= \sum_{j=1}^t \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \right) \leq \sum_{j=1}^t \left(\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \right) \\ &= \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}) + \sum_{i=1}^{t-1} \left(\frac{1+\epsilon}{1-\epsilon} - 1 \right) f(\mathbf{s}^i) \leq \frac{1+\epsilon+2B\epsilon}{1-\epsilon} f(\mathbf{s}) \end{aligned}$$

which completes the proof.

Proof of Claim 2. As $\mathbf{s} \sqsubseteq \mathbf{o}^t$, denote $\{\langle u_1, j_1 \rangle, \dots, \langle u_r, j_r \rangle\}$ as a set of elements and their placement in \mathbf{o}^t that are not in \mathbf{s} .

For each $\langle u_i, j_i \rangle$, denote \mathbf{s}_i as \mathbf{s} when Alg. 1 encounters u_i . As u_i was not added into \mathbf{s} , $\frac{F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle)}{1-\epsilon} \leq (|\mathbf{s}_i| + 1) \frac{\rho}{M}$. While $\frac{F(\mathbf{s}_i)}{1-\epsilon} \geq |\mathbf{s}_i| \frac{\rho}{M}$, we have:

$$\frac{F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle)}{1 - \epsilon} - \frac{F(\mathbf{s}_i)}{1 + \epsilon} \leq \frac{o}{M} + \frac{2\epsilon F(\mathbf{s}_i)}{1 - \epsilon^2}$$

Denote $\mathbf{u}_i = \mathbf{s} \sqcup \{\langle u_1, j_1 \rangle, \dots, \langle u_i, j_i \rangle\}$. We have

$$\begin{aligned} f(\mathbf{o}^t) - f(\mathbf{s}) &= \sum_{i=1}^{B-t} \left(f(\mathbf{u}_i) - f(\mathbf{u}_{i-1}) \right) \leq \sum_{i=1}^{B-t} \left(f(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - f(\mathbf{s}_i) \right) \\ &\leq \sum_{i=1}^{B-t} \left(\frac{1}{1 - \epsilon} F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - \frac{1}{1 + \epsilon} F(\mathbf{s}_i) \right) \leq \sum_{i=1}^{B-t} \left(\frac{o}{M} + \frac{2\epsilon}{1 - \epsilon^2} F(\mathbf{s}_i) \right) \leq \frac{1}{M} f(\mathbf{o}) + \frac{2\epsilon B}{1 - \epsilon} f(\mathbf{s}) \end{aligned}$$

which complete the proof.

C.2. Approximation of DSTREAM when f is non-monotone

Proof of Theorem 2 We still use definition of $\mathbf{o}^j, \mathbf{o}^{j-1/2}, e^j, i^j, \mathbf{s}^j, \mathbf{s}^{j-1/2}$ as in the monotone proof. Also, for simplicity, we first prove the approximation ratio of Alg. 1 (assume $f(\mathbf{o})$ is known).

If $t = B$, $f(\mathbf{s}) \geq \frac{1}{1+\epsilon} \max_{i \in \{1, \dots, B\}} F(\mathbf{s}^i) \geq \frac{1}{1+\epsilon} F(\mathbf{s}^B) \geq \frac{1-\epsilon}{1+\epsilon} B \frac{o}{M} \geq \frac{1-\epsilon}{1+\epsilon} \frac{1}{(1+\gamma)M} f(\mathbf{o})$. The rest of the proof will focus on case $t < B$.

Due to pairwise-monotonicity of a k -submodular function, for any $j \in \{1, \dots, t\}$, there exists no pair $i_1 \neq i_2 \in [k]$ that $\Delta_{e^j, i_1} f(\mathbf{s}^{j-1}) < 0$ and $\Delta_{e^j, i_2} f(\mathbf{s}^{j-1}) < 0$. Therefore $\max_{i \in [k]} f(\mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle) \geq f(\mathbf{s}^{j-1})$. Thus

$$f(\mathbf{s}^j) \geq \frac{F(\mathbf{s}^j)}{1 + \epsilon} = \frac{\max_{i \in [k]} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle)}{1 + \epsilon} \geq \frac{1 - \epsilon}{1 + \epsilon} \max_{i \in [k]} f(\mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle) \geq \frac{1 - \epsilon}{1 + \epsilon} f(\mathbf{s}^{j-1}) \quad (21)$$

Let's consider relation between \mathbf{o}^{j-1} and \mathbf{o}^j , there are 2 cases:

