

# $k$ -means++ under Approximation Stability

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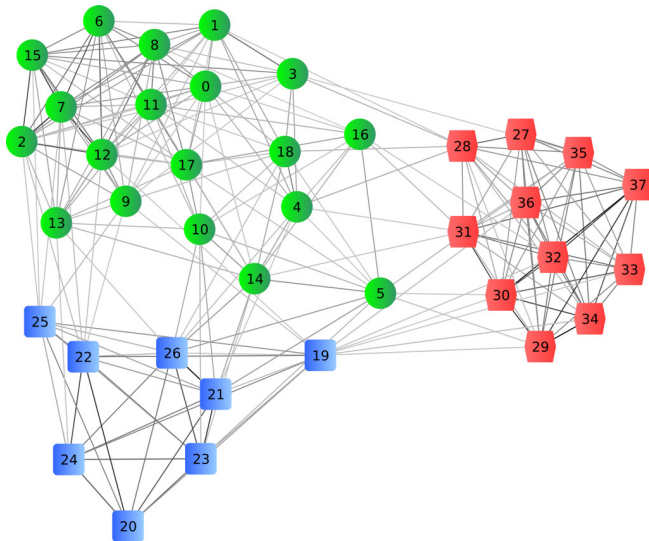
# Agenda

- The clustering problem and  $k$ -means clustering
- Lloyd's algorithm and  $k$ -means++
- Approximation stability and distance between clusterings
- Our contributions
- Analysis of  $k$ -means++
- Conclusion and future work

# The clustering problem

- Given a set of data points, we need to group them together so that *similar* points are in the same group and *dissimilar* points are in different groups.
- Typically, these points live on a *metric space*.
- These groups are called *clusters*.
- There is an *objective function* which has to be optimized.

# Example of a clustering



## $k$ -means clustering problem

- Suppose we have a  $k$ -clustering  $\mathcal{C} = \{C_1, \dots, C_k\}$ .
- The point  $c_i$  is the center of cluster  $C_i$ .
- A point  $x$  is assigned to cluster  $C_i$  if  $d(x, c_i) \leq d(x, c_j)$  for any  $j \neq i$ .
- For the  $k$ -means clustering, we have to minimize the following objective function.

$$\Phi(\mathcal{C}) = \sum_{i=1}^k \sum_{x \in C_i} d(x, c_i)^2.$$

# Lloyd's algorithm

- 1 Choose  $k$  initial centers  $\mathcal{C} = \{c_1, \dots, c_k\}$  arbitrarily.
- 2 For each  $i \in \{1, \dots, k\}$ , set the cluster  $C_i$  to be the set of points in  $\mathcal{X}$  that are closer to  $c_i$  than to  $c_j$  for any  $j \neq i$ .
- 3 For each  $i \in \{1, \dots, k\}$ , set  $c_i$  to be the centroid of all points in  $C_i$ .

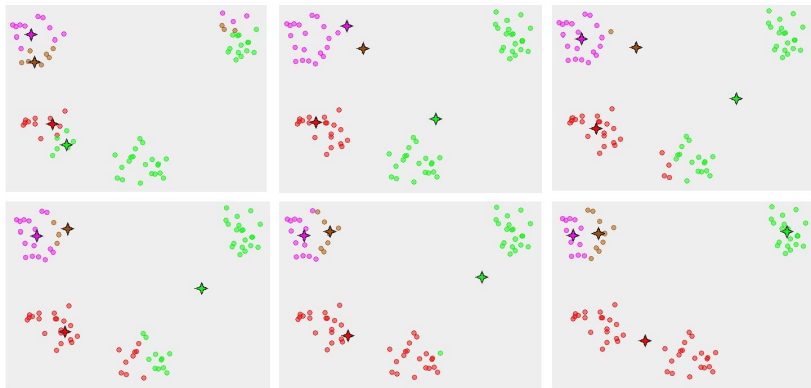
$$c_i = \frac{1}{|C_i|} \sum_{x \in C_i} x.$$

- 4 Repeat Steps 2 and 3 until  $\mathcal{C}$  does not change.

# Problems with Lloyd's algorithm

- Since it is a heuristic algorithm, there is no guarantee that it will converge to the global optimum.
- The result depends on the initial clusters.
- There exist certain point sets (even on the plane), on which the algorithm takes exponential time ( $2^{\Omega(n)}$ ) to converge.
- However, the smoothed running time of  $k$ -means is polynomial.
- $k$ -means assumes that the clusters are spherical that are separable in a way so that the mean value converges towards the cluster center.