- $|supp(\mathbf{o}^{j-1})| < |supp(\mathbf{o}^j)|$, which means $e^j \notin supp(\mathbf{o})$. We randomly pick $i' \in [k]$ and $i' \neq i^j$, we have

$$\begin{aligned} f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) &= f(\mathbf{o}^{j-1} \sqcup \langle e^j, i' \rangle) - f(\mathbf{o}^{j-1}) - (f(\mathbf{o}^j) + f(\mathbf{o}^{j-1} \sqcup \langle e^j, i' \rangle) - 2f(\mathbf{o}^{j-1})) \\ &\leq f(\mathbf{o}^{j-1} \sqcup \langle e^j, i' \rangle) - f(\mathbf{o}^{j-1}) \leq f(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) - f(\mathbf{s}^{j-1}) \\ &\leq \frac{1}{1 - \epsilon} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i' \rangle) - f(\mathbf{s}^{j-1}) \leq \frac{1}{1 - \epsilon} F(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \\ &\leq \frac{1 + \epsilon}{1 - \epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \end{aligned}$$

- $|supp(\mathbf{o}^{j-1})| = |supp(\mathbf{o}^j)|$, which means $e^j \in supp(\mathbf{o})$. We have 2 sub-cases

- $\mathbf{o}^{j-1}(e^j) = i^j$. Then $f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) = 0 \leq \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1})$
- $\mathbf{o}^{j-1}(e^j) \neq i^j$, we have

$$\begin{aligned} f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) &= 2f(\mathbf{o}^{j-1}) - 2f(\mathbf{o}^{j-1/2}) - (f(\mathbf{o}^{j-1}) + f(\mathbf{o}^j) - 2f(\mathbf{o}^{j-1/2})) \\ &\leq 2f(\mathbf{o}^{j-1}) - 2f(\mathbf{o}^{j-1/2}) \leq 2f(\mathbf{s}^{j-1/2}) - 2f(\mathbf{s}^{j-1}) \leq 2 \frac{1 + \epsilon}{1 - \epsilon} f(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1}) \end{aligned}$$

Therefore, in overall $f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \leq 2 \frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}^j) - 2f(\mathbf{s}^{j-1})$. We have

$$f(\mathbf{o}) - f(\mathbf{o}^t) = \sum_{j=1}^t \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) \right) \leq \sum_{j=1}^t 2 \left(\frac{1 + \epsilon}{1 - \epsilon} f(\mathbf{s}^j) - f(\mathbf{s}^{j-1}) \right) \quad (22)$$

$$= 2 \frac{1 + \epsilon}{1 - \epsilon} f(\mathbf{s}^t) + 2 \sum_{i=1}^{t-1} \frac{2\epsilon}{1 - \epsilon} f(\mathbf{s}^i) \leq \frac{2 + 2\epsilon + 4\epsilon B}{1 - \epsilon} \max_{i \leq t} f(\mathbf{s}^i) \leq \frac{(1 + \epsilon)(2 + 2\epsilon + 4\epsilon B)}{(1 - \epsilon)^2} f(\mathbf{s}) \quad (23)$$

Also, similar to monotone case, as $\mathbf{s}^t \sqsubseteq \mathbf{o}^t$, denote $\{\langle u_1, j_1 \rangle, \dots, \langle u_r, j_r \rangle\}$ as a set of elements and their placement in \mathbf{o}^t that are not in \mathbf{s}^t . For each $\langle u_i, j_i \rangle$, denote \mathbf{s}_i as \mathbf{s}^t when Alg. 1 encounters u_i . Denote $\mathbf{u}_i = \mathbf{s}^t \sqcup \{\langle u_1, j_1 \rangle, \dots, \langle u_i, j_i \rangle\}$. We have

$$f(\mathbf{o}^t) - f(\mathbf{s}^t) = \sum_{i=1}^{B-t} \left(f(\mathbf{u}^i) - f(\mathbf{u}^{i-1}) \right) \leq \sum_{i=1}^{B-t} \left(f(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - f(\mathbf{s}_i) \right) \quad (24)$$

$$\leq \sum_{i=1}^{B-t} \left(\frac{F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle)}{1-\epsilon} - \frac{F(\mathbf{s}_i)}{1+\epsilon} \right) \leq \sum_{i=1}^{B-t} \left(\frac{o}{M} + \frac{2\epsilon}{1-\epsilon^2} F(\mathbf{s}^i) \right) \quad (25)$$

$$\leq \frac{1}{M} f(\mathbf{o}) + \frac{2\epsilon B}{1-\epsilon^2} F(\mathbf{s}) \leq \frac{1}{M} f(\mathbf{o}) + \frac{2\epsilon B}{1-\epsilon} f(\mathbf{s}) \quad (26)$$

Combining Equ. 23 and 26, we have

$$\left(1 - \frac{1}{M}\right) f(\mathbf{o}) \leq \left(\frac{(1+\epsilon)(2+2\epsilon+4\epsilon B)}{(1-\epsilon)^2} + \frac{2\epsilon B}{1-\epsilon} + \frac{1+\epsilon}{1-\epsilon} \right) f(\mathbf{s}) = \frac{3+4\epsilon+6\epsilon B+\epsilon^2+2\epsilon^2 B}{(1-\epsilon)^2} f(\mathbf{s})$$

The approximation ratio of DSTREAM when discarding assumption of known $f(\mathbf{o})$ is trivially follows as in the proof of the monotone case.

D. RSTREAM Algorithm

In this part, we present in detail omitted proofs of RSTREAM's approximation ratio. There would be 2 separate sub-parts: one is for when f is **monotone** - which we describe omitted proofs of claims in Theorem 3; the other is for when f is **non-monotone** - which we would fully prove the Theorem 4.

Note that in proofs related to RSTREAM, we abuse the notation and simply write t and \mathbf{s}^i to indicate $t_{j,\epsilon'}$ and $\mathbf{s}_{j,\epsilon'}^i$ in Alg. 3 when o is estimated by $M(1+\gamma)^j$ that satisfies $f(\mathbf{o}) \geq oB \geq \frac{f(\mathbf{o})}{1+\gamma}$; and $\epsilon' = \epsilon$.

Similar to DSTREAM's proof, denote $\langle e^j, i^j \rangle$ as j -th addition of Alg. 3, i.e $\mathbf{s}^j = \mathbf{s}^{j-1} \sqcup \langle e^j, i^j \rangle$. Define:

$$\mathbf{o}^{j-1/2} = (\mathbf{o} \sqcup \mathbf{s}^j) \sqcup \mathbf{s}^{j-1}$$

Furthermore, define $\mathbf{s}^{j-1/2}$ as:

- If $e^j \in \text{supp}(\mathbf{o})$, let i^j be an index of a set containing e^j in \mathbf{o} . Then $\mathbf{s}^{j-1/2} = \mathbf{s}^{j-1} \sqcup \langle e^j, i^j \rangle$
- Otherwise, $\mathbf{s}^{j-1/2} = \mathbf{s}^{j-1}$

D.1. Approximation ratio of RSTREAM when f is monotone

Proof of Lemma 1 We have:

$$\begin{aligned} F(\mathbf{s}) &= F(\mathbf{s}^B) = \sum_{i=1}^B \left(F(\mathbf{s}^i) - F(\mathbf{s}^{i-1}) \right) = \sum_{i=1}^B \left(F(\mathbf{s}^i) - \frac{1-\epsilon}{1+\epsilon} F(\mathbf{s}^{i-1}) \right) - \frac{2\epsilon}{1+\epsilon} \sum_{i=1}^{B-1} F(\mathbf{s}^i) \\ &\geq (1-\epsilon) \sum_{i=1}^B \frac{o}{M} - 2\epsilon \sum_{i=1}^{B-1} f(\mathbf{s}^i) \geq \frac{(1-\epsilon)f(\mathbf{o})}{(1+\gamma)M} - 2\epsilon B f(\mathbf{s}) \end{aligned}$$

Therefore, $\frac{(1+\gamma)(1+\epsilon+2\epsilon B)M}{1-\epsilon} f(\mathbf{s}) \geq f(\mathbf{o})$

Proof of Claim 3 Let's consider Alg. 3 before adding e^j into \mathbf{s}^{j-1} . Denote $d_i = \frac{1}{1-\epsilon} F(\mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle) - \frac{1}{1+\epsilon} F(\mathbf{s}^{j-1})$. We consider the following cases:

- $|supp(\mathbf{o}^{j-1})| = |supp(\mathbf{o}^j)|$, which means $e^j \in supp(\mathbf{o})$. Let $p = \mathbf{o}^{j-1}(e^j)$.

Let $I \subseteq [k]$ be a set of values of i that $d_i > \frac{o}{M}$, $T = |I|$.

We define \mathbf{o}_i^j as a k -set that $\mathbf{o}_i^j(e) = \mathbf{o}^j(e) \forall e \in V \setminus \{e^j\}$ and $\mathbf{o}_i^j(e^j) = i$. Define $\mathbf{o}^{j-1/2}$ as a k -set that $\mathbf{o}^{j-1/2}(e) = \mathbf{o}^j(e) \forall e \in V \setminus \{e^j\}$ and $\mathbf{o}^{j-1/2}(e^j) = 0$.