# $k$ -means convergence to a local minimum





## $k$ -means++: Initialization of cluster centers

- 1 Choose the first center  $c_1$  uniformly at random from  $\mathcal{X}$ .
- 2 Choose the next center  $c_i$  with probability  $\frac{D(c_i)^2}{\sum_{x \in S} D(x)^2}$ .
- 3 Here  $D(x)$  is the shortest distance from a point  $x$  to the closest center we have already chosen.
- 4 Repeat Step 2, until  $k$  centers are chosen.

# Performance of $k$ -means++

- $k$ -means++ is  $O(\log k)$ -competitive in expectation.
- There are examples on which  $k$ -means++ is  $\Omega(\log k)$ -competitive in expectation.
- So, this is a tight analysis.
- Can  $k$ -means++ do better if the data has additional properties?

## Distance between two clusterings

- Suppose we have two  $k$ -clusterings  $\mathcal{C} = \{C_1, \dots, C_k\}$  and  $\mathcal{C}' = \{C'_1, \dots, C'_k\}$  of a point set  $\mathcal{X}$ .
- Distance between  $\mathcal{C}$  and  $\mathcal{C}'$  is the fraction of points on which they disagree under the optimal matching of clusters in  $\mathcal{C}$  to clusters in  $\mathcal{C}'$ .
- Formally,

$$\text{dist}(\mathcal{C}, \mathcal{C}') = \min_{\sigma \in \mathcal{S}_k} \frac{1}{n} \sum_{i=1}^k |C_i \setminus C'_{\sigma(i)}|,$$

where  $\mathcal{S}_k$  is the set of all permutations  $\sigma : \{1, \dots, k\} \mapsto \{1, \dots, k\}$ .

- Two clusterings  $\mathcal{C}$  and  $\mathcal{C}'$  are  $\epsilon$ -close if  $\text{dist}(\mathcal{C}, \mathcal{C}') < \epsilon$ .

# Approximation stability

- Suppose we are given an objective function  $\Phi$  such as  $k$ -means or  $k$ -median.
- The point set  $\mathcal{X}$  satisfies  $(c, \epsilon)$ -approximation stability if all clusterings  $\mathcal{C}$  with  $\Phi(\mathcal{C}) \leq c \cdot \Phi_{OPT}$  are  $\epsilon$ -close to the target clustering  $\mathcal{C}_T$ .
- At most  $\epsilon$  fraction of points have to be reassigned in  $\mathcal{C}$  to match  $\mathcal{C}_T$ .
- We can assume w.l.o.g that  $\mathcal{C}_T$  is the optimal clustering.

## Our results for large clusters

- Let  $0 < \epsilon, \alpha \leq 1$ . If a dataset satisfies  $(1 + \alpha, \epsilon)$ -approximation stability and each optimal cluster has size at least  $\frac{60\epsilon n}{\alpha^2}$ , then the  $k$ -means++ algorithm gives an  $\delta$ -approximation to the  $k$ -means objective with probability  $\Omega(\frac{1}{k})$ .
- Let  $0 < \epsilon \leq 1$  and  $\alpha > 1$ . If a dataset satisfies  $(1 + \alpha, \epsilon)$ -approximation stability and each optimal cluster has size at least  $70\epsilon n$ , then the  $k$ -means++ algorithm gives an  $\delta$ -approximation to the  $k$ -means objective with probability  $\Omega(\frac{1}{k})$ .
- We also generalize these results for  $k$ -medians with respect to distance measures that satisfy approximate symmetry and approximate triangle inequality.

## Lower bound example for small clusters

- We show that there exists a dataset  $\mathcal{X} \in \mathbb{R}^d$  such that the following holds:
  - $\mathcal{X}$  satisfies the  $(1 + \alpha, \epsilon)$  approximation stability property.
  - $k$ -means++ achieves an approximation factor of  $\frac{1}{2} \log k$  with probability at most  $e^{-\sqrt{k}-o(1)}$ .