Let $\mathbf{s}_i^j = \mathbf{s}^{j-1} \sqcup \langle e^j, i \rangle$. We have two following sub-cases:

- $d_p \geq \frac{o}{M}$, then if $T = 1$, $\mathbf{o}^{j-1} = \mathbf{o}^j$ and $f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j) = 0 \leq (1 - \frac{1}{k})(\frac{1+\epsilon}{1-\epsilon}E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon}f(\mathbf{s}^{j-1}))$. Thus, we assume $T > 1$, in which we have

$$f(\mathbf{o}^{j-1}) - E[f(\mathbf{o}^j)] = \frac{1}{D} \sum_{i \in I \setminus \{p\}} (f(\mathbf{o}^{j-1}) - f(\mathbf{o}_i^j)) d_i^{T-1} \leq \frac{1}{D} \sum_{i \in I \setminus \{p\}} (f(\mathbf{o}^{j-1}) - f(\mathbf{o}^{j-1/2})) d_i^{T-1} \quad (27)$$

$$\leq \frac{1}{D} \sum_{i \in I \setminus \{p\}} (f(\mathbf{s}_p^j) - f(\mathbf{s}^{j-1})) d_i^{T-1} \leq \frac{1}{D} \sum_{i \in I \setminus \{p\}} d_p d_i^{T-1} \quad (28)$$

$$\leq \frac{1}{D} \left(1 - \frac{1}{T}\right) \sum_{i \in I} d_i^T \quad (29)$$

$$= \frac{1}{D} \left(1 - \frac{1}{T}\right) \sum_{i \in I} \left(\frac{1}{1-\epsilon} F(\mathbf{s}_i^j) - \frac{1}{1+\epsilon} F(\mathbf{s}^{j-1})\right) d_i^{T-1} \quad (30)$$

$$\leq \frac{1}{D} \left(1 - \frac{1}{T}\right) \sum_{i \in I} \left(\frac{1+\epsilon}{1-\epsilon} f(\mathbf{s}_i^j) - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})\right) d_i^{T-1} \quad (31)$$

$$\leq \left(1 - \frac{1}{k}\right) \left(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})\right) \quad (32)$$

Inequality 29 comes from AM-GM inequality, defined as:

Theorem 7. (Hirschhorn, 2007) Given n non-negative numbers x_1, \dots, x_n

$$x_1 + \dots + x_n \geq n \sqrt[n]{x_1 \times \dots \times x_n}$$

Thus, directly apply the theorem given us $d_p d_i^{T-1} \leq \frac{1}{T} (d_p^T + (T-1)d_i^T)$.

- $d_p < \frac{o}{M}$, which also means $d_p \leq \frac{o}{M} \leq E[d_{ij}] \leq \frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})$. So:

$$\begin{aligned} f(\mathbf{o}^{j-1}) - E[f(\mathbf{o}^j)] &\leq f(\mathbf{s}_p^j) - f(\mathbf{s}^{j-1}) \leq d_p \leq \left(1 - \frac{1}{k}\right) E[d_{ij}] + \frac{1}{k} \frac{o}{M} \\ &\leq \left(1 - \frac{1}{k}\right) \left(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})\right) + \frac{1}{k} \frac{o}{M} \end{aligned}$$

- $|supp(\mathbf{o}^{j-1})| < |supp(\mathbf{o}^j)|$, then $e^j \notin supp(\mathbf{o})$. Due to monotonicity of f .

$$f(\mathbf{o}^{j-1}) - E[f(\mathbf{o}^j)] \leq 0 \leq \left(1 - \frac{1}{k}\right) \left(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})\right)$$

Therefore, in overall $f(\mathbf{o}^{j-1}) - E[f(\mathbf{o}^j)] \leq (1 - \frac{1}{k})(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})) + \frac{o}{kM}$

Proof of Claim 4 We have

$$\begin{aligned} f(\mathbf{o}) - E[f(\mathbf{o}^t)] &= \sum_{j=1}^t E[f(\mathbf{o}^{j-1}) - f(\mathbf{o}^j)] \leq \sum_{j=1}^t \left(\left(1 - \frac{1}{k}\right) \left(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} E[f(\mathbf{s}^{j-1})]\right) + \frac{o}{kM} \right) \\ &\leq \left(1 - \frac{1}{k}\right) \left(\frac{1+\epsilon}{1-\epsilon} E[f(\mathbf{s})] + \sum_{j=1}^{t-1} \frac{4\epsilon}{1-\epsilon^2} E[f(\mathbf{s}^j)]\right) + \frac{oB}{kM} \\ &\leq \left(1 - \frac{1}{k}\right) \frac{(1+\epsilon)^2 + 4B\epsilon}{1-\epsilon^2} E[f(\mathbf{s})] + \frac{1}{kM} f(\mathbf{o}) \end{aligned}$$

which completes the proof.