## An important result [BBG09]

- Let  $C_1^*, \dots, C_k^*$  denote the optimal  $k$  clusters with respect to the  $k$ -means objective function and let  $c_1^*, \dots, c_k^*$  denote the centroids of these optimal clusters.
- For a point  $x \in \mathcal{X}$ , let  $w(x)$  be its distance from the closest center and  $w_2(x)$  be its distance from the second closest center.
- Suppose OPT is the cost of the optimal clustering.
- If the dataset satisfies  $(1 + \alpha, \epsilon)$ -approximation-stability for the  $k$ -means objective, then
  - 1 If  $\forall i, |C_i^*| \geq 2\epsilon n$ , then less than  $\epsilon n$  points have  $w_2^2(x) - w^2(x) \leq \frac{\alpha \cdot \text{OPT}}{\epsilon n}$ .
  - 2 For any  $t > 0$ , at most  $t\epsilon n$  points have  $w^2(x) \geq \frac{\text{OPT}}{t\epsilon n}$ .

# Preliminaries

- Let  $c_1, \dots, c_i$  be the centers chosen by the first  $i$  iterations of  $k$ -means++.
- Suppose  $j_1, \dots, j_i$  are the indices of the optimal clusters to which these centers belong.
- Define  $J_i = \{j_1, \dots, j_i\}$  and  $\bar{J}_i = \{1, \dots, k\} \setminus J_i$ .
- $J_i$  is the set of indices of the clusters that are covered at the end of the  $i^{\text{th}}$  iteration.



- Let  $B_1$  be the subset of points in  $\bar{\mathcal{X}}_i$  such that for any point  $x \in B_1$ ,  $w_2^2(x) - w^2(x) \leq \frac{\alpha \cdot \text{OPT}}{\epsilon n}$ .
- Let  $B_2$  denote the subset of points in  $\bar{\mathcal{X}}_i$  such that for every point  $x \in B_2$ ,  $w^2(x) \geq \frac{\alpha^2 \cdot \text{OPT}}{6\epsilon n}$ .
- We know that  $|B_1| \leq \epsilon n$  and  $|B_2| \leq \frac{6\epsilon n}{\alpha^2}$ .
- Let  $B = B_1 \cup B_2$  and  $\bar{B} = \bar{\mathcal{X}}_i \setminus B$ .
- We know that  $|B| \leq \frac{7\epsilon n}{\alpha^2}$ .

## A key lemma

### Lemma

Let  $\beta = \frac{1-\frac{\alpha}{2}}{6+\alpha}$ . For any  $x \in \bar{B}$  we have,  $D^2(x, c_t) \geq \beta \cdot D^2(x, c_{j_t}^*)$ .

- Proof: Let  $j$  be the index of the optimal cluster to which  $x$  belongs.
- Note that  $w^2(x) = D^2(x, c_j^*)$  and  $w_2^2(x) \leq D^2(x, c_{j_t}^*)$ .
- For any  $x \in \bar{B}$ , we have:

$$\begin{aligned} w_2^2(x) - w^2(x) &\geq \frac{\alpha \cdot \text{OPT}}{\epsilon n} \geq \frac{6w^2(x)}{\alpha} \\ \Rightarrow w_2^2(x) &\geq \left(1 + \frac{6}{\alpha}\right) \cdot w^2(x) \end{aligned} \quad (1)$$

- Suppose that  $D^2(x, c_t) < \beta \cdot D^2(x, c_{j_t}^*)$ .

- Then we get the following inequalities.

$$2 \cdot D^2(x, c_j^*) + 2 \cdot D^2(x, c_t) \geq D^2(c_t, c_j^*) \quad (\Delta \text{ inequality})$$

$$\Rightarrow 2 \cdot D^2(x, c_j^*) + 2 \cdot D^2(x, c_t) \geq D^2(c_t, c_{j_t}^*) \quad (D^2(c_t, c_j^*) \geq D^2(c_t, c_{j_t}^*))$$

$$\Rightarrow 2 \cdot D^2(x, c_j^*) + 2 \cdot D^2(x, c_t) \geq \frac{1}{2} \cdot D^2(x, c_{j_t}^*) - D^2(x, c_t)$$

$$\Rightarrow 3 \cdot D^2(x, c_t) \geq \frac{1}{2} \cdot D^2(x, c_{j_t}^*) - 2 \cdot D^2(x, c_j^*)$$

$$\Rightarrow 3\beta \cdot D^2(x, c_{j_t}^*) > \frac{1}{2} \cdot D^2(x, c_{j_t}^*) - 2 \cdot D^2(x, c_j^*)$$

(using assumption  $D^2(x, c_t) < \beta \cdot D^2(x, c_{j_t}^*)$ )

$$\Rightarrow D^2(x, c_j^*) > \frac{1 - 6\beta}{4} \cdot D^2(x, c_{j_t}^*)$$

$$\Rightarrow w^2(x) > \frac{1}{1 + \frac{6}{\alpha}} \cdot w_2^2(x) \quad (D^2(x, c_{j_t}^*) \geq w_2^2(x) \text{ and } \beta = \frac{1 - \frac{\alpha}{2}}{6 + \alpha})$$