Proof of Claim 5 Denote $\{\langle u_1, j_1 \rangle, \dots, \langle u_r, j_r \rangle\}$ as a set of elements and their placements in \mathbf{o}^t that are not in \mathbf{s} .

For each $\langle u_i, j_i \rangle$, denote \mathbf{s}_i as \mathbf{s} when Alg. 3 encounters u_i . As u_i was not added into \mathbf{s}_i , $\frac{1}{1-\epsilon}F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - \frac{1}{1+\epsilon}F(\mathbf{s}_i) < \frac{o}{M}$

Let $\mathbf{u}_i = \mathbf{s} \sqcup \{\langle u_1, j_1 \rangle, \dots, \langle u_i, j_i \rangle\}$. We have

$$\begin{aligned} \mathbb{E}[f(\mathbf{o}^t)] - \mathbb{E}[f(\mathbf{s})] &= \sum_{i=1}^{B-t} \mathbb{E}[f(\mathbf{u}_i) - f(\mathbf{u}_{i-1})] \leq \sum_{i=1}^{B-t} \mathbb{E}[f(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - f(\mathbf{s}^i)] \\ &\leq \sum_{i=1}^{B-t} \mathbb{E}\left[\frac{1}{1-\epsilon}F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - \frac{1}{1+\epsilon}F(\mathbf{s}_i)\right] \leq \sum_i^{B-t} \frac{o}{M} \leq \frac{1}{M}f(\mathbf{o}) \end{aligned}$$

which completes the proof.

D.2. Approximation ratio of RSTREAM when f is non-monotone

Proof of Theorem 4 Similar to monotone case, the key of our proof is to show that:

$$\max(\alpha(M), \beta(M))\mathbb{E}[f(\mathbf{s}^t)] \geq f(\mathbf{o}) \quad (33)$$

If $t = B$, we have:

$$f(\mathbf{s}) \geq \frac{F(\mathbf{s})}{1+\epsilon} \geq \frac{F(\mathbf{s}^B)}{1+\epsilon} = \frac{1}{1+\epsilon} \sum_{i=1}^B (F(\mathbf{s}^i) - F(\mathbf{s}^{i-1})) \quad (34)$$

$$= \frac{1}{1+\epsilon} \sum_{i=1}^B \left(F(\mathbf{s}^i) - \frac{1-\epsilon}{1+\epsilon}F(\mathbf{s}^{i-1}) \right) - \frac{2\epsilon}{(1+\epsilon)^2} \sum_i^{B-1} F(\mathbf{s}^i) \quad (35)$$

$$\geq \frac{1-\epsilon}{1+\epsilon} \sum_{i=1}^B \frac{o}{M} - \frac{2\epsilon B}{(1+\epsilon)^2} F(\mathbf{s}) \geq \frac{(1-\epsilon)f(\mathbf{o})}{(1+\epsilon)(1+\gamma)M} - \frac{2\epsilon B}{1+\epsilon} f(\mathbf{s}) \quad (36)$$

Therefore, with $t = B$, $\frac{(1+\epsilon+2\epsilon B)(1+\gamma)M}{1-\epsilon}f(\mathbf{s}) \geq f(\mathbf{o})$. The rest of the proof would focus on when $t < B$. We re-use notations $\mathbf{o}_i^j, \mathbf{s}_i^j$ as in the monotone proof and still compare \mathbf{o}^{j-1} and \mathbf{o}^j , we have two following cases.

- $|\text{supp}(\mathbf{o}^{j-1})| < |\text{supp}(\mathbf{o}^j)|$, which means $e^j \notin \text{supp}(\mathbf{o})$. We consider 2 sub-cases:

– If $T < k$, which means there exists $i' \in [k]$ such that $\frac{F(\mathbf{s}_{i'}^j)}{1-\epsilon} - \frac{F(\mathbf{s}_{i'}^{j-1})}{1+\epsilon} < \frac{o}{M} \leq \frac{\mathbb{E}[F(\mathbf{s}^j)]}{1-\epsilon} - \frac{F(\mathbf{s}^{j-1})}{1+\epsilon}$. Then

$$\begin{aligned} f(\mathbf{o}^{j-1}) - \mathbb{E}[f(\mathbf{o}^j)] &= f(\mathbf{o}_{i'}^j) - f(\mathbf{o}^{j-1}) - \left(\mathbb{E}[f(\mathbf{o}^j)] + f(\mathbf{o}_{i'}^j) - 2f(\mathbf{o}^{j-1}) \right) \\ &\leq f(\mathbf{o}_{i'}^j) - f(\mathbf{o}^{j-1}) \leq f(\mathbf{s}_{i'}^j) - f(\mathbf{s}^{j-1}) \leq \frac{F(\mathbf{s}_{i'}^j)}{1-\epsilon} - \frac{F(\mathbf{s}^{j-1})}{1+\epsilon} \\ &\leq \frac{\mathbb{E}[F(\mathbf{s}^j)]}{1-\epsilon} - \frac{F(\mathbf{s}^{j-1})}{1+\epsilon} \leq \frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1}) \end{aligned}$$

– If $T = k$, define a permutation $\pi : [k] \rightarrow [k]$ such that $\pi(i) \neq i$ for all $i \in [k]$. Then

$$\begin{aligned}
 f(\mathbf{o}^{j-1}) - \mathbb{E}[f(\mathbf{o}^j)] &= \frac{1}{D} \sum_{i \in [k]} \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}_i^j) \right) d_i^{T-1} \\
 &= \frac{1}{D} \sum_{i \in [k]} \left(f(\mathbf{o}_{\pi(i)}^j) - f(\mathbf{o}^{j-1}) - (f(\mathbf{o}_{\pi(i)}^j) + f(\mathbf{o}_i^j) - 2f(\mathbf{o}^{j-1})) \right) d_i^{T-1} \\
 &\leq \frac{1}{D} \sum_{i \in [k]} \left(f(\mathbf{o}_{\pi(i)}^j) - f(\mathbf{o}^{j-1}) \right) d_i^{T-1} \leq \frac{1}{D} \sum_i \left(f(\mathbf{s}_{\pi(i)}^j) - f(\mathbf{s}^{j-1}) \right) d_i^{T-1} \\
 &\leq \frac{1}{D} \sum_{i \in [k]} d_{\pi(i)} d_i^{T-1} \leq \frac{1}{D} \sum_{i \in [k]} d_i^T = \frac{1}{1-\epsilon} \mathbb{E}[F(\mathbf{s}^j)] - \frac{1}{1+\epsilon} F(\mathbf{s}^{j-1}) \\
 &\leq \frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1})
 \end{aligned}$$

• $|\text{supp}(\mathbf{o}^{j-1})| = |\text{supp}(\mathbf{o}^j)|$. We re-use notation $p, I, \mathbf{o}^{j-1/2}$ as in the monotone case and consider other two sub-cases:

– If $d_p \geq \frac{o}{M}$, then

$$\begin{aligned}
 f(\mathbf{o}^{j-1}) - \mathbb{E}[f(\mathbf{o}^j)] &= \frac{1}{D} \sum_{i \in I \setminus \{p\}} \left(f(\mathbf{o}^{j-1}) - f(\mathbf{o}_i^j) \right) d_i^{T-1} \\
 &= \frac{1}{D} \sum_{i \in I \setminus \{p\}} \left(2f(\mathbf{o}_p^j) - 2f(\mathbf{o}^{j-1/2}) - (f(\mathbf{o}_p^j) + f(\mathbf{o}_i^j) - 2f(\mathbf{o}^{j-1/2})) \right) d_i^{T-1} \\
 &\leq \frac{2}{D} \sum_{i \in I \setminus \{p\}} \left(f(\mathbf{s}_p^j) - f(\mathbf{s}^{j-1}) \right) d_i^{T-1} \leq \frac{2}{D} \sum_{i \in I \setminus \{p\}} d_p d_i^{T-1} \leq \frac{2}{D} \left(1 - \frac{1}{T} \right) \sum_{i \in I} d_i^T \\
 &\leq \left(2 - \frac{2}{k} \right) \left(\frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1}) \right)
 \end{aligned}$$