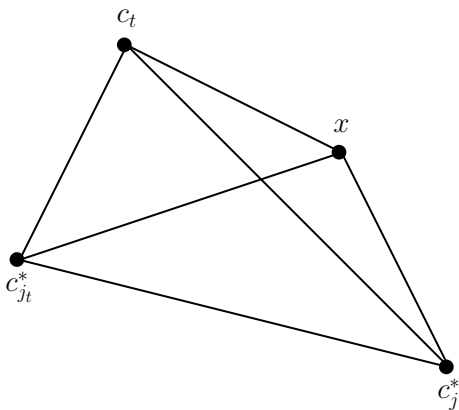


Figure:  $x$  belongs to the uncovered cluster  $j$ .

- This contradicts with Equation (1). Hence, for any  $x \in \bar{B}$  and any  $t \in \{1, \dots, i\}$ , we have  $D^2(x, c_t) \geq \beta \cdot D^2(x, c_{j_t}^*)$ .

- Let  $W_{min} = \min_{t \in [k]} \left( \sum_{x \in C_t^*, x \in \bar{B}} w_2^2(x) \right)$ .
- Let  $C_i$  denote the set of centers  $\{c_1, \dots, c_i\}$  that are chosen in the first  $i$  iterations of  $k$ -means++.
- Let  $\mathcal{X}_i = \cup_{t \in J_i} C_t^*$  and  $\bar{\mathcal{X}}_i = \mathcal{X} \setminus \mathcal{X}_i$ .
- $\mathcal{X}_i$  denotes the points that are covered by the algorithm after step  $i$ .
- For any subset of points  $Y \subseteq \mathcal{X}$ ,  $\phi_{C_i}(Y)$  is the cost of the points in  $Y$  with respect to the centers  $C_i$ , i.e.,  

$$\phi_{C_i}(Y) = \sum_{x \in Y} \min_{c \in C_i} D^2(x, c).$$
- We have  $\phi_{\{c_1, \dots, c_i\}}(\bar{\mathcal{X}}_i) \geq \beta \cdot (k - i) \cdot W_{min}$ .

- Let  $E_i$  denote the event that the set  $J_i$  contains  $i$  distinct indices from  $\{1, \dots, k\}$ .
- This means that the first  $i$  sampled centers cover  $i$  optimal clusters.
- The next Lemma is from [AV07] and shows that given that event  $E_i$  happens, the expected cost of points in  $\mathcal{X}_i$  with respect to  $C_i$  is at most some constant times the optimal cost of  $\mathcal{X}_i$  with respect to  $\{c_1^*, \dots, c_k^*\}$ .
- $\forall i, \mathbf{E}[\phi_{\{c_1, \dots, c_i\}}(\mathcal{X}_i) | E_i] \leq 4 \cdot \phi_{\{c_1^*, \dots, c_k^*\}}(\mathcal{X}_i)$ .

- From the last lemma, we get

$$\Pr \left[ \phi_{\{c_1, \dots, c_k\}}(\mathcal{X}) \leq 8 \cdot \phi_{\{c_1^*, \dots, c_k^*\}}(\mathcal{X}) \right] \geq \frac{1}{2} \Pr[E_k].$$

- We also show that  $\Pr[E_{i+1}|E_i] \geq \frac{k-i}{k-i+1}$ .
- This gives  $\Pr[E_k] \geq \frac{1}{k}$ .
- Hence,  $\Pr \left[ \phi_{\{c_1, \dots, c_k\}}(\mathcal{X}) \leq 8 \cdot \phi_{\{c_1^*, \dots, c_k^*\}}(\mathcal{X}) \right] \geq \frac{1}{2k}$ .
- Thus, the  $k$ -means++ algorithm gives an 8-approximation to the  $k$ -means objective with probability  $\Omega(\frac{1}{k})$ .

## Conclusion and future work

- In this work, we showed that the  $k$ -means++ algorithm gives a constant factor approximation to the  $k$ -means and  $k$ -median objective with probability  $\Omega(\frac{1}{k})$ , provided all the clusters are large.
- We also showed that for small clusters, there is a dataset on which  $k$ -means++ can't achieve a constant factor approximation.
- Can we improve the upper and lower bounds in the analysis?