– If $d_p = 0$, which also means $\frac{1}{1-\epsilon} F(\mathbf{s}_p^j) - \frac{1}{1+\epsilon} F(\mathbf{s}^{j-1}) < \frac{o}{M} \leq \frac{1}{1-\epsilon} \mathbb{E}[F(\mathbf{s}^j)] - \frac{1}{1+\epsilon} F(\mathbf{s}^{j-1})$. So:

$$\begin{aligned}
 f(\mathbf{o}^{j-1}) - \mathbb{E}[f(\mathbf{o}^j)] &= 2f(\mathbf{o}_p^j) - 2f(\mathbf{o}^{j-1/2}) - \left(f(\mathbf{o}_p^j) + \mathbb{E}[f(\mathbf{o}^j)] - 2f(\mathbf{o}^{j-1/2}) \right) \\
 &\leq 2f(\mathbf{o}_p^j) - 2f(\mathbf{o}^{j-1/2}) \leq 2f(\mathbf{s}_p^j) - 2f(\mathbf{s}^{j-1}) \\
 &\leq \frac{2}{1-\epsilon} F(\mathbf{s}_p^j) - \frac{2}{1+\epsilon} F(\mathbf{s}^{j-1}) \leq \left(2 - \frac{2}{k} \right) \left(\frac{\mathbb{E}[F(\mathbf{s}^j)]}{1-\epsilon} - \frac{F(\mathbf{s}^{j-1})}{1+\epsilon} \right) + \frac{2o}{kM} \\
 &\leq \left(2 - \frac{2}{k} \right) \left(\frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1}) \right) + \frac{2o}{kM}
 \end{aligned}$$

Since $k \geq 2$, in overall, $f(\mathbf{o}^{j-1}) - \mathbb{E}[f(\mathbf{o}^j)] \leq \left(2 - \frac{2}{k} \right) \left(\frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1}) \right) + \frac{2o}{kM}$. We have

$$f(\mathbf{o}) - \mathbb{E}[f(\mathbf{o}^t)] = \sum_{j=1}^t \left(\mathbb{E}[f(\mathbf{o}^{j-1})] - \mathbb{E}[f(\mathbf{o}^j)] \right) \leq \sum_{j=1}^t \left(\left(2 - \frac{2}{k} \right) \left(\frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^j)] - \frac{1-\epsilon}{1+\epsilon} f(\mathbf{s}^{j-1}) \right) + \frac{2o}{kM} \right) \quad (37)$$

$$= \left(2 - \frac{2}{k} \right) \left(\frac{1+\epsilon}{1-\epsilon} \mathbb{E}[f(\mathbf{s}^t)] + \sum_{i=1}^{t-1} \frac{4\epsilon}{1-\epsilon^2} \mathbb{E}[f(\mathbf{s}^i)] \right) + \frac{2oB}{kM} \quad (38)$$

$$\leq \left(2 - \frac{2}{k} \right) \frac{(1+\epsilon)^2 + 4\epsilon B}{1-\epsilon^2} \max_{i \leq t} \mathbb{E}[f(\mathbf{s}^i)] + \frac{2}{kM} f(\mathbf{o}) \quad (39)$$

$$\leq \left(2 - \frac{2}{k} \right) \frac{(1+\epsilon)^2 + 4\epsilon B}{(1-\epsilon)^2} \mathbb{E}[f(\mathbf{s})] + \frac{2}{kM} f(\mathbf{o}) \quad (40)$$

Also, similar to monotone case, as $\mathbf{s}^t \sqsubseteq \mathbf{o}^t$, denote $\{\langle u_1, j_1 \rangle, \dots, \langle u_r, j_r \rangle\}$ as a set of elements and their placement in \mathbf{o}^t that are not in \mathbf{s}^t . For each $\langle u_i, j_i \rangle$, denote \mathbf{s}_i as \mathbf{s}^t when Alg. 1 encounters u_i . Denote $\mathbf{u}_i = \mathbf{s}^t \sqcup \{\langle u_i, j_i \rangle\}$. We have

$$\mathbb{E}[f(\mathbf{o}^t)] - \mathbb{E}[f(\mathbf{s}^t)] = \sum_{i=1}^{B-t} \left(\mathbb{E}[f(\mathbf{u}^i) - f(\mathbf{u}^{i-1})] \right) \leq \sum_{i=1}^{B-t} \left(\mathbb{E}[f(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle) - f(\mathbf{s}_i)] \right) \quad (41)$$

$$\leq \sum_{i=1}^{B-t} \mathbb{E} \left[\frac{F(\mathbf{s}_i \sqcup \langle u_i, j_i \rangle)}{1 - \epsilon} - \frac{F(\mathbf{s}_i)}{1 + \epsilon} \right] \leq \sum_{i=1}^{B-t} \frac{o}{M} \leq \frac{1}{M} f(\mathbf{o}) \quad (42)$$

Combining Equ. 40 and 42, we have

$$f(\mathbf{o}) \leq \frac{M}{kM - k - 2} \frac{(3k - 2)(1 + \epsilon)^2 + (8k - 8)\epsilon B}{(1 - \epsilon)^2} \mathbb{E}[f(\mathbf{s})]$$

The approximation ratio of RSTREAM when discarding assumption of known $f(\mathbf{o})$ trivially follows as in the proof of the monotone case.

E. Sampling Method for Influence Maximization with k topics

In this part, we would present the sampling method that helps obtaining $F(\mathbf{s})$ satisfying $(1 - \epsilon)\mathbb{I}(\mathbf{s}) \leq F(\mathbf{s}) \leq (1 + \epsilon)\mathbb{I}(\mathbf{s})$ with high probability. We adopt a concept of Reverse Influence Sampling (RIS) (Borgs et al., 2014) to the problem as follows:

Given a social network $G = (V, E)$ and $w_{u,v}^i$ is a weight of edge (u, v) on topic i , a random Reverse Reachable (RR) sample $\mathcal{R} = \{R_1, \dots, R_k\}$ is generated from G by: (1) selecting a random node $v \in V$; (2) generating sample graph $\{g_1, g_2, \dots, g_k\}$ from G , where g_i is generated from weights $\{w_{u,v}^i\}$ for all $(u, v) \in E$; (3) return $\mathcal{R} = \{R_1, \dots, R_k\}$ where R_i is a set of nodes that can reach v in g_i . A k -set $\mathbf{s} = \{S_1, \dots, S_k\}$ would activate the random sample \mathcal{R} iff there exists $i \in [k]$ that $S_i \cap R_i \neq \emptyset$. For simplicity, denote $\mathbf{s} \triangleright \mathcal{R}$ as an indicator variable whether \mathbf{s} activates \mathcal{R} . Using a similar proof as **Observation 3.2** (Borgs et al., 2014), we have $\mathbb{I}(\mathbf{s}) = |V| \cdot \Pr_{\mathcal{R}}[\mathbf{s} \triangleright \mathcal{R}]$.

To estimate $\mathbb{I}(\mathbf{s})$, we generate multiple samples $\mathcal{R}_1, \dots, \mathcal{R}_n$ and let $F(\mathbf{s}) = \frac{|V|}{n} \sum_{i=1}^n \mathbf{s} \triangleright \mathcal{R}_i$. We apply Chernoff Bound² to bound the number of samples, which guarantee $F(\mathbf{s})$ is an ϵ -estimate of $\mathbb{I}(\mathbf{s})$ with high probability. The Chernoff Bound theorem is stated as follows.

Theorem 8. (Chernoff bound) Suppose X_1, \dots, X_n are independent random variables taking values in $\{0, 1\}$. Let X denote their sum and let $\mu = \mathbb{E}[X]$. Then for $\epsilon \in [0, 1]$ we have

$$\Pr\left(X \leq (1 - \epsilon)\mu\right) \leq e^{-\frac{\epsilon^2\mu}{2}}$$

$$\Pr\left(X \geq (1 + \epsilon)\mu\right) \leq e^{-\frac{\epsilon^2\mu}{3}}$$

Therefore, the number of samples that helps us obtaining ϵ -estimate of $\mathbb{I}(\mathbf{s})$ is stated in the following lemma.

Lemma 3. Given a seed set \mathbf{s} , by generating at least $n = \frac{3|V|}{\epsilon^2|\mathbf{s}|} \ln \frac{1}{1 - \sqrt{\lambda}}$ RR samples, $F(\mathbf{s})$ is ϵ -estimate of $\mathbb{I}(\mathbf{s})$ with probability at least λ .

Proof. A slight change in algebra in Chernoff bound helps us obtaining:

$$\Pr\left((1 - \epsilon)\mathbb{I}(\mathbf{s}) \leq F(\mathbf{s}) \leq (1 + \epsilon)\mathbb{I}(\mathbf{s})\right) \geq \left(1 - e^{-\frac{\epsilon^2\mathbb{I}(\mathbf{s})n}{3|V|}}\right)^2 \quad (43)$$

Since $\mathbb{I}(\mathbf{s}) \geq |\mathbf{s}|$, $n \geq \frac{3|V|}{\epsilon^2|\mathbf{s}|} \ln \frac{1}{1 - \sqrt{\lambda}}$ guarantees $F(\mathbf{s})$ is ϵ -estimate of $\mathbb{I}(\mathbf{s})$ with probability at least λ \square

In experiment, we set $\lambda = 0.8$ whenever the algorithms query $\mathbb{I}(\cdot)$.

²<http://math.mit.edu/~goemans/18310S15/chernoff-notes.pdf>